

18-791 Lecture #18
FAST FOURIER TRANSFORM INVERSES
AND ALTERNATE IMPLEMENTATIONS

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Introduction

- In our lecture last Tuesday we described and discussed the basic decimation-in-time Cooley-Tuckey fast Fourier transform algorithm for DFT sizes that are integer powers of 2 (radix 2)
- Today we will discuss some variations and extensions of the basic FFT algorithm:
 - Alternate forms of the FFT structure
 - Computation of the inverse DFT
 - The decimation-in-frequency FFT algorithm
 - FFT structures for DFT sizes that are not an integer power of 2

Alternate FFT structures

- We developed the basic decimation-in-time (DIT) FFT structure in the last lecture, but other forms are possible simply by rearranging the branches of the signal flowgraph
- Consider the rearranged signal flow diagrams on the following panels

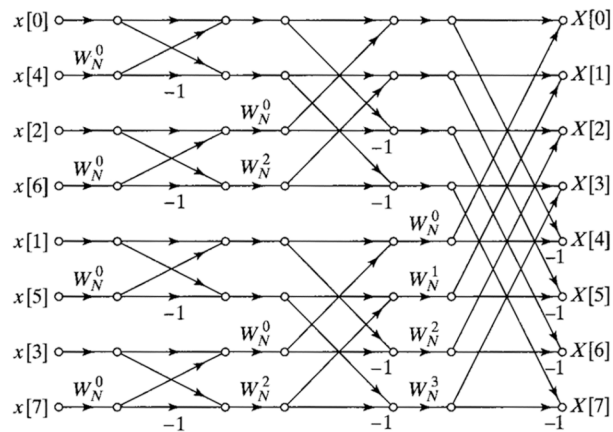


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Alternate DIT FFT structures (continued)

- DIT structure with input bit-reversed, output natural (OSB 9.10)

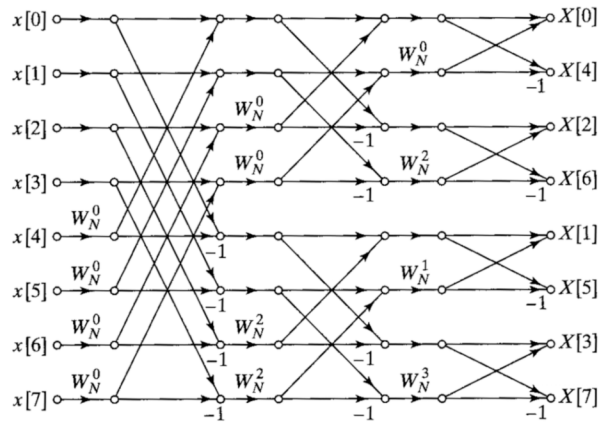


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Alternate DIT FFT structures (continued)

- DIT structure with input natural, output bit-reversed (OSB 9.14)

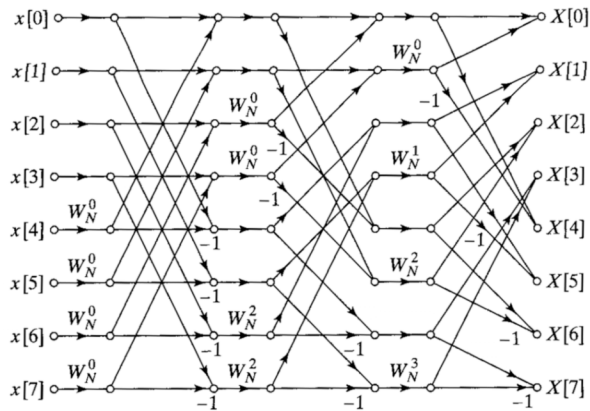


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Alternate DIT FFT structures (continued)

- DIT structure with both input and output natural (OSB 9.15):

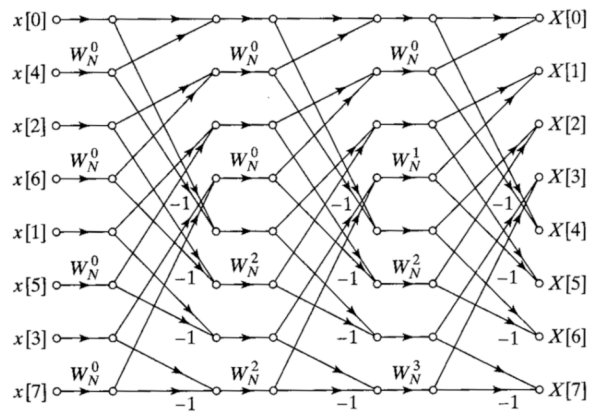


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Alternate DIT FFT structures (continued)

■ DIT structure with same structure for each stage (OSB 9.16):



Comments on alternate FFT structures

■ A method to avoid bit-reversal in filtering operations is:

- Compute forward transform using natural input, bit-reversed output (as in OSB 9.10)
- Multiply DFT coefficients of input and filter response (both in bit-reversed order)
- Compute inverse transform of product using bit-reversed input and natural output (as in OSB 9/14)

■ Latter two topologies (as in OSB 9.15 and 9.16) are now rarely used

Using FFTs for inverse DFTs

- We've always been talking about forward DFTs in our discussion about FFTs what about the inverse FFT?

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}$$

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j2\pi kn/N}$$

- One way to modify FFT algorithm for the inverse DFT computation is:
 - Replace $e^{-j2\pi kn/N}$ by $e^{j2\pi kn/N}$ wherever it appears
 - Multiply final output by $1/N$
- This method has the disadvantage that it requires modifying the internal code in the FFT subroutine

A better way to modify FFT code for inverse DFTs

- Taking the complex conjugate of both sides of the IDFT equation and multiplying by N :

$$N x^*[n] = \sum_{k=0}^{N-1} X^*[k] e^{-j2\pi kn/N}$$

$$X[k] = \sum_{n=0}^{N-1} N x^*[n] e^{-j2\pi kn/N}$$

- This suggests that we can modify the FFT algorithm for the inverse DFT computation by the following:
 - Complex conjugate the input DFT coefficients
 - Compute the *forward* FFT
 - Complex conjugate the output of the FFT and multiply by $1/N$
- This method has the advantage that the internal FFT code is undisturbed; it is widely used.

The decimation-in-frequency (DIF) FFT algorithm

- Introduction: Decimation in frequency is an alternate way of developing the FFT algorithm
- It is different from decimation in time in its development, although it leads to a very similar structure

The decimation in frequency FFT (continued)

- Consider the original DFT equation

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}$$

- Separate the first half and the second half of time samples:

$$X[k] = \sum_{n=0}^{N/2-1} x[n] e^{-j2\pi kn/N} + \sum_{n=N/2}^{N-1} x[n] e^{-j2\pi kn/N}$$

$$= \sum_{n=0}^{N/2-1} x[n] e^{-j2\pi kn/N} + \sum_{n=0}^{N/2-1} x[n+N/2] e^{-j2\pi k(n+N/2)/N}$$

$$= \sum_{n=0}^{N/2-1} x[n] e^{-j2\pi kn/N} + \sum_{n=0}^{N/2-1} x[n+N/2] e^{-j2\pi kn/N} e^{-j2\pi kN/2N}$$

$$= \sum_{n=0}^{N/2-1} (x[n] + (-1)^k x[n+N/2]) e^{-j2\pi kn/N}$$

- Note that these are not $N/2$ -point DFTs

Continuing with decimation in frequency ...



- For k even, let $k = 2r$



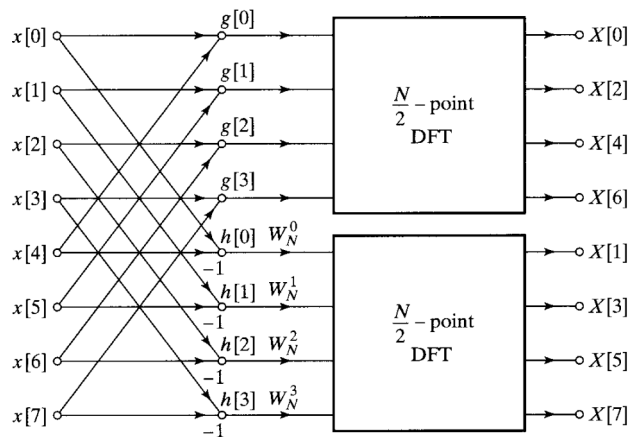
- For k odd, let $k = 2r + 1$



- These expressions are the $N/2$ -point DFTs of



These equations describe the following structure:



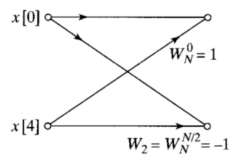
The DIF FFT is the *transpose* of the DIT FFT

- To obtain flowgraph transposes:
 - Reverse direction of flowgraph arrows
 - Interchange input(s) and output(s)

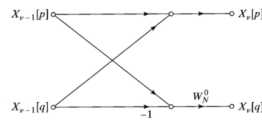
- Reverse direction of flowgraph arrows

- Interchange input(s) and output(s)

- DIT butterfly:



- DIF butterfly:



- Comment:

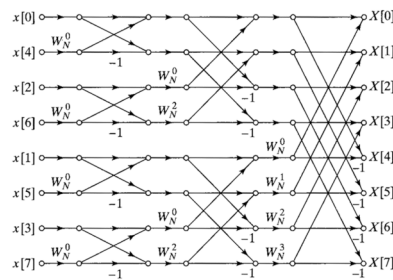
- We will revisit transposed forms again in our discussion of filter implementation



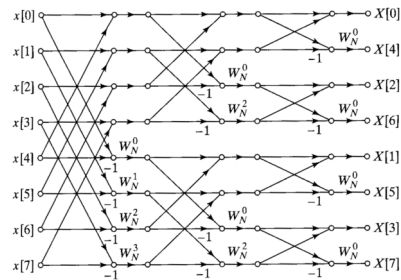
The DIF FFT is the *transpose* of the DIT FFT

- Comparing DIT and DIF structures:

- DIT FFT structure:



- DIF FFT structure:

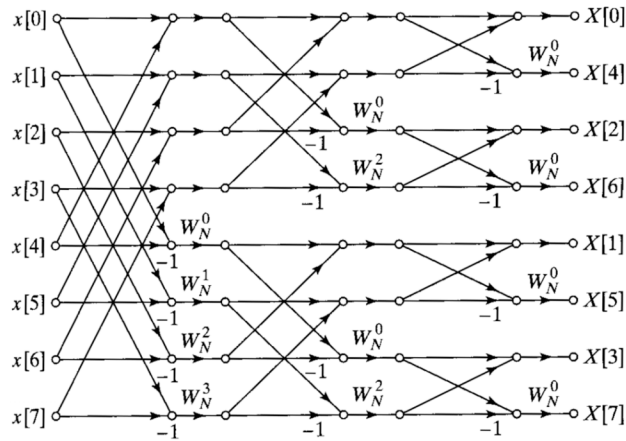


- Alternate forms for DIF FFTs are similar to those of DIT FFTs



Alternate DIF FFT structures

- DIF structure with input natural, output bit-reversed (OSB 9.20)

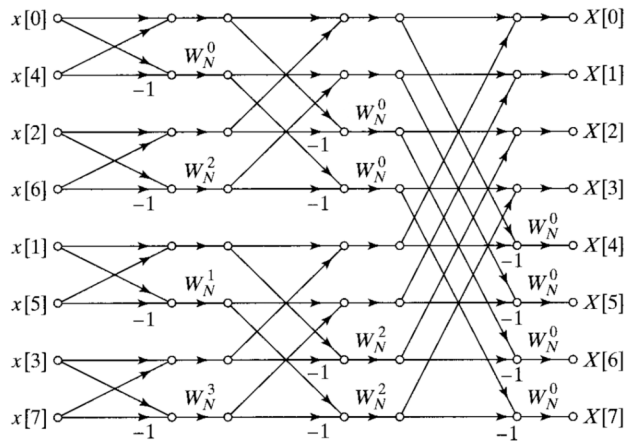


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Alternate DIF FFT structures (continued)

- DIF structure with input bit-reversed, output natural (OSB 9.22)

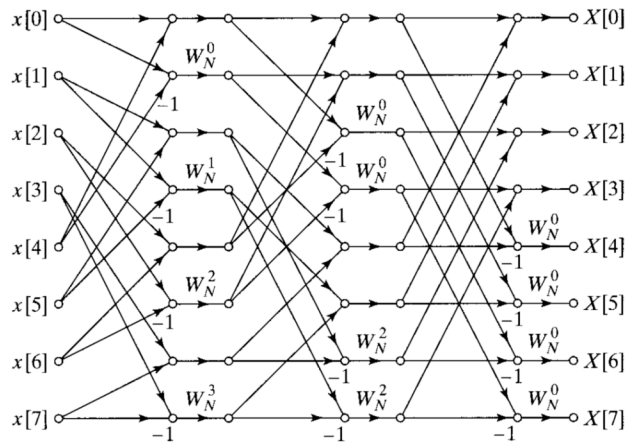


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Alternate DIF FFT structures (continued)

- DIF structure with both input and output natural (OSB 9.23):



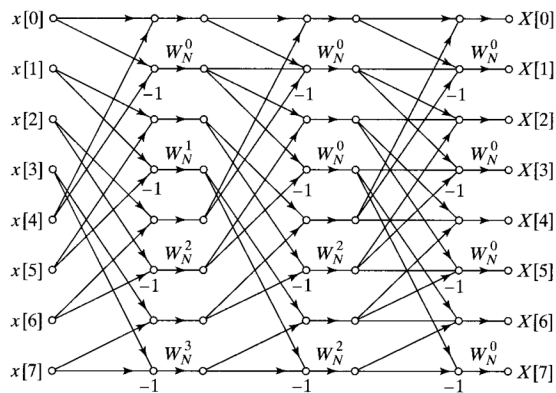
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Alternate DIF FFT structures (continued)

- DIF structure with same structure for each stage (OSB 9.24):



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FFT structures for other DFT sizes

- Can we do anything when the DFT size N is not an integer power of 2 (the non-radix 2 case)?
- Yes! Consider a value of N that is not a power of 2, but that still is highly factorable ...

$$x[n] = \sum_{k=0}^{N-1} X[k] e^{j2\pi kn/N}$$

- Then let

$$x[n] = \sum_{k=0}^{N-1} X[k] e^{j2\pi kn/N} = \sum_{k=0}^{N-1} X[k] e^{j2\pi k n / (N_1 N_2)} = \sum_{k=0}^{N-1} X[k] e^{j2\pi k n_1 n_2 / (N_1 N_2)}$$



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Non-radix 2 FFTs (continued)

- An arbitrary term of the sum on the previous panel is

$$X[k] e^{j2\pi k n_1 n_2 / (N_1 N_2)}$$

- This is, of course, a DFT of size N_1 of points spaced by N_2



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Non-radix 2 FFTs (continued)

- In general, for the first decomposition we use

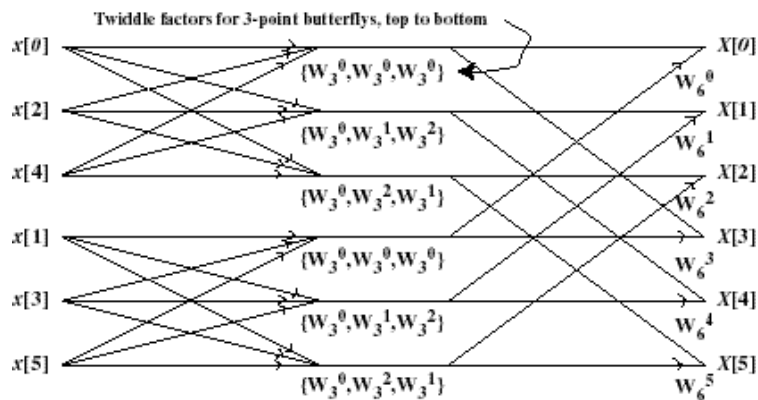


- **Comments:**

- This procedure can be repeated for subsequent factors of N
- The amount of computational savings depends on the extent to which N is “composite”, able to be factored into small integers
- Generally the smallest factors possible used, with the exception of some use of radix-4 and radix-8 FFTs

An example The 6-point DIT FFT

- $P_1 = 2; P_2 = 3;$



Summary

- **This morning we considered a number of alternative ways of computing the FFT:**
 - Alternate implementation structures
 - The decimation-in-frequency structure
 - FFTs for sizes that are non-integer powers of 2
 - Using standard FFT structures for inverse FFTs

- **Starting on Tuesday we will begin to discuss digital filter implementation structures**