

18734: Foundations of Privacy

Fairness in Classification

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With many slides from Moritz Hardt

Fall 2015

Fairness in Classification

Advertising 

Education 

Financial aid

Health

Banking

Care 

Insurance



Taxation

many more...

Concern: Discrimination

- Certain attributes should be *irrelevant!*
- Population includes minorities
 - Ethnic, religious, medical, geographic
- Protected by law, policy, ethics



“Big Data: Seizing Opportunities, Preserving Values”

THE 90-DAY REVIEW
FOR BIG DATA



"big data technologies can cause societal harms beyond damages to privacy"

Other notions of “fairness” in CS

- Fair scheduling
- Distributed computing
- Envy-free division (cake cutting)
- Stable matching





Discrimination arises even when nobody's *evil*

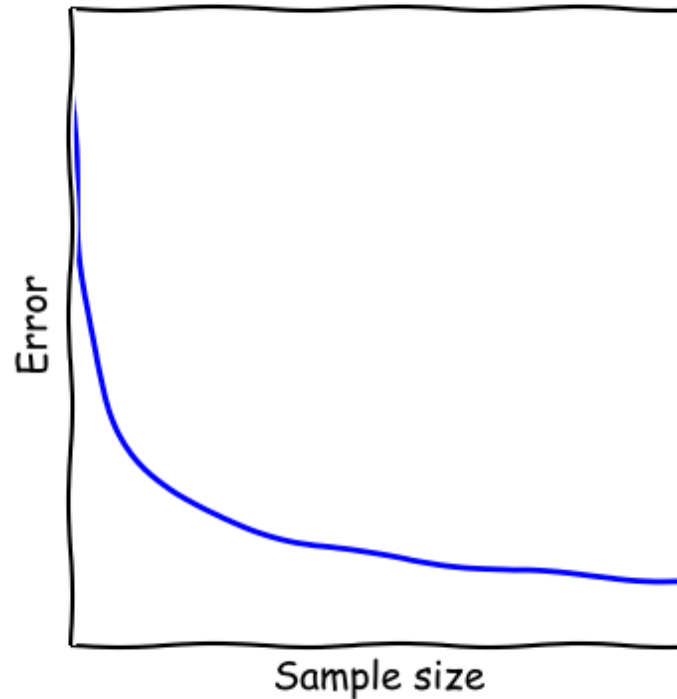
- Google+ tries to classify real vs fake names
- Fairness problem:
 - Most training examples standard white American names: John, Jennifer, Peter, Jacob, ...
 - Ethnic names often unique, much fewer training examples

Likely outcome: Prediction accuracy
worse on ethnic names

“Due to Google's ethnocentricity I was prevented from using my real last name (my nationality is: Tungus and Sami)”

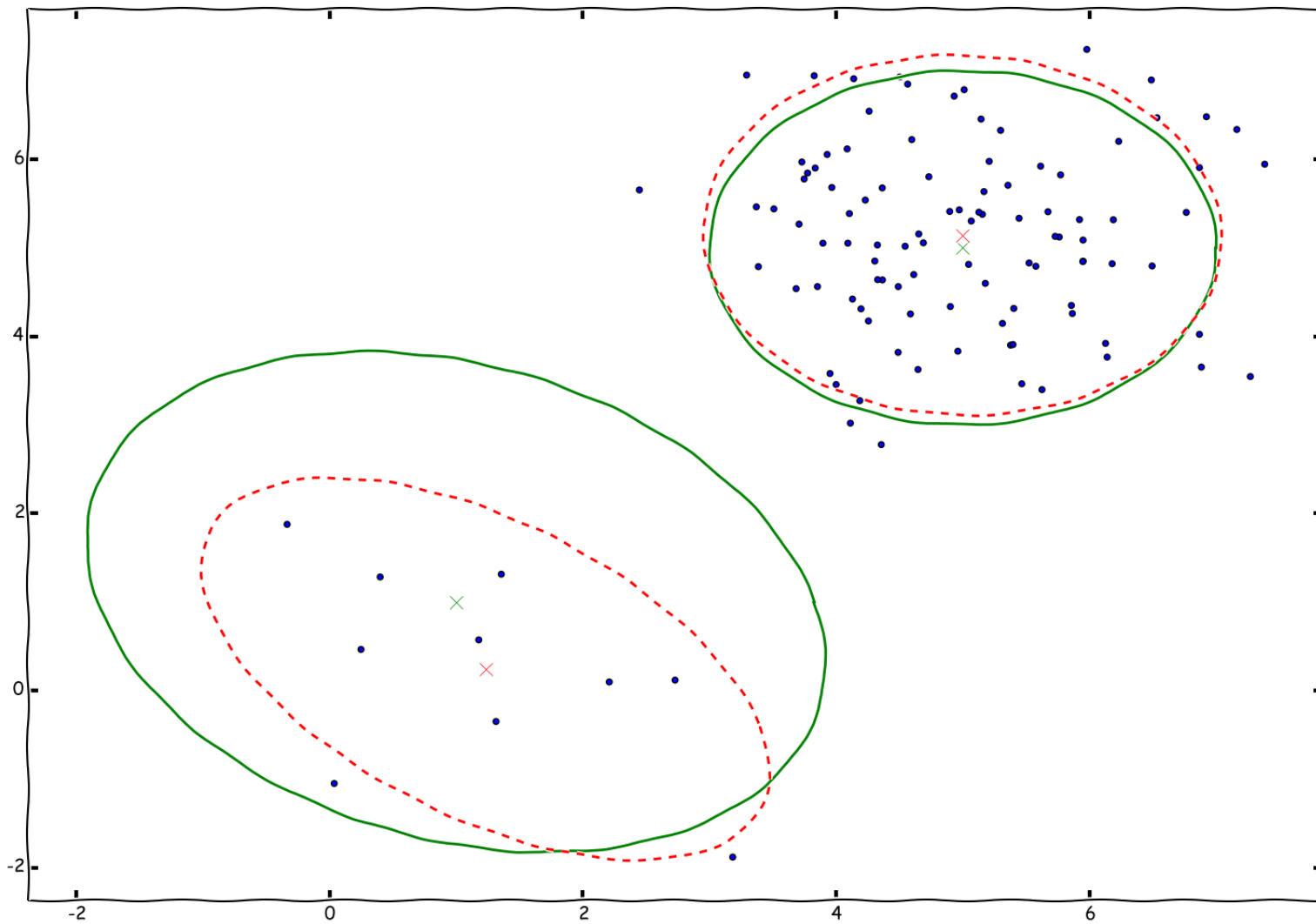
- Katya Casio. Google Product Forums.

Error vs sample size



Sample Size Disparity:

In a heterogeneous population, smaller groups face larger error



Credit Application



**More miles
and no annual fee**

Earn trips faster with VentureOneSM

Get Started 

only at **Capital One CARD LAB**

Capital One Card Lab
Platinum Prestige Credit Card

Capital One Card Lab
VentureOne Card

Savings Accounts
Earn With Great Rates

The advertisement features a yellow Capital One VentureOne Visa Signature credit card. The card displays the name 'VENTURE', the number '4000 1234 5678 9010', the expiration date '12/12', and the name 'DER'. The card is set against a background of a tropical island with palm trees and a blue sky. The text 'More miles and no annual fee' is prominently displayed in blue and orange. Below this, it says 'Earn trips faster with VentureOneSM'. A green button with a white arrow points to the right, labeled 'Get Started'. At the bottom, there are three white boxes with blue borders containing text: 'only at Capital One CARD LAB', 'Capital One Card Lab Platinum Prestige Credit Card', 'Capital One Card Lab VentureOne Card', and 'Savings Accounts Earn With Great Rates'.

User visits `capitalone.com`

Capital One uses tracking information provided by the tracking network [x+1] to personalize offers

Concern: Steering minorities into higher rates (illegal)

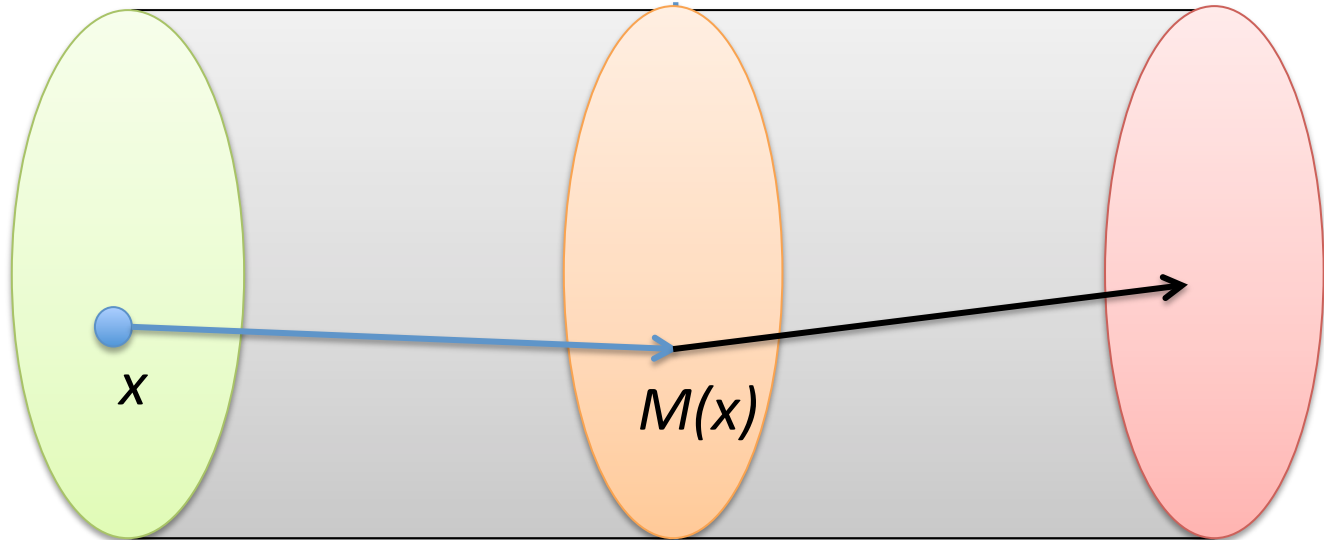
WSJ 2010

Classifier
(eg. ad network)

$$M: V \rightarrow O$$

Vendor
(eg. capital one)

$$f: O \rightarrow A$$



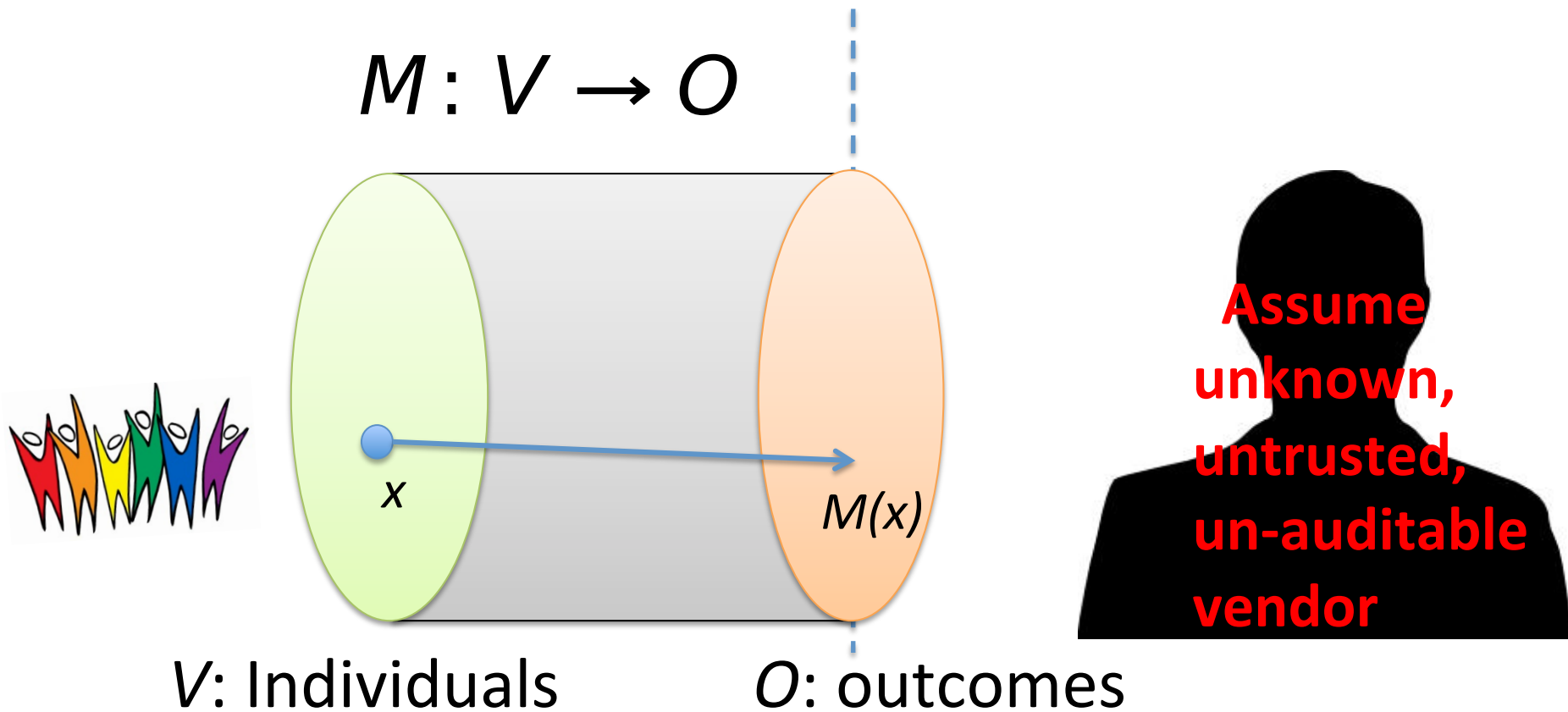
V: Individuals

O: outcomes

A: actions

Goal:

Achieve Fairness in the classification step



First attempt...

Fairness through Blindness



Fairness through Blindness

Ignore all irrelevant/protected attributes

“We don’t even look at ‘race’!”

Point of Failure

You don't need to *see* an attribute to be able to *predict* it with high accuracy

E.g.: User visits `artofmanliness.com`
... 90% chance of being male

Fairness through Privacy?

“It's Not Privacy, and It's Not Fair”

Cynthia Dwork & Deirdre K. Mulligan. Stanford Law Review.

Privacy is no Panacea: Can't hope to have privacy solve our fairness problems.

“At worst, **privacy solutions can hinder efforts to identify classifications that unintentionally produce objectionable outcomes**—for example, differential treatment that tracks race or gender—by limiting the availability of data about such attributes.”

Second attempt...

Statistical Parity (Group Fairness)

Equalize two groups S , T at the level of outcomes

– E.g. $S = \text{minority}$, $T = S^c$

$$\Pr[\text{outcome } o \mid S] = \Pr[\text{outcome } o \mid T]$$

“Fraction of people in S getting credit same as in T .”

Not strong enough as a notion of fairness

– Sometimes desirable, but can be abused

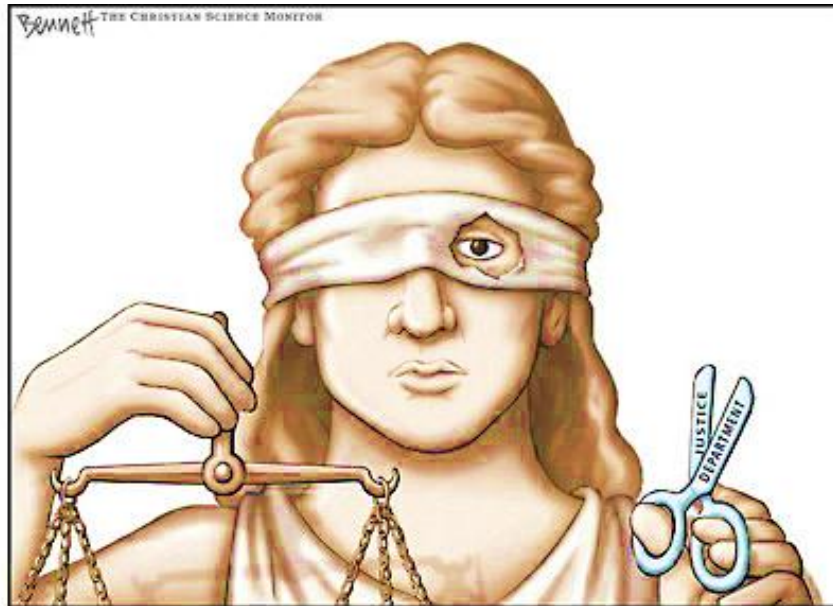
- **Self-fulfilling prophecy:** Select smartest students in T , random students in S

– *Students in T will perform better*

Lesson: Fairness is *task-specific*

Fairness requires understanding of classification task and protected groups

“Awareness”



Individual Fairness Approach

Individual Fairness

Treat *similar* individuals *similarly*



Similar for the purpose of
the classification task



Similar distribution
over outcomes



The Similarity Metric

Metric

- Assume *task-specific similarity metric*
 - Extent to which two individuals are similar w.r.t. the classification task at hand
- Ideally captures *ground truth*
 - Or, society's best approximation
- Open to public discussion, refinement
 - In the spirit of Rawls
- Typically, does not suggest classification!

Examples

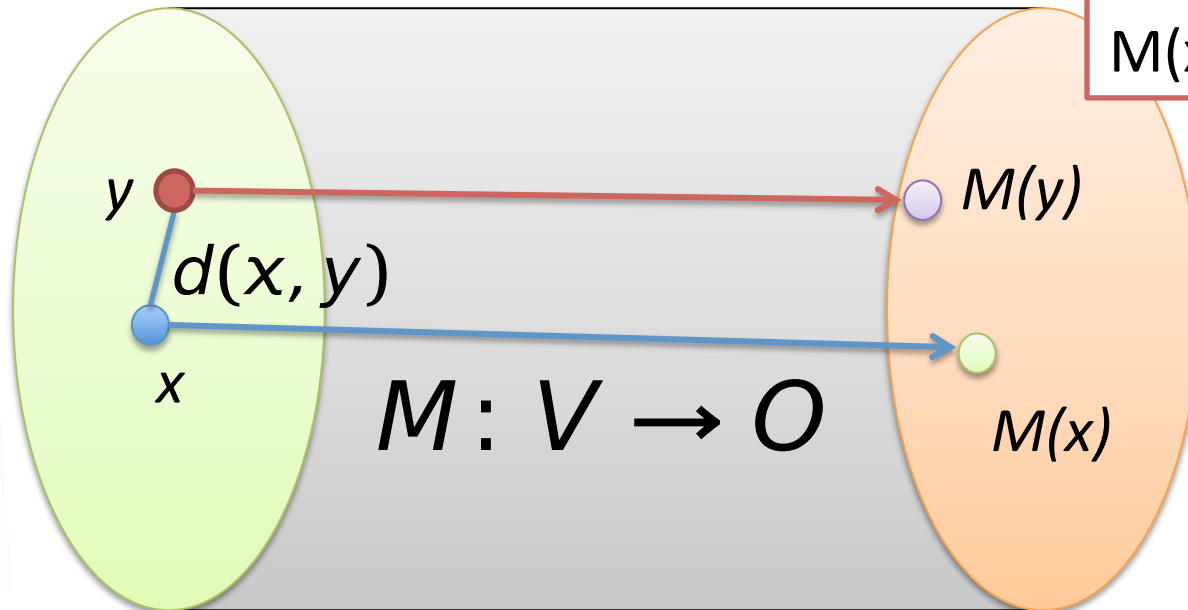
- Financial/insurance risk metrics
 - Already widely used (though secret)
- **AALIM health care metric**
 - health metric for treating similar patients similarly
- Roemer's relative effort metric
 - Well-known approach in Economics/Political theory

Maybe not so much science fiction after all...

How to formalize this?

Think of V as space
with metric $d(x,y)$
similar = small $d(x,y)$

How can we
compare
 $M(x)$ with $M(y)$?

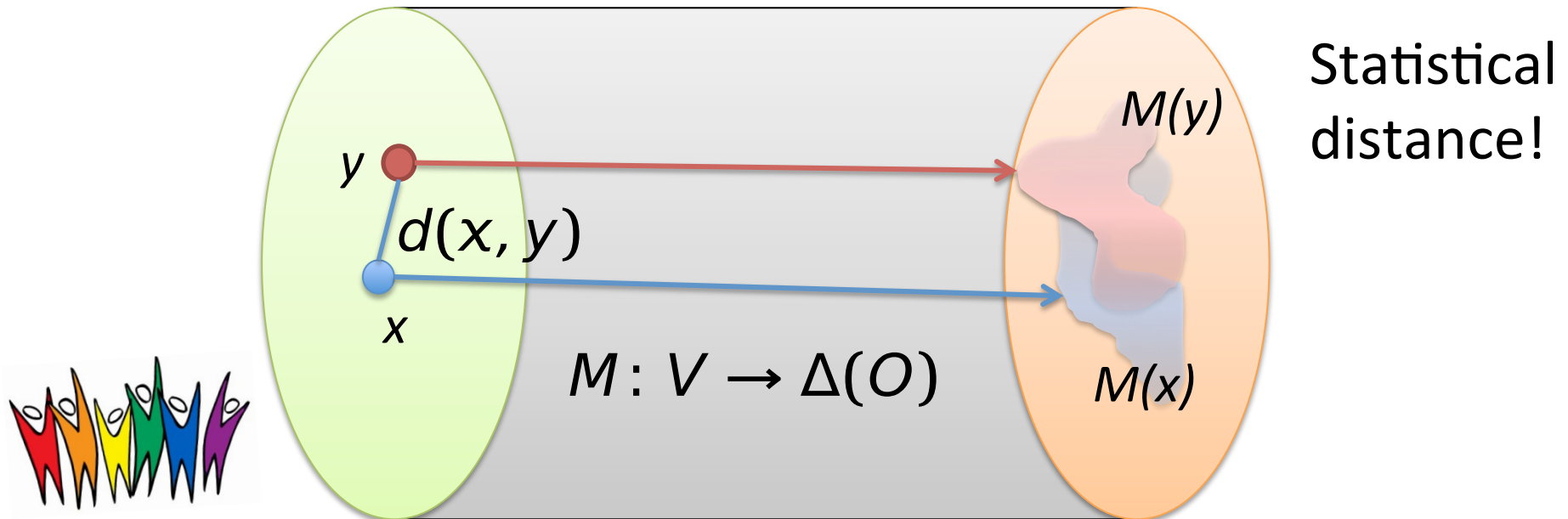


V : Individuals

O : outcomes

Distributional outcomes

How can we compare $M(x)$ with $M(y)$?



V : Individuals

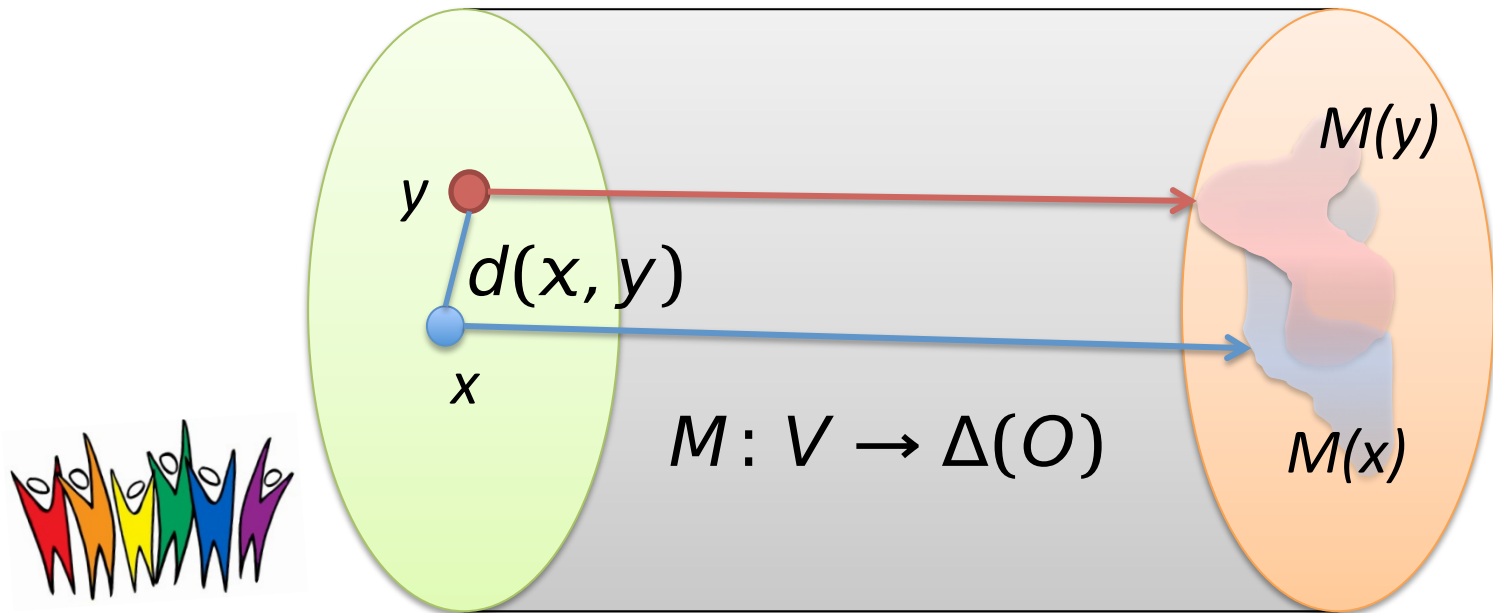
O : outcomes

Metric $d: V \times V \rightarrow \mathbb{R}$

Lipschitz condition $\|M(x) - M(y)\| \leq d(x, y)$

This talk: Statistical distance

in $[0,1]$



V : Individuals

O : outcomes

Statistical Distance

P, Q denote probability measures on a finite domain A . The *statistical distance* between P and Q is denoted by

$$D_{\text{tv}}(P, Q) = \frac{1}{2} \sum_{a \in A} |P(a) - Q(a)|.$$

Notation match:

$$M(x) = P$$

$$M(y) = Q$$

$$O = A$$

Statistical Distance

P, Q denote probability measures on a finite domain A . The *statistical distance* between P and Q is denoted by

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Example: High D

$$A = \{0, 1\}$$

$$P(0) = 1, P(1) = 0$$

$$Q(0) = 0, Q(1) = 1$$

$$D(P, Q) = 1$$

Statistical Distance

P, Q denote probability measures on a finite domain A . The *statistical distance* between P and Q is denoted by

$$D_{\text{tv}}(P, Q) = \frac{1}{2} \sum_{a \in A} |P(a) - Q(a)|.$$

Example: Low D

$$A = \{0, 1\}$$

$$P(0) = 1, P(1) = 0$$

$$Q(0) = 1, Q(1) = 0$$

$$D(P, Q) = 0$$

Statistical Distance

P, Q denote probability measures on a finite domain A . The *statistical distance* between P and Q is denoted by

$$D_{\text{tv}}(P, Q) = \frac{1}{2} \sum_{a \in A} |P(a) - Q(a)|.$$

Example: Mid D

$$A = \{0, 1\}$$

$$P(0) = P(1) = \frac{1}{2}$$

$$Q(0) = \frac{3}{4}, Q(1) = \frac{1}{4}$$

$$D(P, Q) = \frac{1}{4}$$

Existence Proof

There exists a classifier that satisfies the Lipschitz condition

- Idea: Map all individuals to the same distribution over outcomes
- Are we done?

Key elements of approach...

Utility Maximization

Vendor can specify **arbitrary utility function**

$$U: V \times O \rightarrow \mathbb{R}$$

$U(v,o)$ = Vendor's utility of giving individual v
the outcome o

Maximize vendor's expected utility subject to Lipschitz condition

$$\max_{M(x)} \mathbb{E}_{x \sim V} \mathbb{E}_{o \sim M(x)} U(x, o)$$

s.t. M is d -Lipschitz

$$\|M(x) - M(y)\| \leq d(x, y)$$

Linear Program Formulation

- Objective function is linear
 - $U(x, o)$ is constant for fixed x, o
 - Distribution over V is known
 - $\{M(x)\}_{x \in V}$ are only variables to be computed
- Lipschitz condition is linear when using statistical distance
- Linear program can be solved efficiently

Discrimination Harms

Information use

- Explicit discrimination
 - Explicit use of race/gender for employment
- Redundant encoding/proxy attributes

Practices

- Redlining
- Self-fulfilling prophecy
- Reverse tokenism

When does Individual Fairness imply Group Fairness?

Suppose we enforce a metric d .

Question: Which *groups of individuals* receive (approximately) equal outcomes?

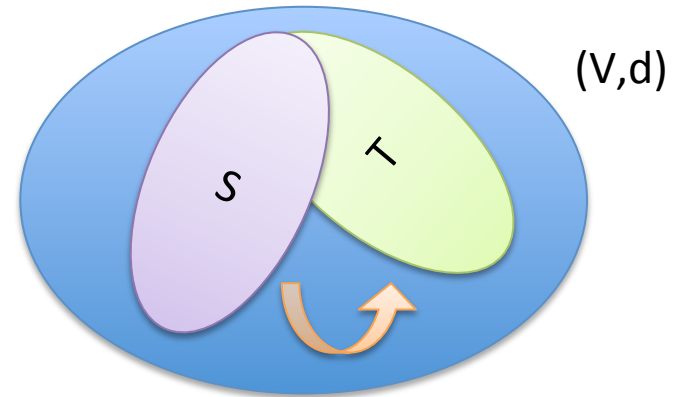
Theorem:

Answer is given by **Earthmover distance** (w.r.t. d) between the two groups.



How different are S and T ?

Earthmover Distance:
Cost of transforming
uniform distribution on S to
uniform distribution on T



$$\sigma_{EM}(S, T) \stackrel{\text{def}}{=} \min \sum_{x, y \in V} h(x, y) \sigma(x, y)$$

subject to

$$\sum_{y \in V} h(x, y) = S(x)$$
$$\sum_{y \in V} h(y, x) = T(x)$$
$$h(x, y) \geq 0$$

$$\sigma_{\text{EM}}(S, T) \stackrel{\text{def}}{=} \min \sum_{x, y \in V} h(x, y) \sigma(x, y)$$

subject to

$$\sum_{y \in V} h(x, y) = S(x)$$

$$\sum_{y \in V} h(y, x) = T(x)$$

$$h(x, y) \geq 0$$

$\text{bias}(d, S, T) =$ largest violation of statistical parity between S and T that any d -Lipschitz mapping can create

Theorem:

$$\text{bias}(d, S, T) = d_{\text{EM}}(S, T)$$

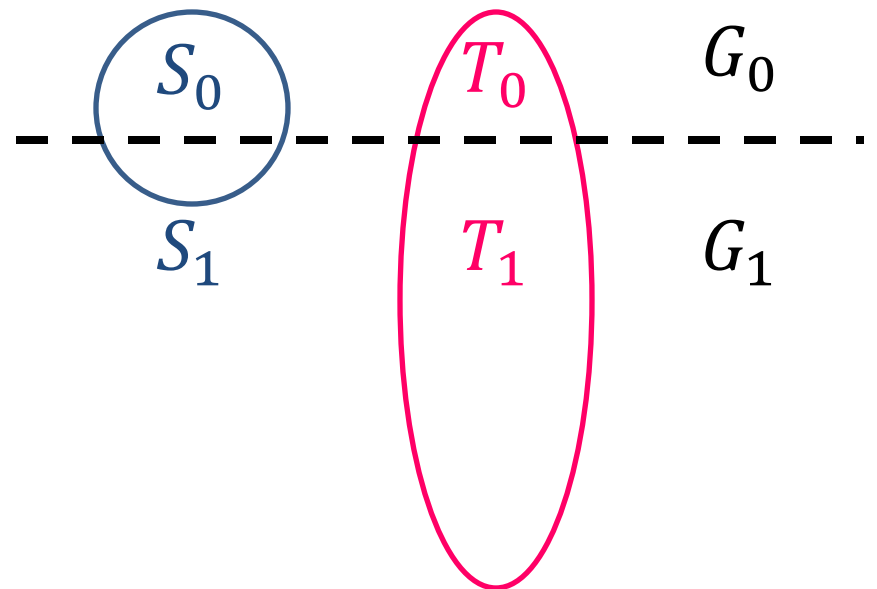


The Story So Far...

- Group fairness
- Individual fairness
- Group fairness does not imply individual fairness
- Individual fairness implies group fairness if earthmover distance small
- What if earthmover distance large?

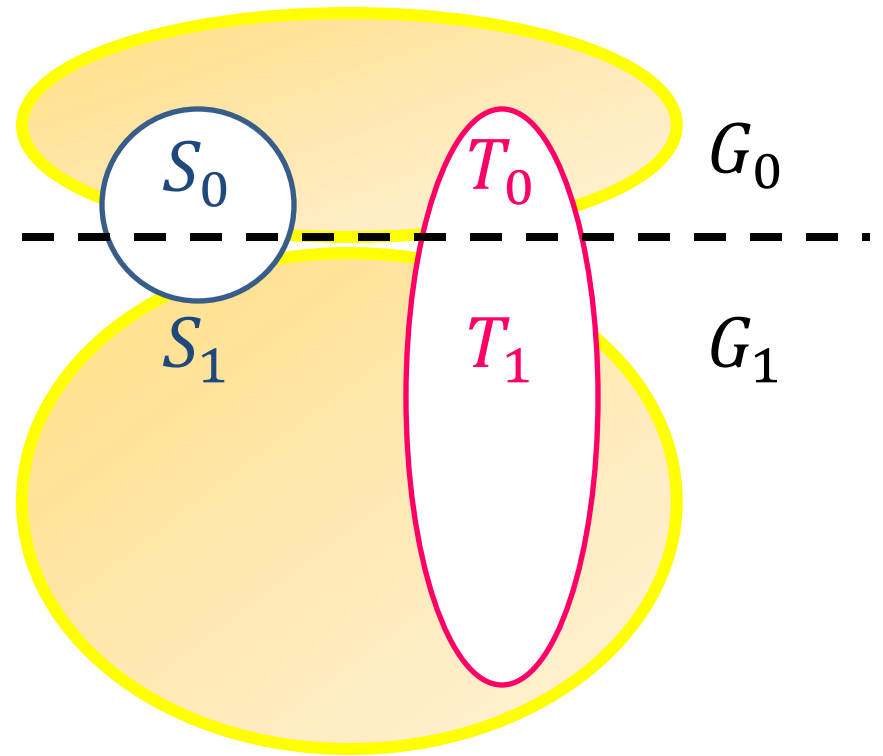
Toward Fair Affirmative Action: When $EM(S,T)$ is Large

- G_0 is unqualified
- G_1 is qualified



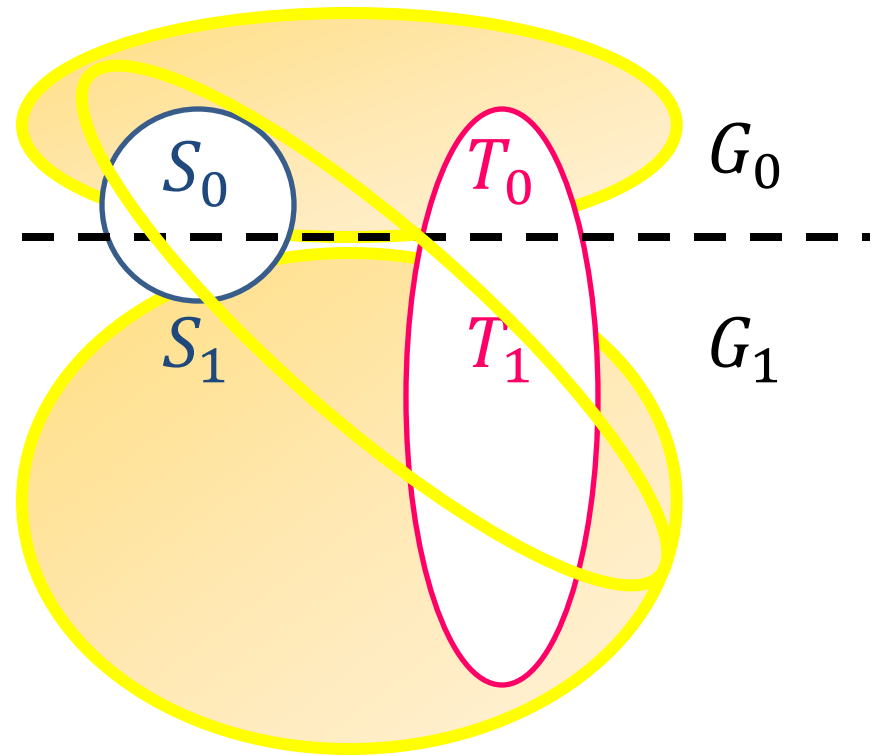
Toward Fair AA: When $EM(S,T)$ is Large

- Lipschitz \Rightarrow
All in G_i treated same



Toward Fair AA: When $EM(S,T)$ is Large

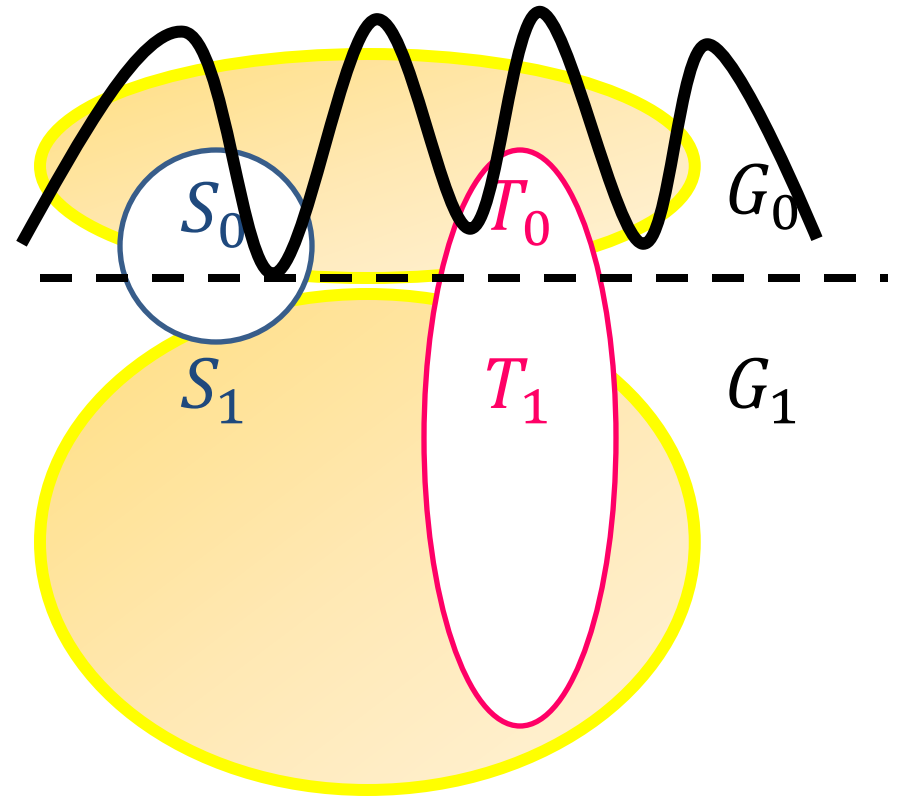
- Lipschitz \Rightarrow
All in G_i treated same
- **Statistical Parity** \Rightarrow
much of S_0 must be
treated the same as
much of T_1



Toward Fair AA: When $EM(S,T)$ is Large

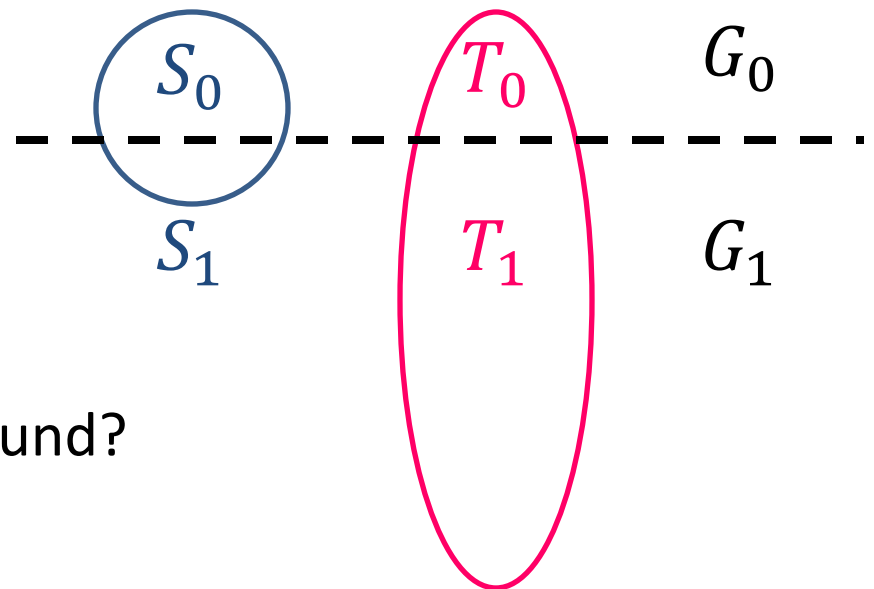
- Lipschitz \Rightarrow
All in G_i treated same

Failure to Impose Parity \Rightarrow
anti- S vendor can target G_0
with blatant hostile ad f_u .
Drives away almost all of S
while keeping most of T .



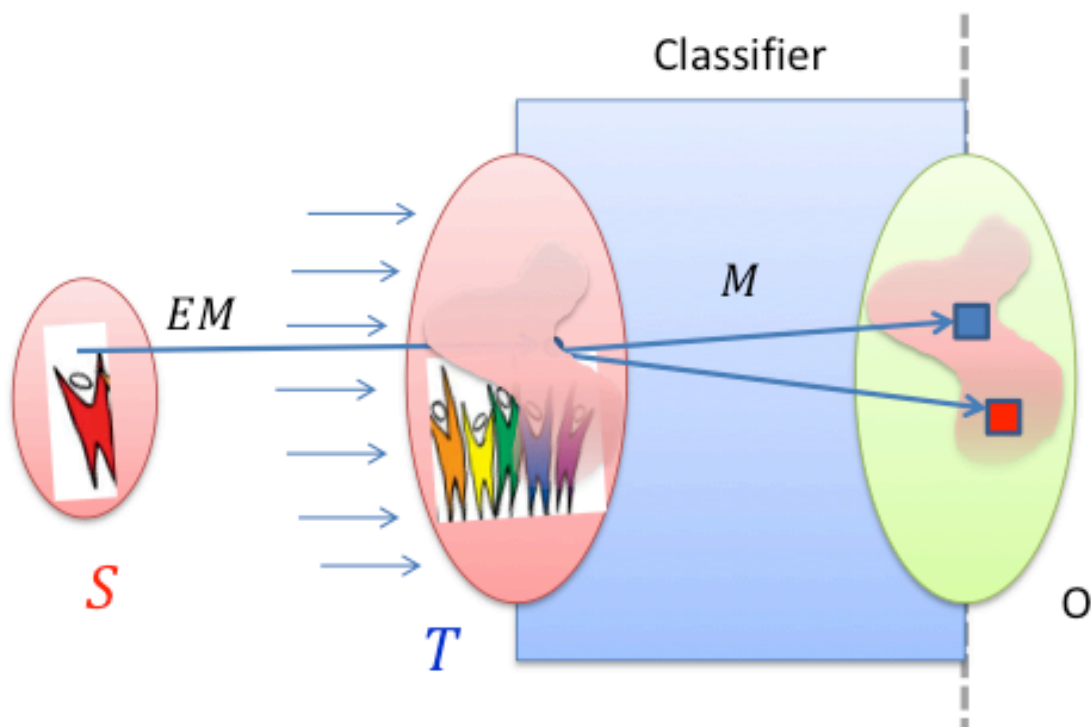
Dilemma: What to Do When $EM(S,T)$ is Large?

- Imposing parity causes collapse
- Failing to impose parity permits blatant discrimination



How can we form a middle ground?

Fair Affirmative Action

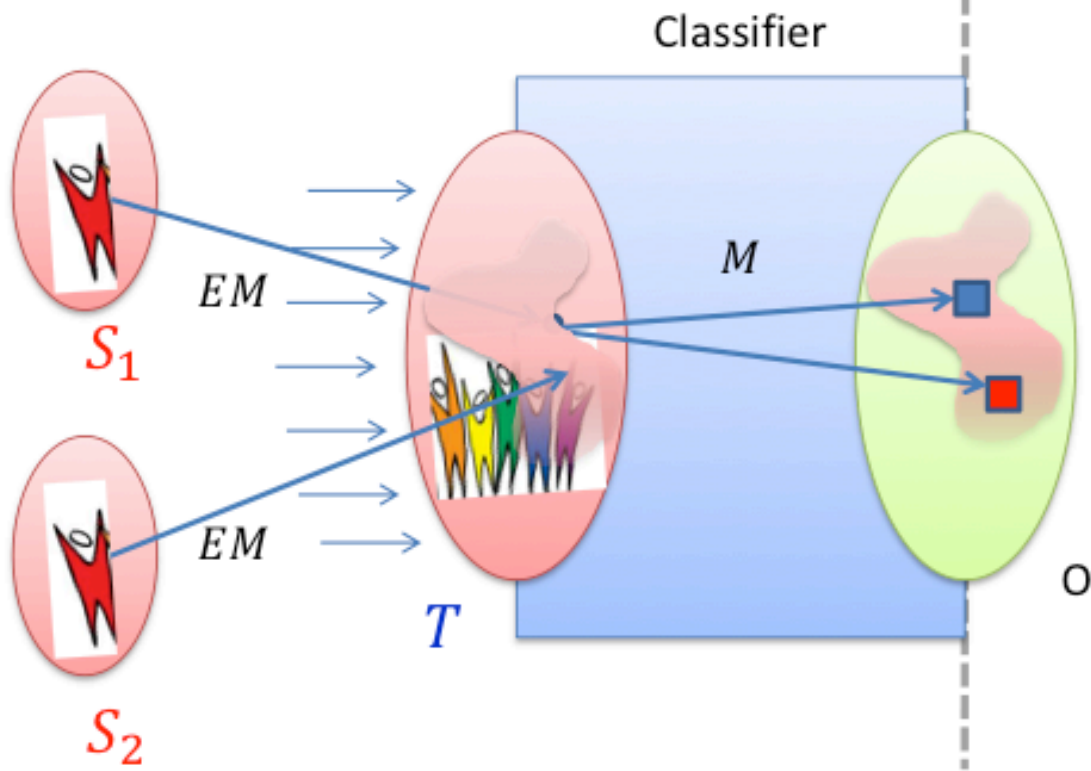


Earthmover mapping from S to T + Lipschitz mapping from T to O

Achieves:

- Lipschitz on $S \times S, T \times T$, on average on $S \times T$
- statistical parity between S and T
- no collapse

Fair Affirmative Action



- ▶ Immediately suggests a method of dealing with multiple disjoint S 's

Connection to differential privacy

- Close connection between individual fairness and **differential privacy** [Dwork-McSherry-Nissim-Smith'06]

DP: Lipschitz condition on set of databases

IF: Lipschitz condition on set of individuals

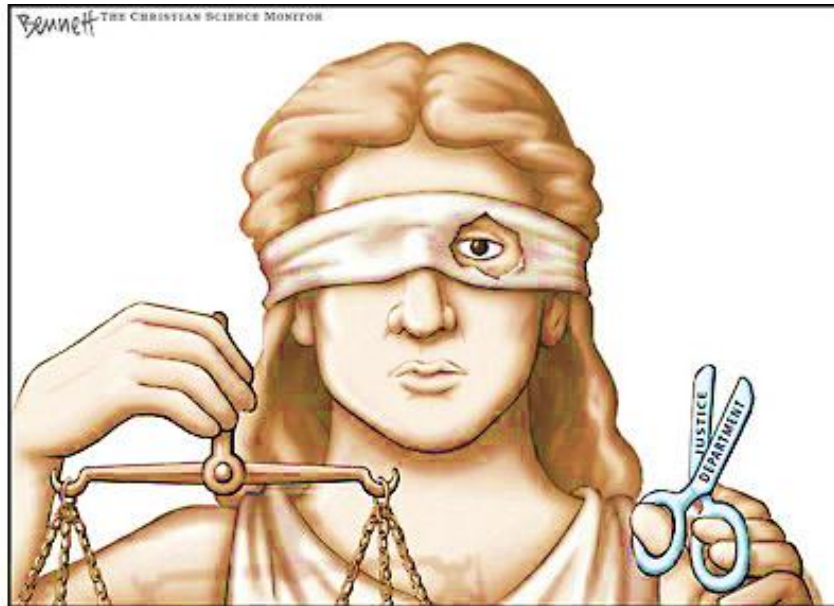
	Differential Privacy	Individual Fairness
Objects	Databases	Individuals
Outcomes	Output of statistical analysis	Classification outcome
Similarity	General purpose metric	Task-specific metric

Summary: Individual Fairness

- Formalized fairness property based on treating similar individuals similarly
 - Incorporates vendor's utility
- Explored relationship between individual fairness and group fairness
 - Earthmover distance
- Approach to fair affirmative action based on Earthmover solution

Lots of open problems/direction

- **Metric**
 - Social aspects, who will define them?
 - How to generate metric (semi-)automatically?
- **Earthmover characterization** when probability metric is not statistical distance (but infinity-div)
- Explore connection to **Differential Privacy**
- Connection to **Economics** literature/problems
 - Rawls, Roemer, Fleurbaey, Peyton-Young, Calsamiglia
- **Case Study**
- **Quantitative trade-offs** in concrete settings



Questions?

Metric

A metric on a set X is a function $d : X \times X \rightarrow \mathbb{R}^+$ (where \mathbb{R}^+ is the set of non-negative real numbers). For all x, y, z in X , this function is required to satisfy the following conditions:

- $d(x, y) \geq 0$ (non-negativity)
- $d(x, y) = 0$ if and only if $x = y$ (identity of indiscernibles.
Note that condition 1 and 2 together produce positive definiteness)
- $d(x, y) = d(y, x)$ (symmetry)
- $d(x, z) \leq d(x, y) + d(y, z)$ (subadditivity / triangle inequality).

Another Statistical Distance

$$D_{\infty}(P, Q) = \sup_{a \in A} \log \left(\max \left\{ \frac{P(a)}{Q(a)}, \frac{Q(a)}{P(a)} \right\} \right)$$

Partial Proof Idea

Theorem:

$$\text{bias}(d, S, T) \leq d_{EM}(S, T)$$

- $d_{EM}(S, T)$ cost of best coupling between the two distributions subject to the penalty function $d(x, y) = E d(x, y)$

Proof Sketch: LP Duality

- $EM_d(S,T)$ is an LP by definition
- Can write $\text{bias}(d,S,T)$ as an LP:

$\max \Pr(M(x) = 0 \mid x \text{ in } S) - \Pr(M(x) = 0 \mid x \text{ in } T)$

subject to:

- (1) $M(x)$ is a probability distribution for all x in V
- (2) M satisfies all d -Lipschitz constraints

Program dual to Earthmover LP!

Fair Affirmative Action (1)

- (a) First we compute a mapping from elements in S to distributions over T which transports the uniform distribution over S to the uniform distribution over T , while minimizing the total distance traveled. Additionally the mapping preserves the Lipschitz condition between elements within S .
 - (b) This mapping gives us the following new loss function for elements of T : For $y \in T$ and $a \in A$ we define a new loss, $L'(y, a)$, as

$$L'(y, a) = \sum_{x \in S} \mu_x(y) L(x, a) + L(y, a),$$

where $\{\mu_x\}_{x \in S}$ denotes the mapping computed in step (a). L' can be viewed as a reweighting of the loss function L , taking into account the loss on S (indirectly through its mapping to T).

2. Run the Fairness LP only on T , using the new loss function L' .

Fair Affirmative Action (2)

Formally, we can express the first step of this alternative approach as a restricted Earthmover problem defined as

$$\begin{aligned} d_{\text{EM+L}}(S, T) &\stackrel{\text{def}}{=} \min \mathbb{E}_{x \in S} \mathbb{E}_{y \sim \mu_x} d(x, y) & (15) \\ \text{subject to } & D(\mu_x, \mu_{x'}) \leq d(x, x') \quad \text{for all } x, x' \in S \\ & D_{\text{tv}}(\mu_S, U_T) \leq \epsilon \\ & \mu_x \in \Delta(T) \quad \text{for all } x \in S \end{aligned}$$

Here, U_T denotes the uniform distribution over T . Given $\{\mu_x\}_{x \in S}$ which minimizes (15) and $\{\nu_x\}_{x \in T}$ which minimizes the original fairness LP (2) restricted to T , we define the mapping $M: V \rightarrow \Delta(A)$ by putting

$$M(x) = \begin{cases} \nu_x & x \in T \\ \mathbb{E}_{y \sim \mu_x} \nu_y & x \in S \end{cases} . & (16)$$

Fair Affirmative Action (3)

Proposition 4.1. *The mapping M defined in (16) satisfies*

- 1. statistical parity between S and T up to bias ϵ ,*
- 2. the Lipschitz condition for every pair $(x, y) \in (S \times S) \cup (T \times T)$.*

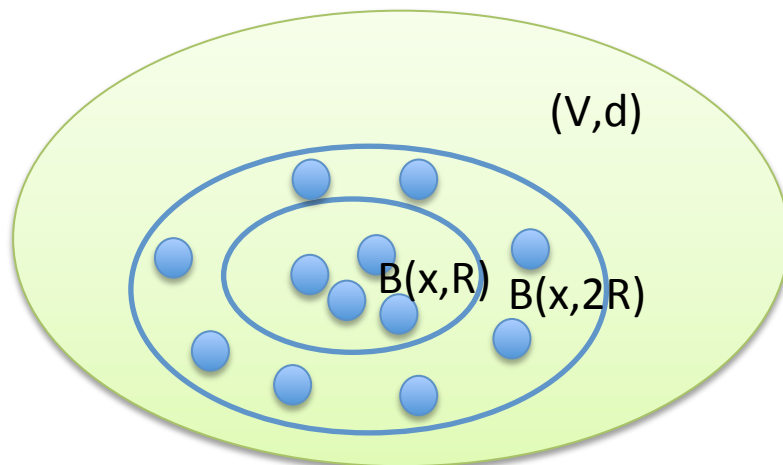
Proposition 4.2. *Suppose $D = D_{\text{tv}}$ in (15). Then, the resulting mapping M satisfies*

$$\mathbb{E} \max_{x \in S, y \in T} \left[D_{\text{tv}}(M(x), M(y)) - d(x, y) \right] \leq d_{\text{EM+L}}(S, T).$$

Can we import techniques from Differential Privacy?

Theorem: Fairness mechanism with “high utility” in metric spaces (V, d) of bounded doubling dimension

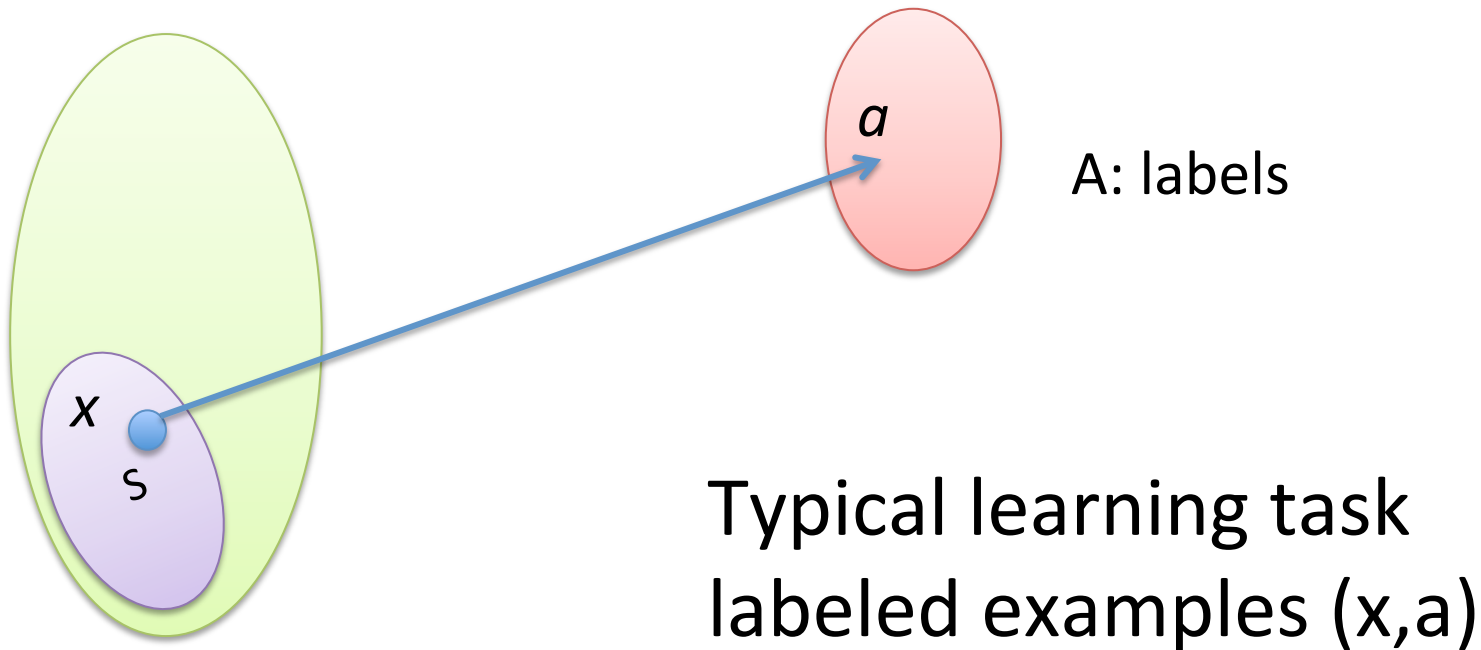
Based on exponential mechanism [MT'07]



$$|B(x, R)| \leq O(|B(x, 2R)|)$$

Some recent work

- Zemel-Wu-Swersky-Pitassi-Dwork
“Learning Fair Representations” (ICML 2013)



V : Individuals

S : protected set

Web Fairness Measurement

How do we measure the **“fairness of the web”**?

- Need to model/understand user browsing behavior
- Evaluate how web sites respond to different behavior/attributes
- Cope with noisy measurements
- Exciting progress by Datta, Datta, Tschantz