

Course Description

Lecture 5: Pipelines and Hazards

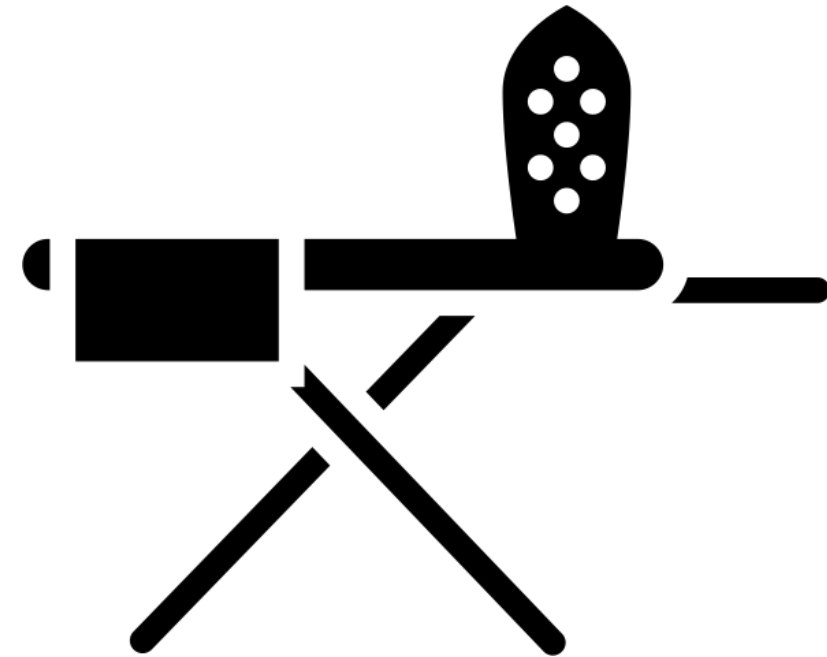
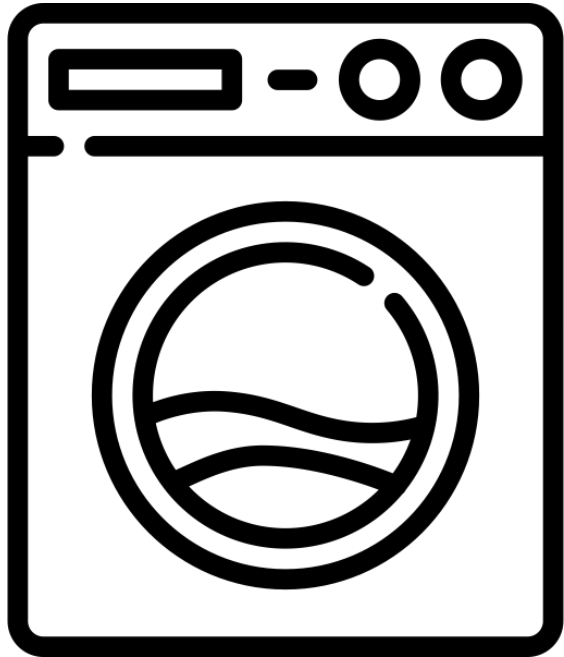
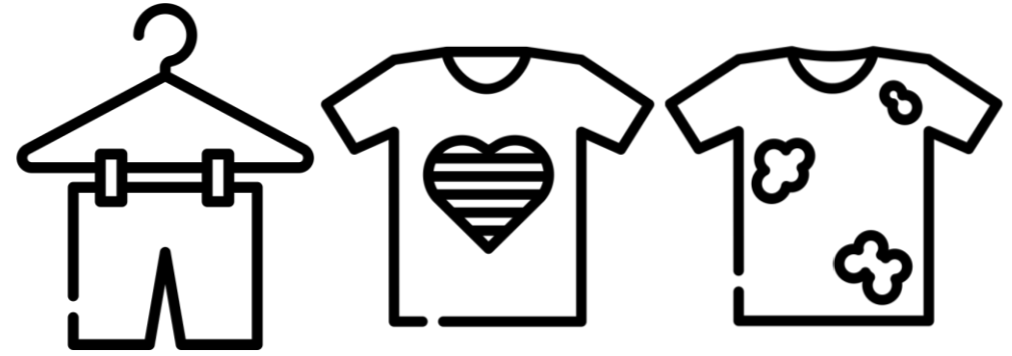
This course covers the design and implementation of computer systems from the perspective of the hardware software interface. The purpose of this course is for students to understand the relationship between the operating system, software, and computer architecture. Students that complete the course will have learned operating system fundamentals, computer architecture fundamentals, compilation to hardware abstractions, and how software actually executes from the perspective of the hardware software/boundary. The course will focus especially on understanding the relationships between software and hardware, and how those relationships influence the design of a computer system's software and hardware. The course will convey these topics through a series of practical, implementation-oriented lab assignments.

Credit: Brandon Lucia

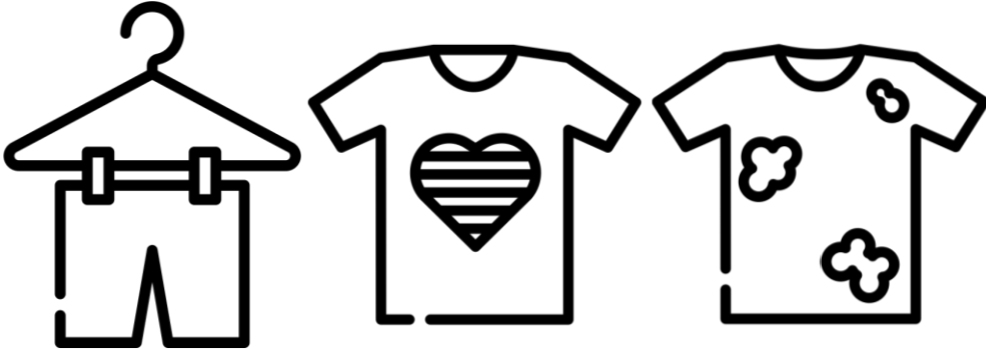
Pipelined Microarchitectural Implementation

- Pipelining for Instruction-Level Parallelism (ILP)
- Pipelined microarchitecture design sketch
- Control hazards
- Branch prediction for dealing with control hazards

Pipelining in the abstract: A laundry efficiency problem



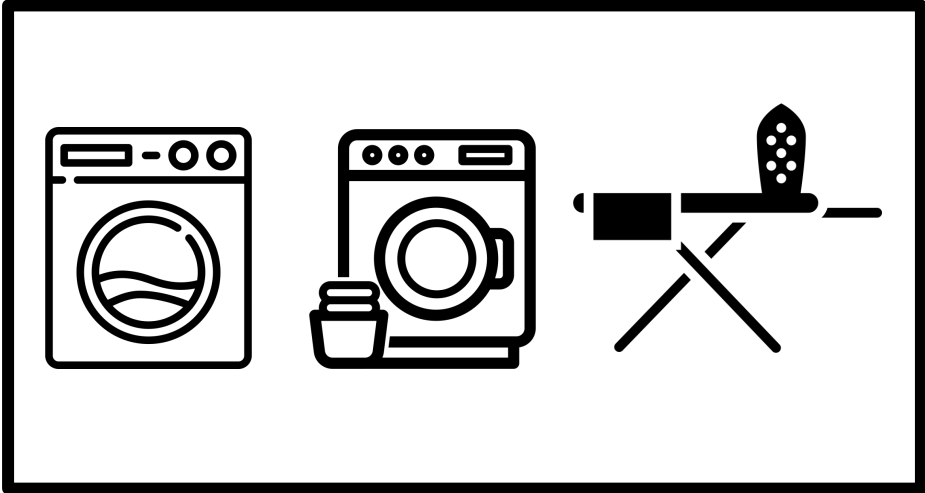
Pipelining in the abstract: A laundry efficiency problem



To be washed



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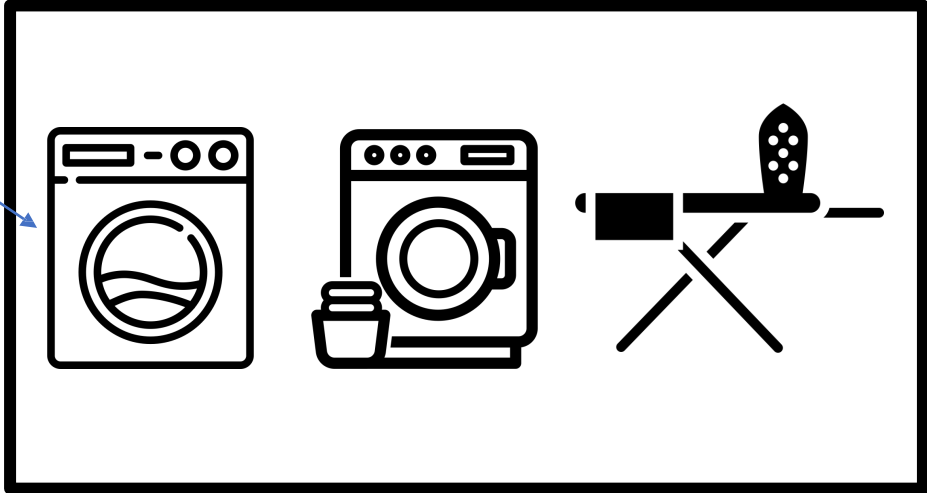


Private Laundry Room Model: only one person at a time allowed in laundry room

Pipelining in the abstract: A laundry efficiency problem



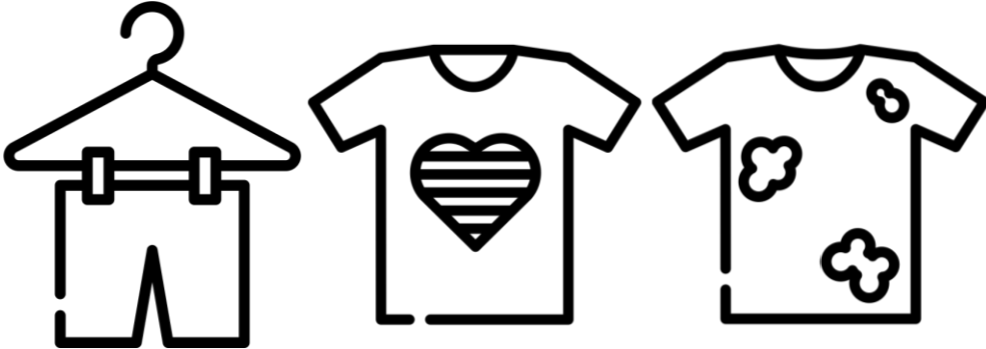
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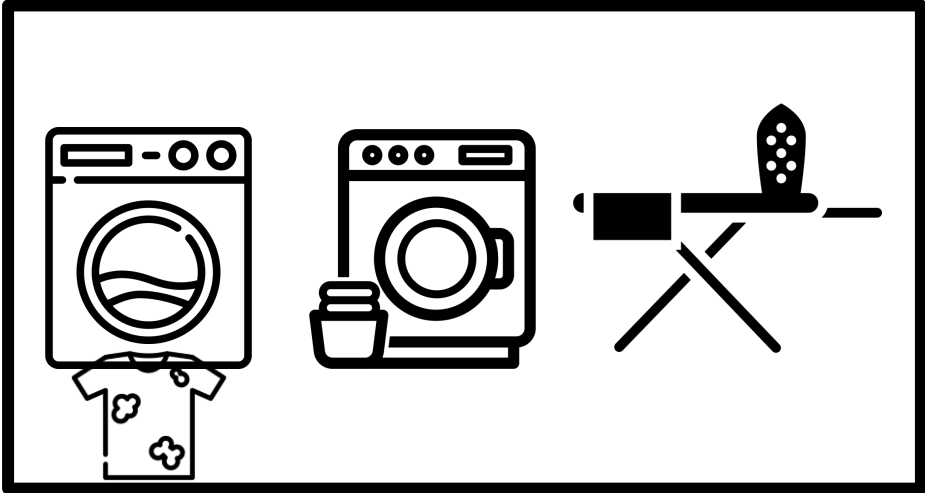
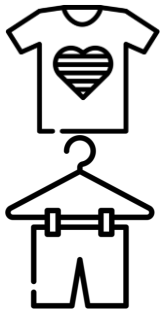
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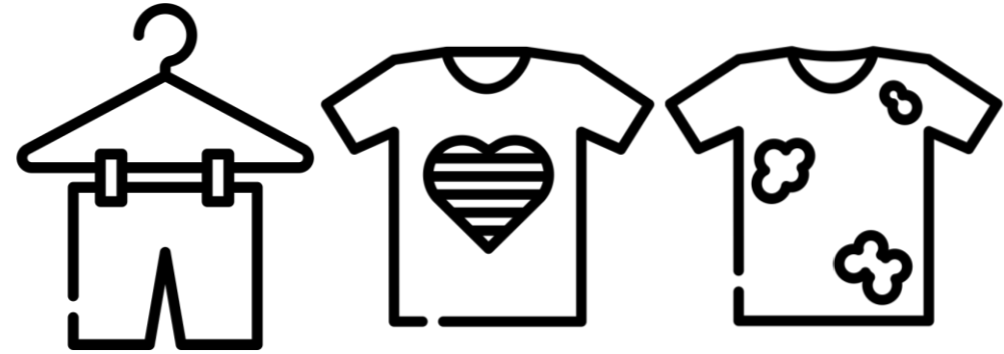
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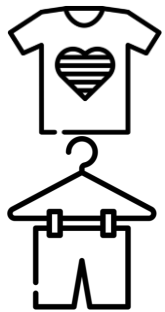


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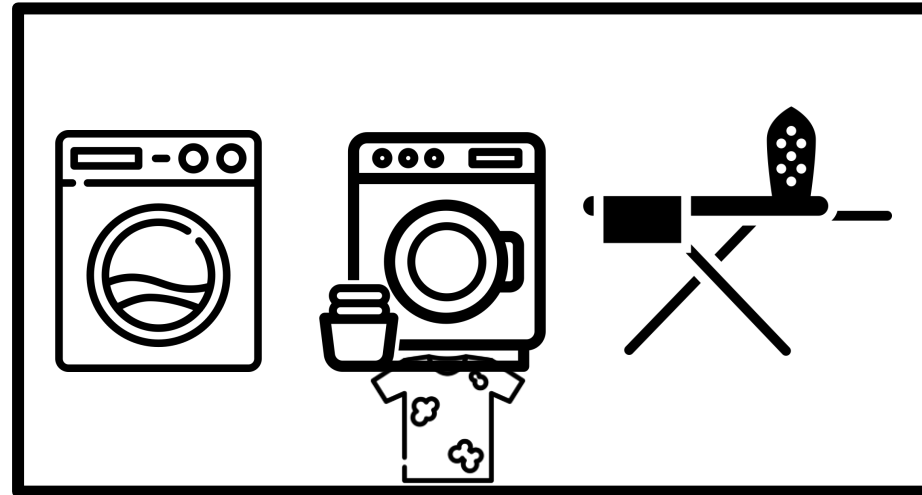
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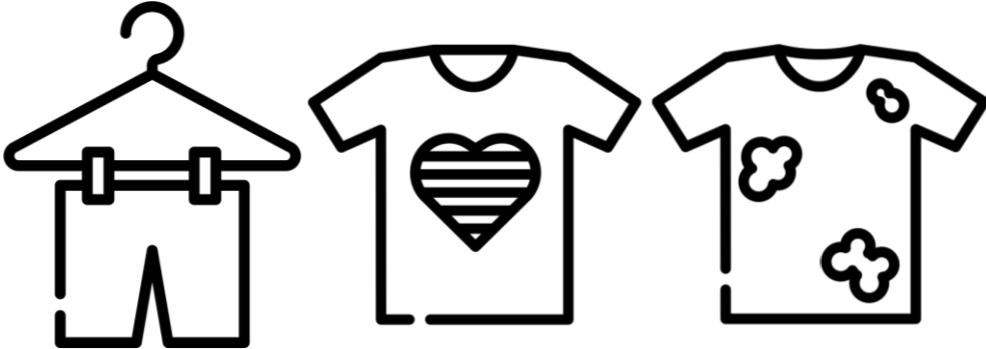


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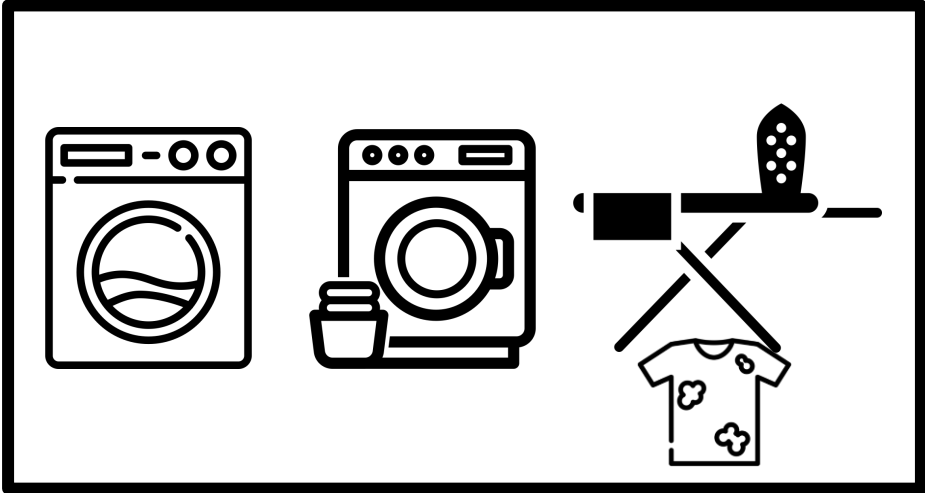
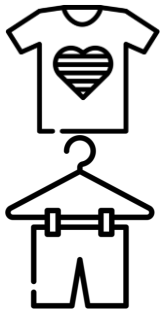
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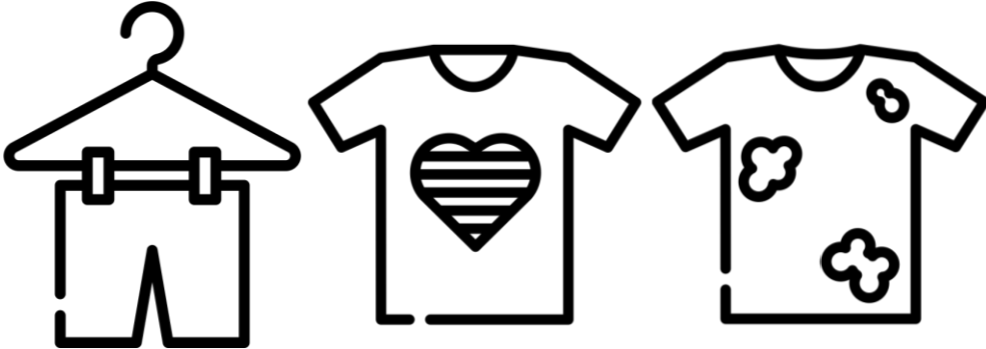
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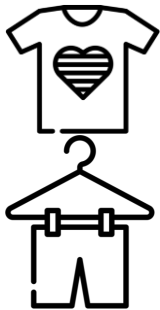


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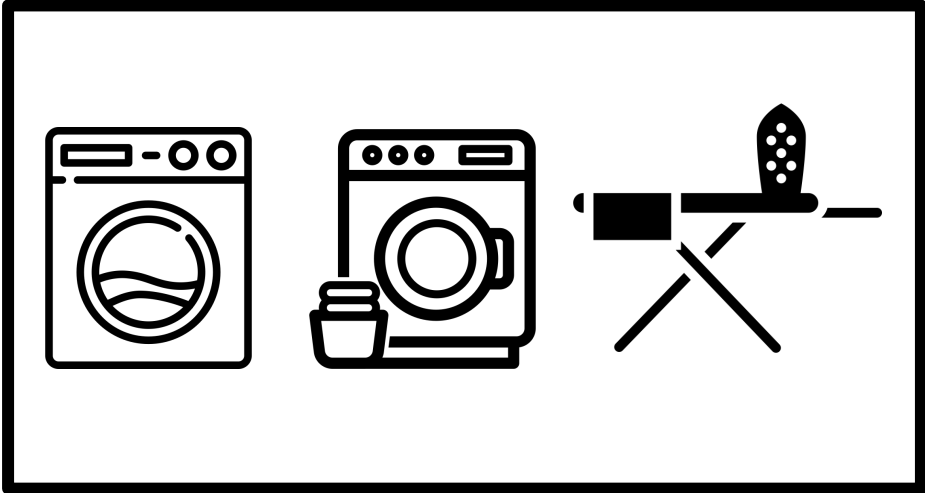
Pipelining in the abstract: A laundry efficiency problem



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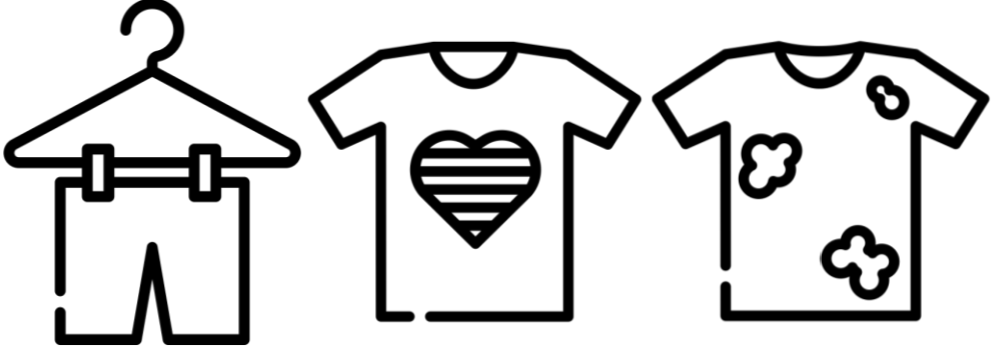
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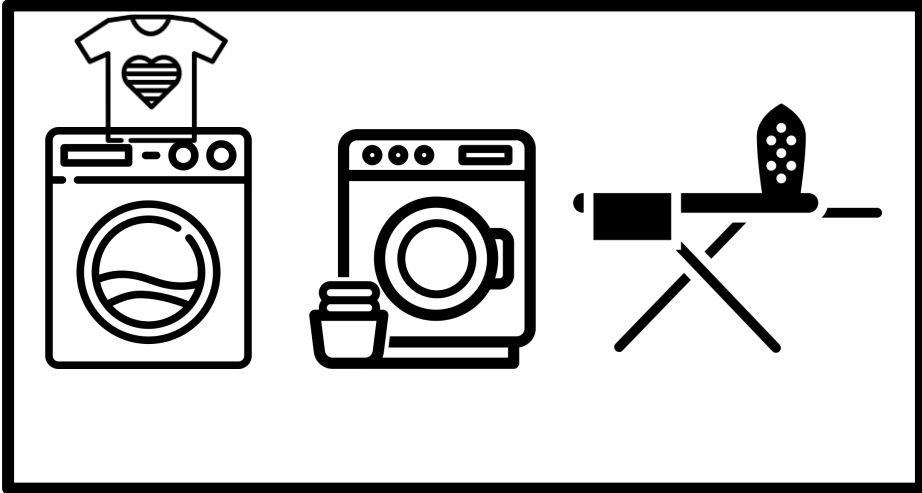
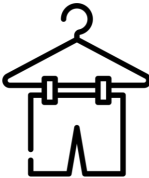


Pipelining in the abstract: A laundry efficiency problem



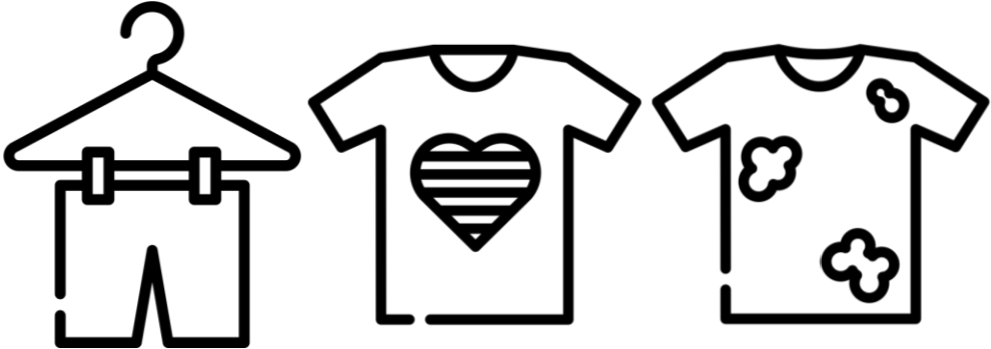
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Done being washed



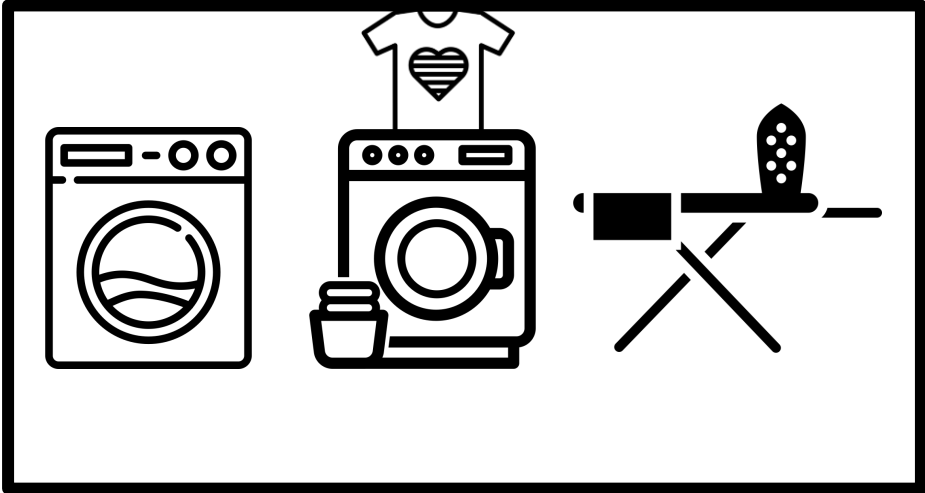
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Pipelining in the abstract: A laundry efficiency problem



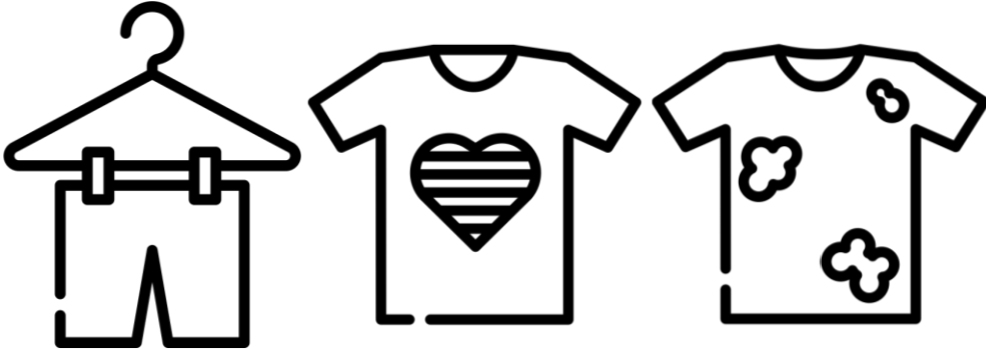
To be washed

Done being washed



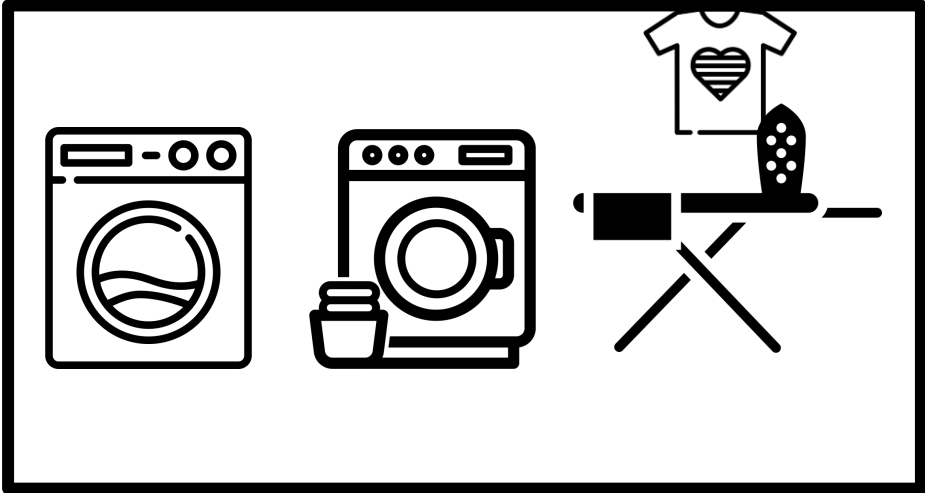
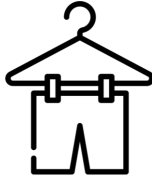
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Pipelining in the abstract: A laundry efficiency problem



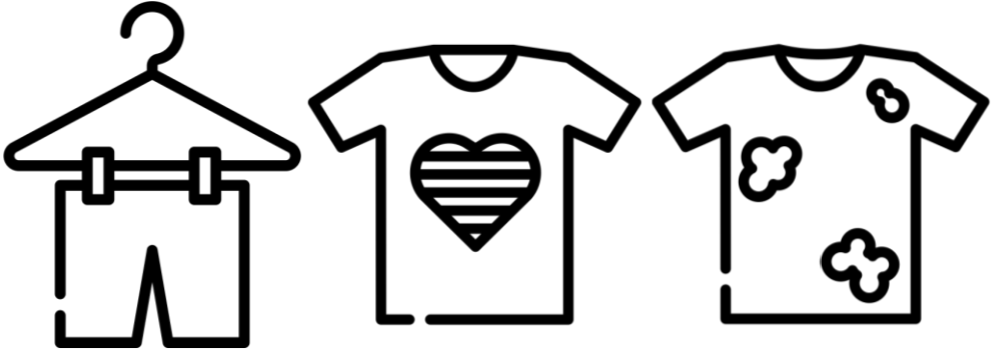
To be washed

Done being washed



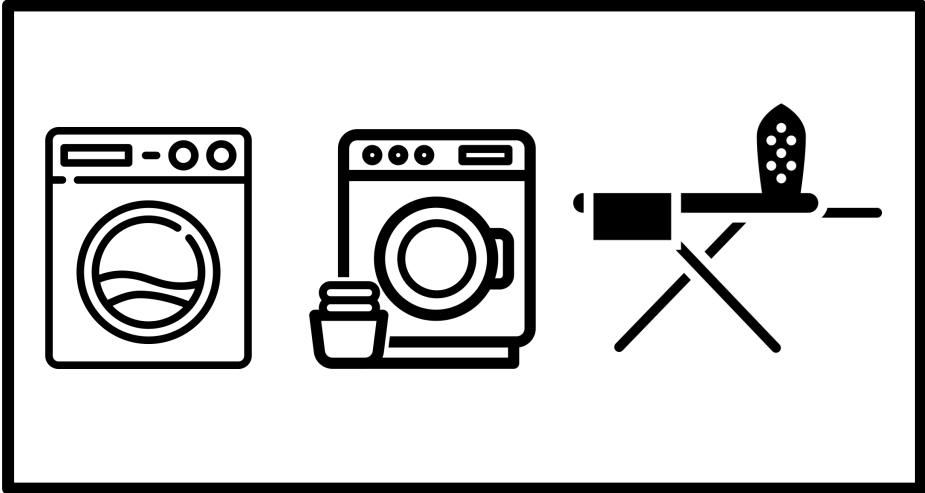
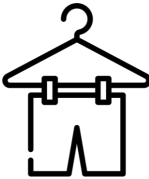
Time = 7

Pipelining in the abstract: A laundry efficiency problem



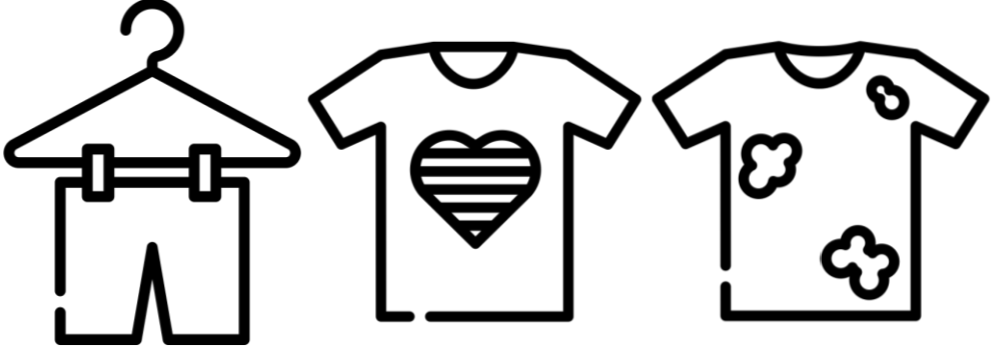
To be washed

Done being washed

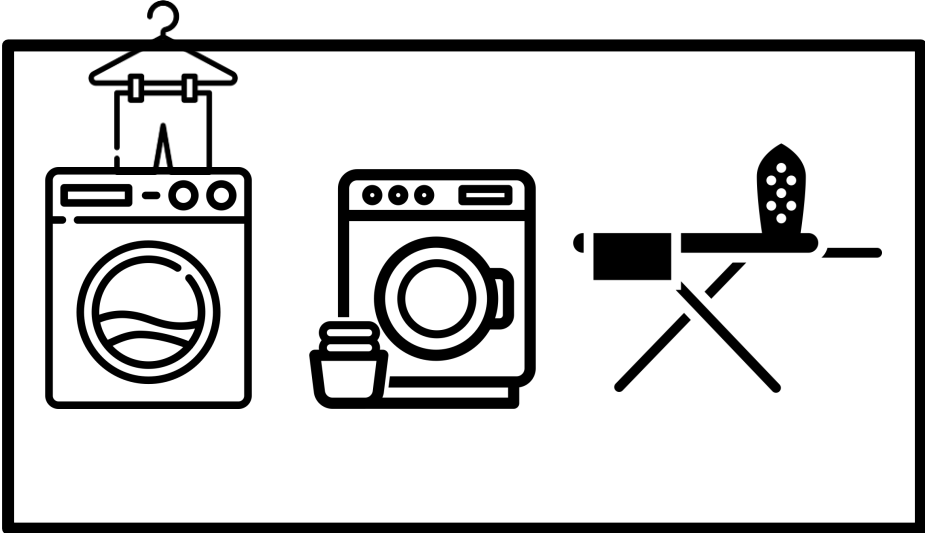


Time = 8

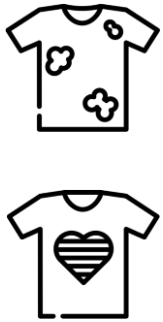
Pipelining in the abstract: A laundry efficiency problem



To be washed

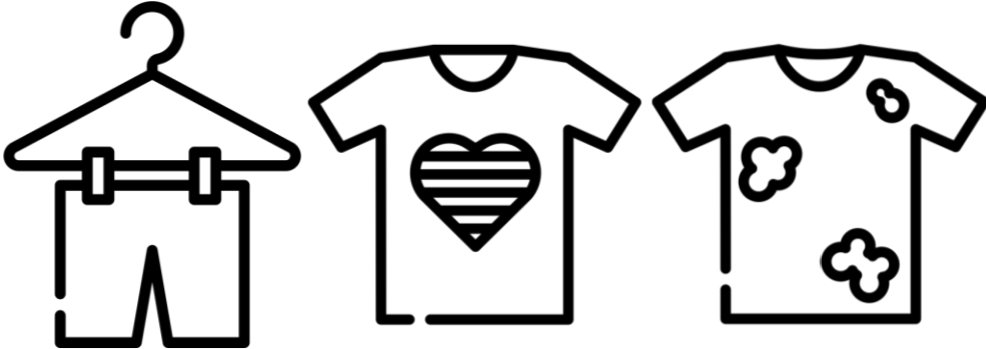


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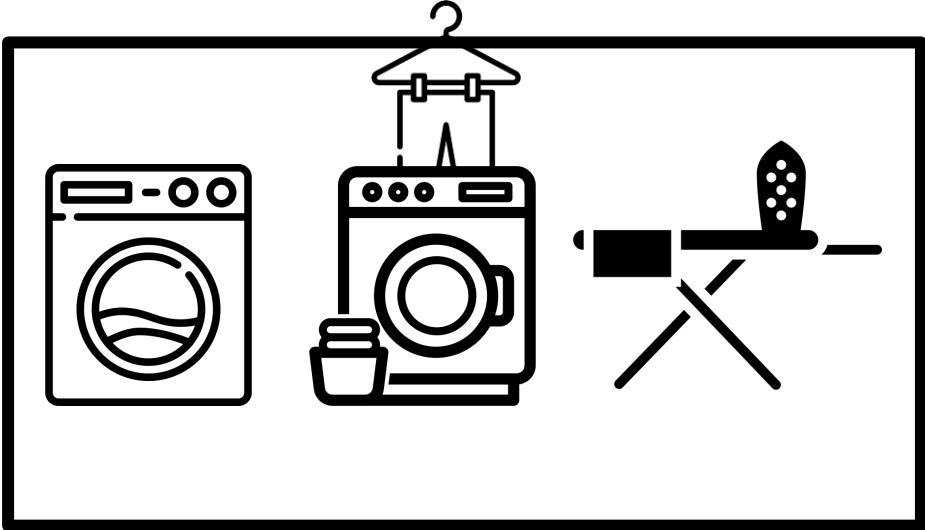


Time = 9

Pipelining in the abstract: A laundry efficiency problem



To be washed

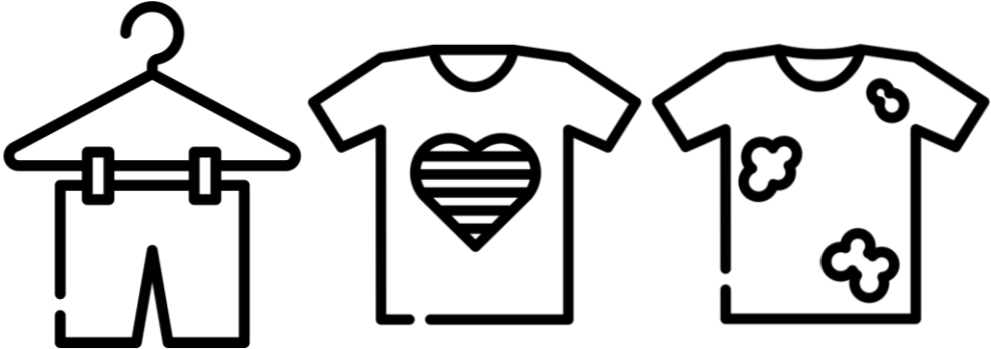


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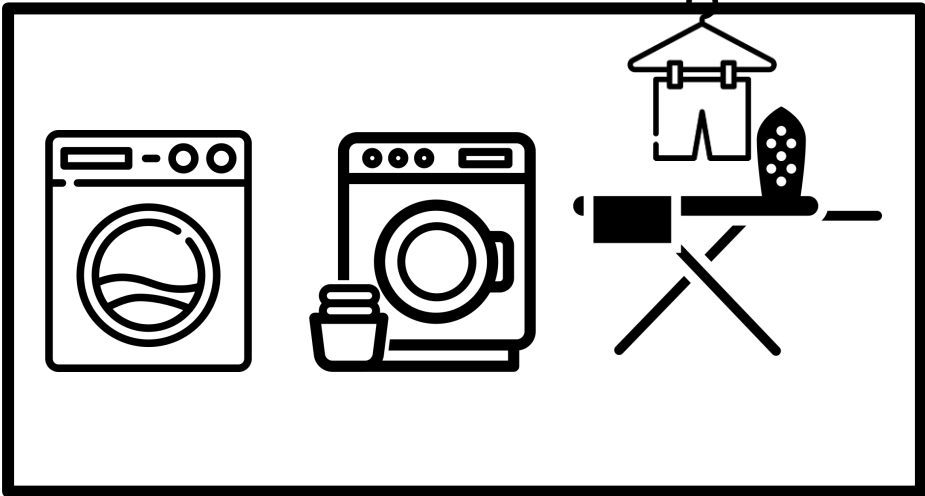
Time = 10

Pipelining in the abstract: A laundry efficiency problem



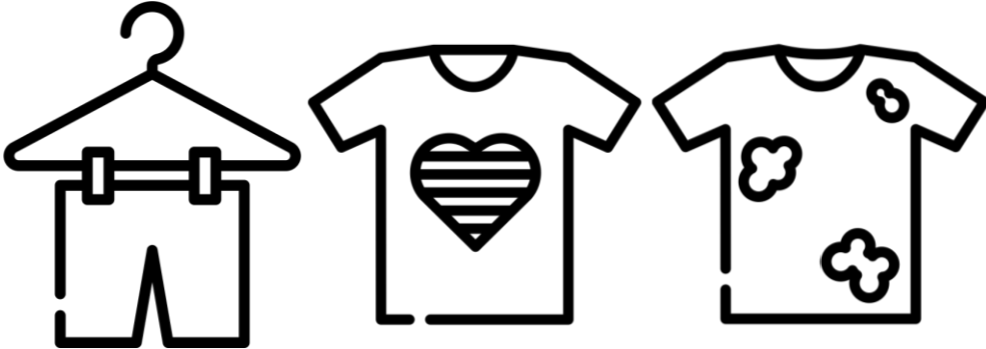
To be washed

Done being washed

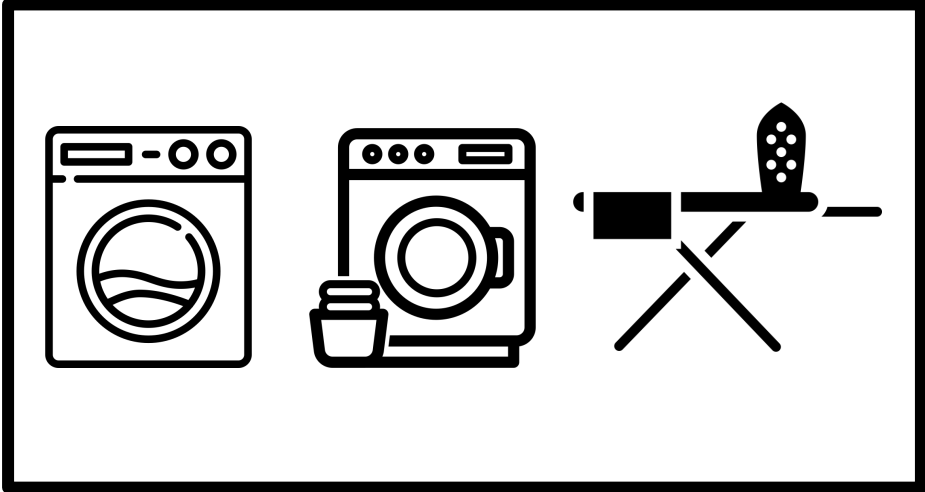


Time = 11

Pipelining in the abstract: A laundry efficiency problem



To be washed

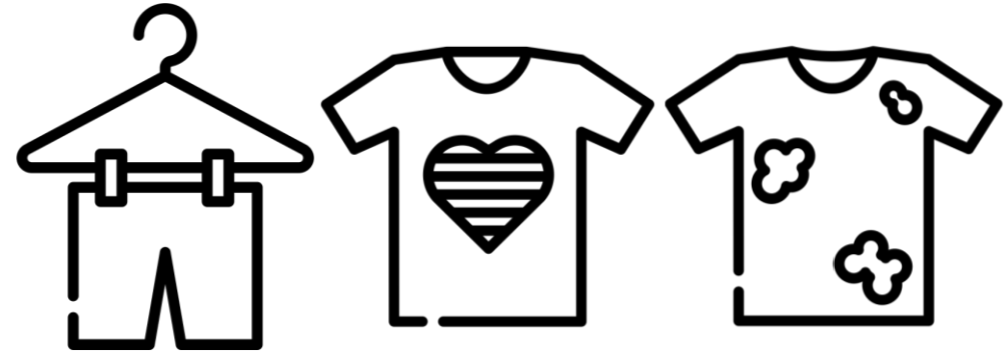


Done being washed

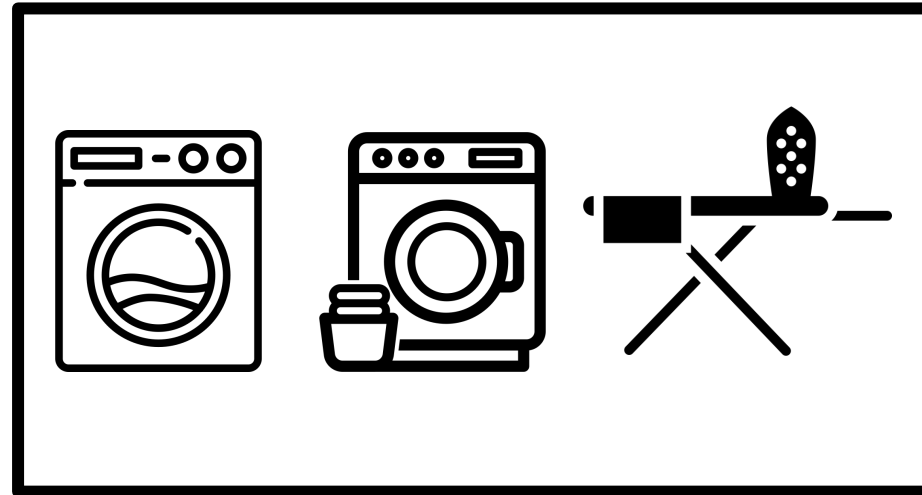


Time = 12

Pipelining in the abstract: A laundry efficiency problem



To be washed

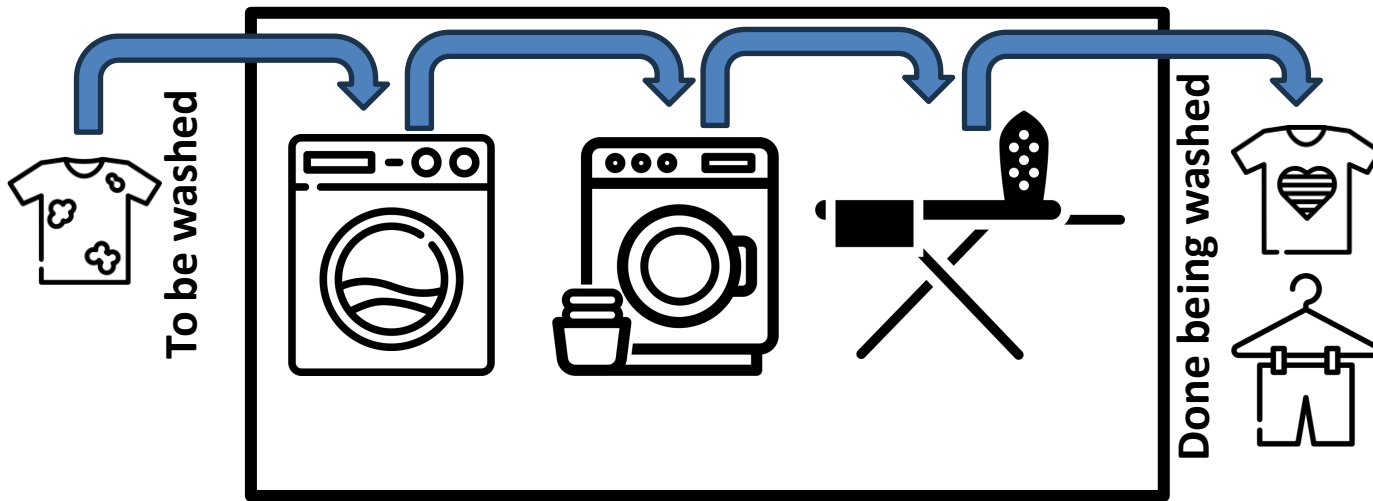


Done being washed



Analysis: With 3 resources ( ,  , ) and 3 units of work ( ,  , ) our laundry took 12 units of time

Pipelining in the abstract: A laundry efficiency problem



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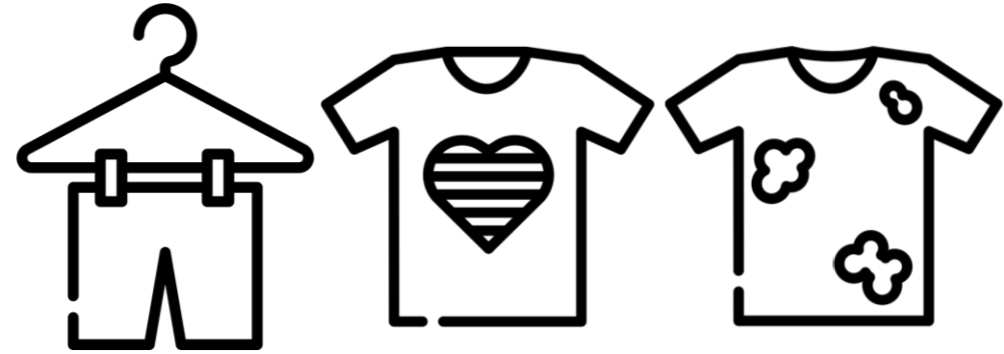
12 units of time?

Why 12 units of time vs 9 units of time overall?

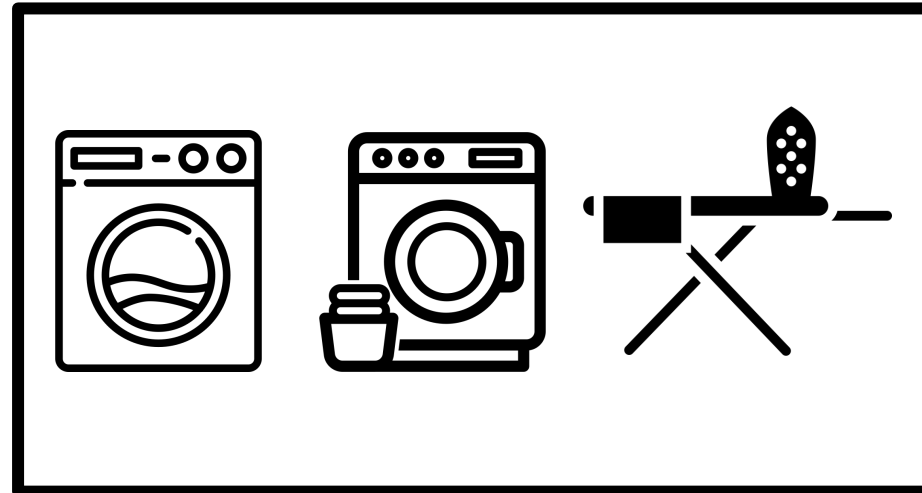
Why 4 units of time per load vs 3 units of time?

- Processors and their workings are triggered devices.
- It takes 4 triggers for the dirty laundry pile to be washed, dried, folded, and available.

Pipelining in the abstract: A laundry efficiency problem



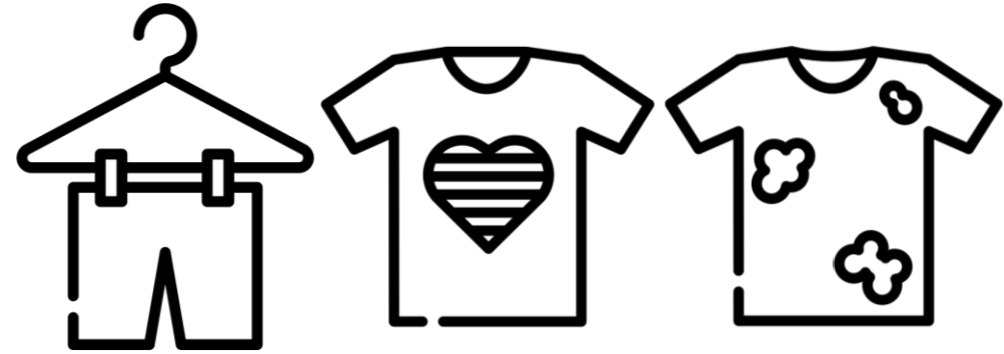
To be washed



Done being washed

Let's redesign our laundry room to make it more efficient

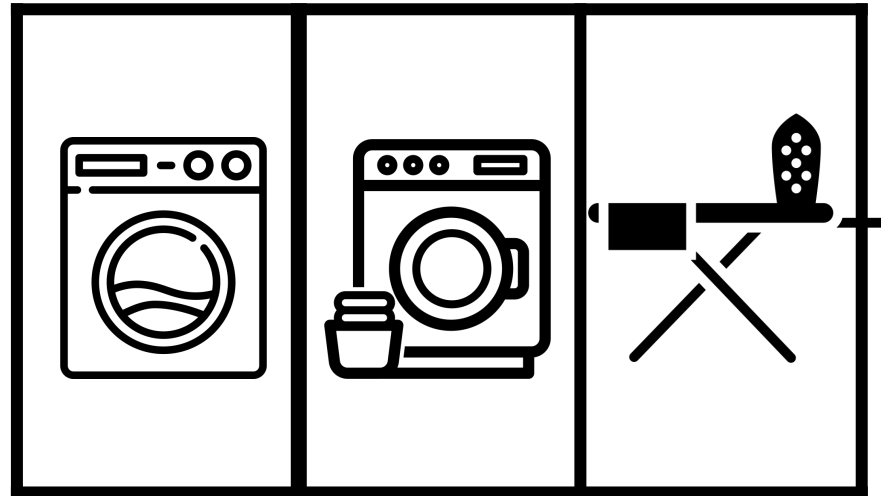
Pipelining in the abstract: A laundry efficiency problem



To be washed

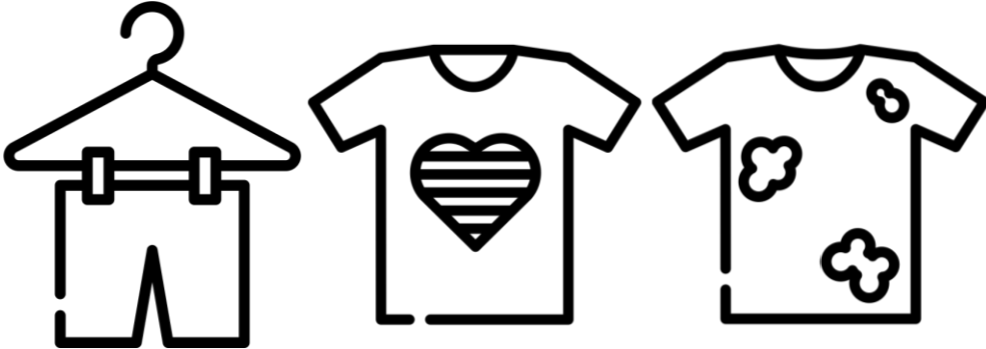


Done being washed



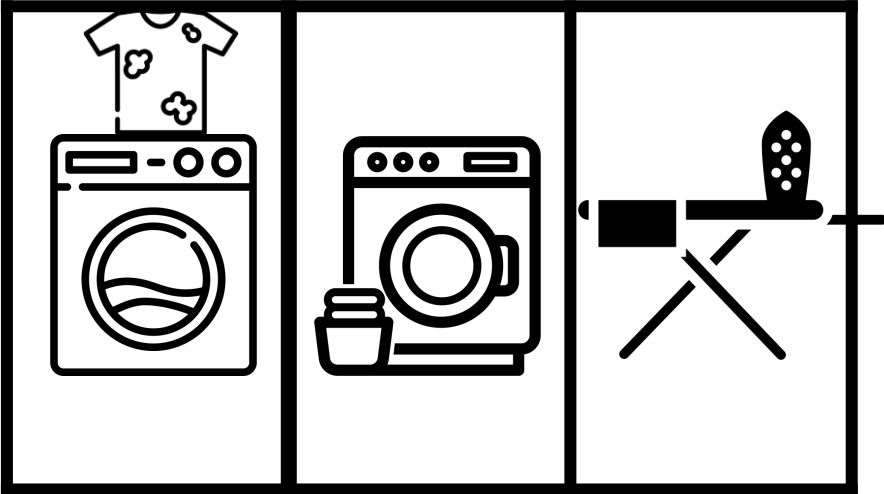
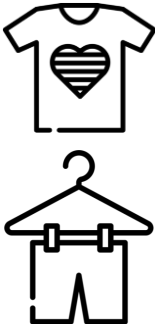
Shared Laundry Room: single laundry task uses single machine at a time, not entire room. Multiple roommates allowed in at once.

Pipelining in the abstract: A laundry efficiency problem



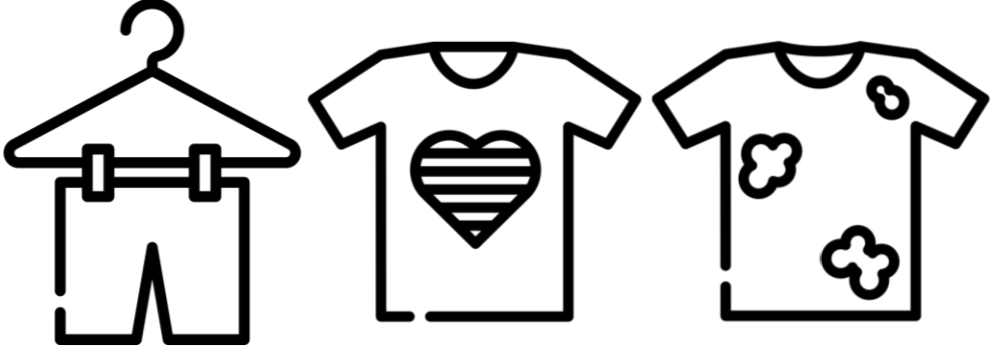
To be washed

Done being washed



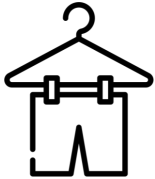
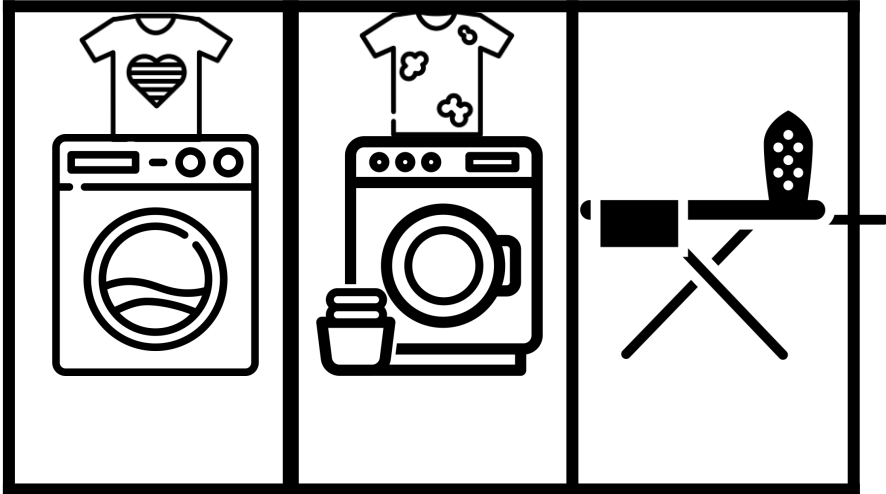
Time = 1

Pipelining in the abstract: A laundry efficiency problem



To be washed

Done being washed



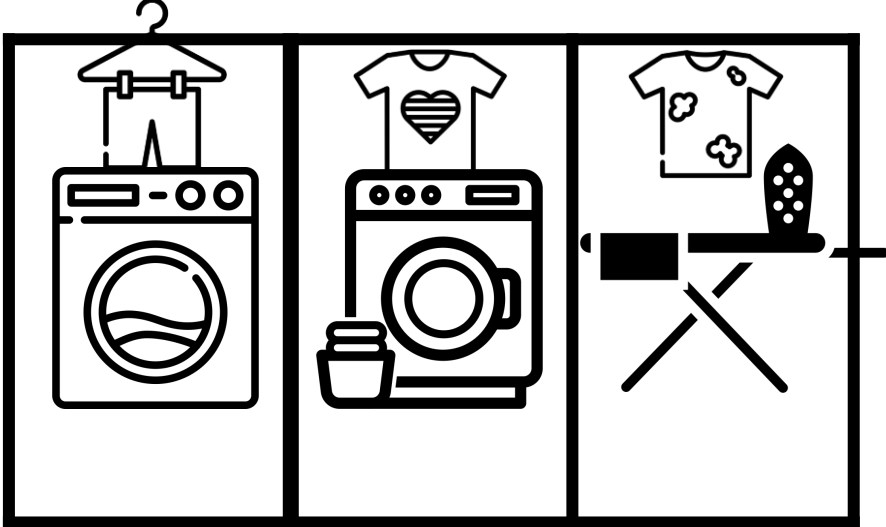
Time = 2

Pipelining in the abstract: A laundry efficiency problem



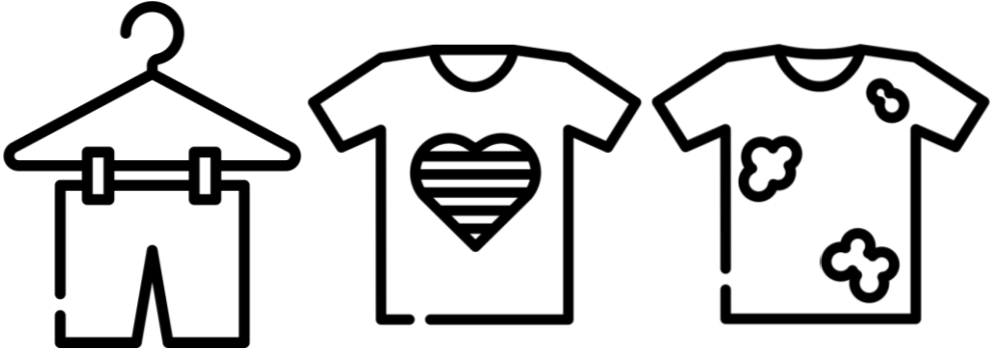
To be washed

Done being washed

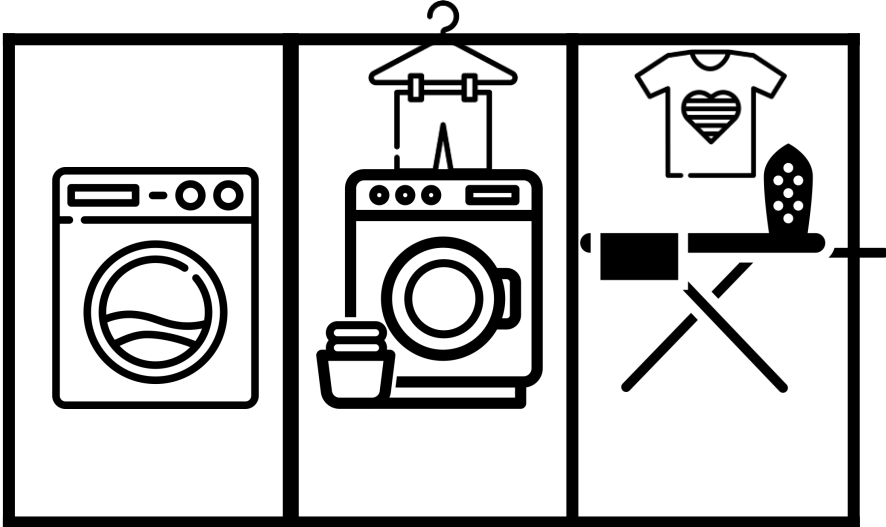


Time = 3

Pipelining in the abstract: A laundry efficiency problem



To be washed

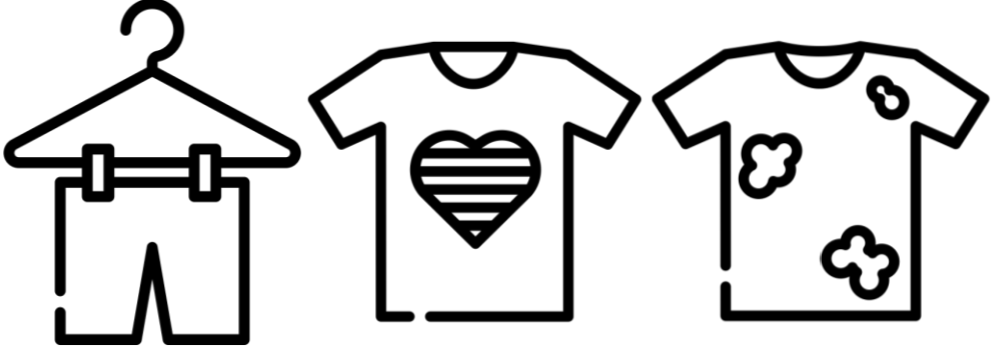


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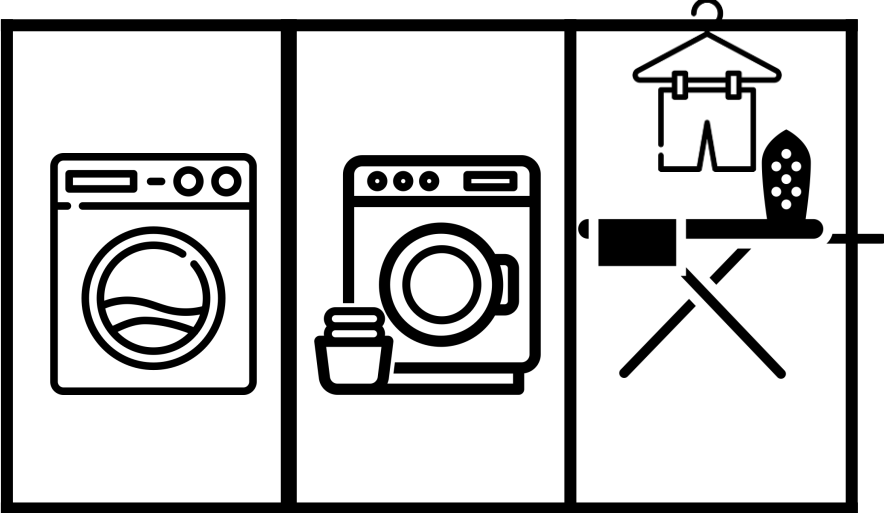


Time = 4

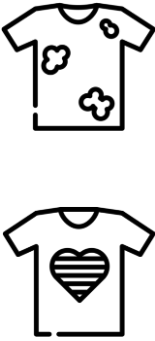
Pipelining in the abstract: A laundry efficiency problem



To be washed

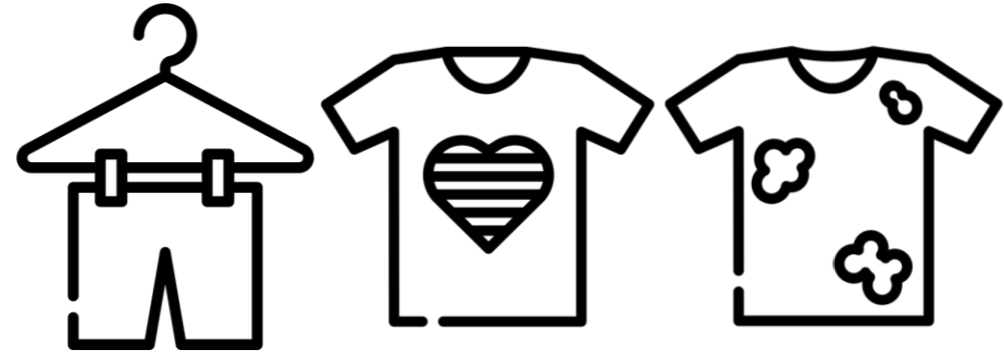


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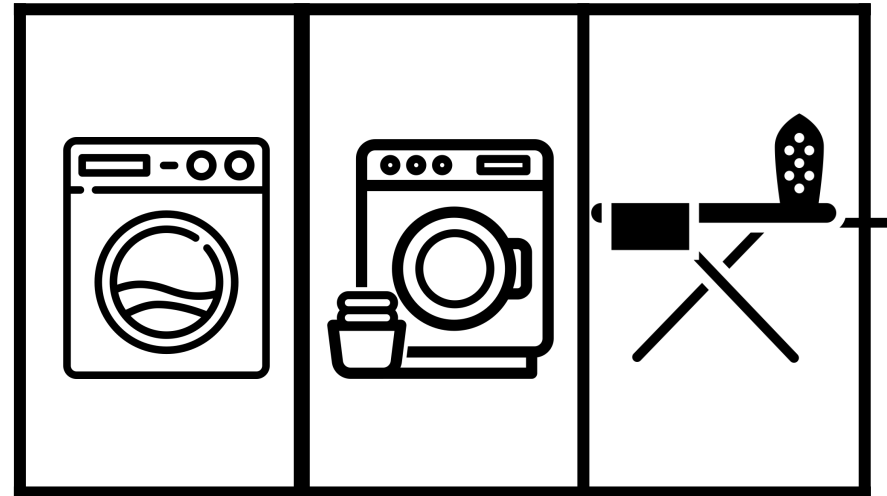


Time = 5

Pipelining in the abstract: A laundry efficiency problem



To be washed

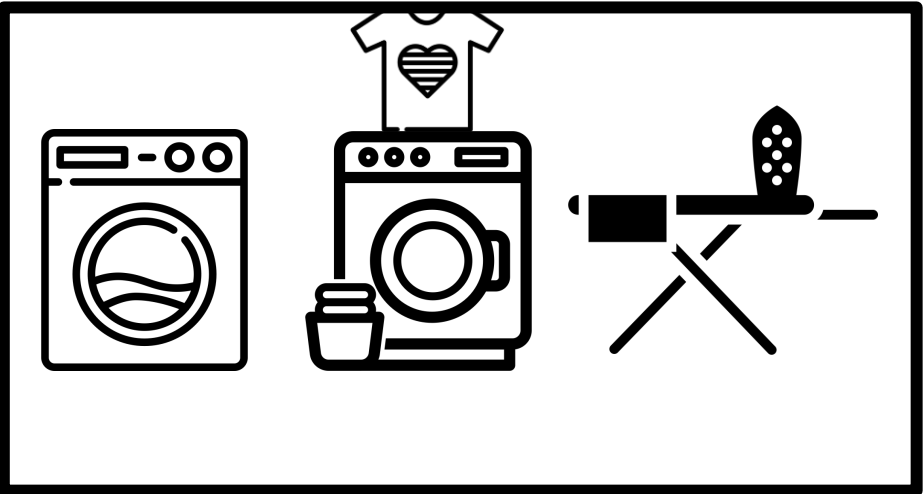
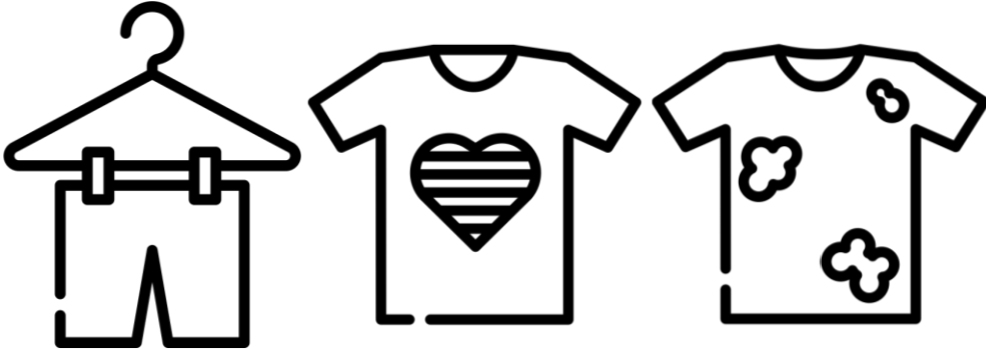


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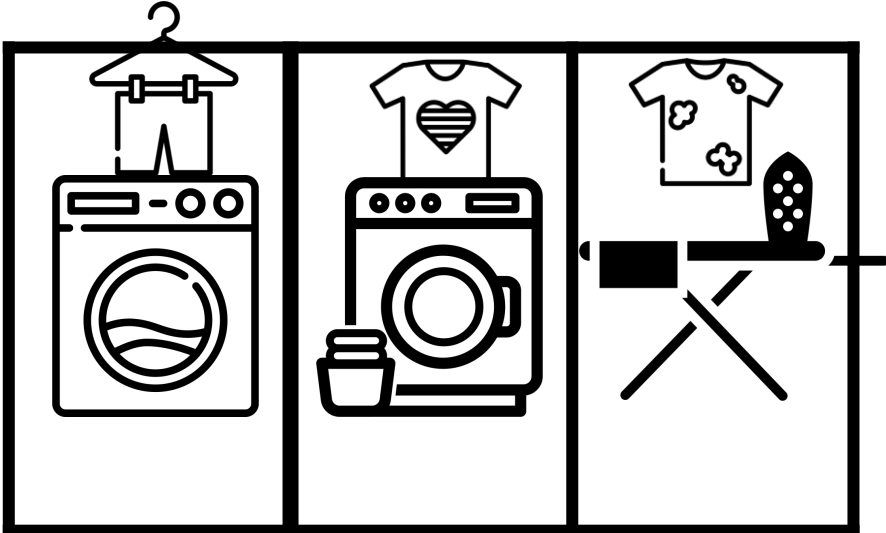


Analysis: With 3 resources ( ,  , ) and 3 units of work ( ,  , ) our laundry took 6 units of time

Pipelining in the abstract: A laundry efficiency problem

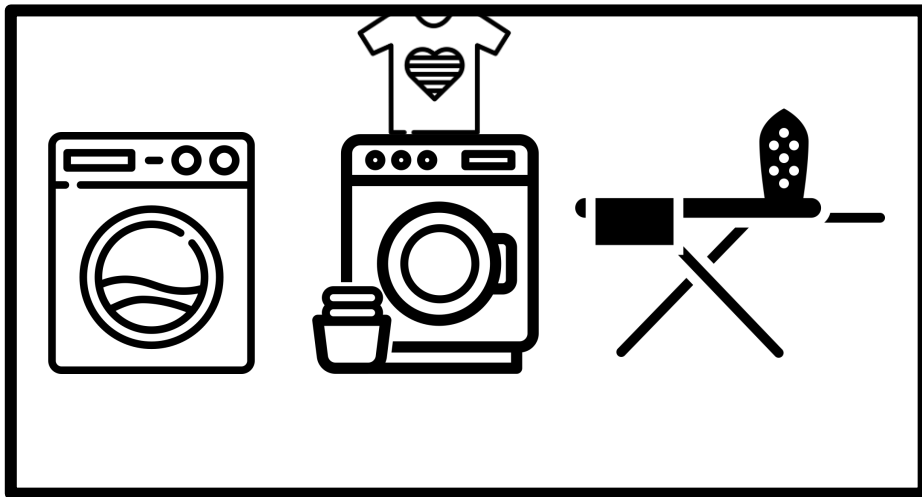
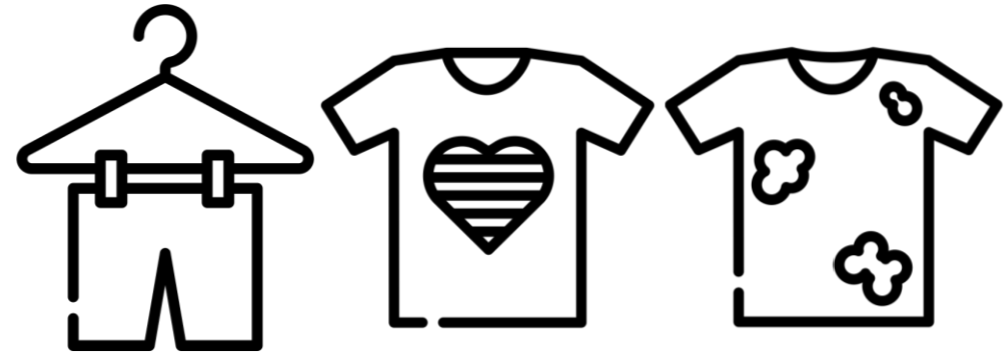


Vs.

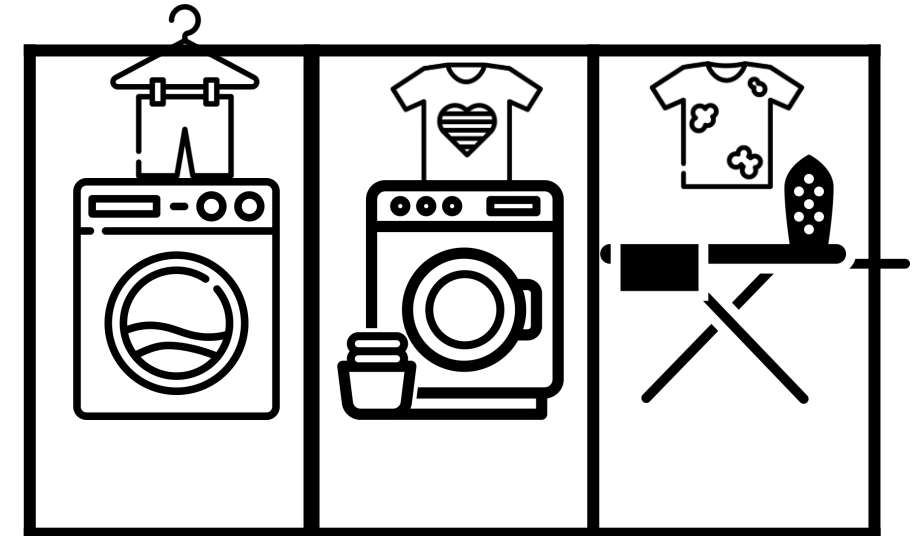


General observations about private laundry room model vs. shared laundry room model?

Pipelining in the abstract: A laundry efficiency problem



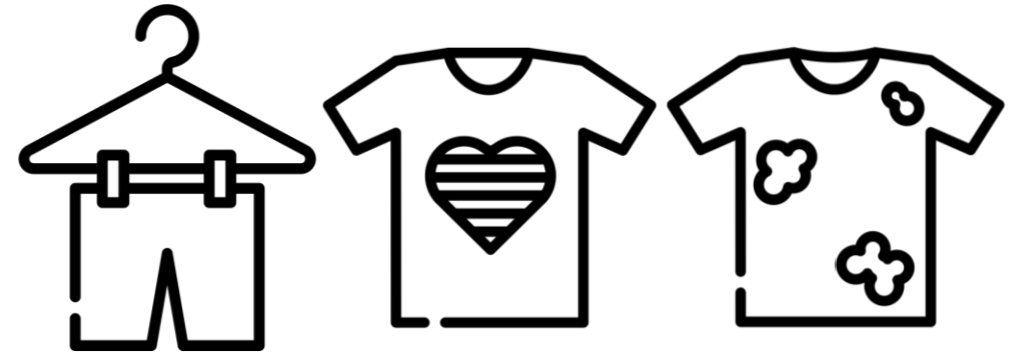
Vs.



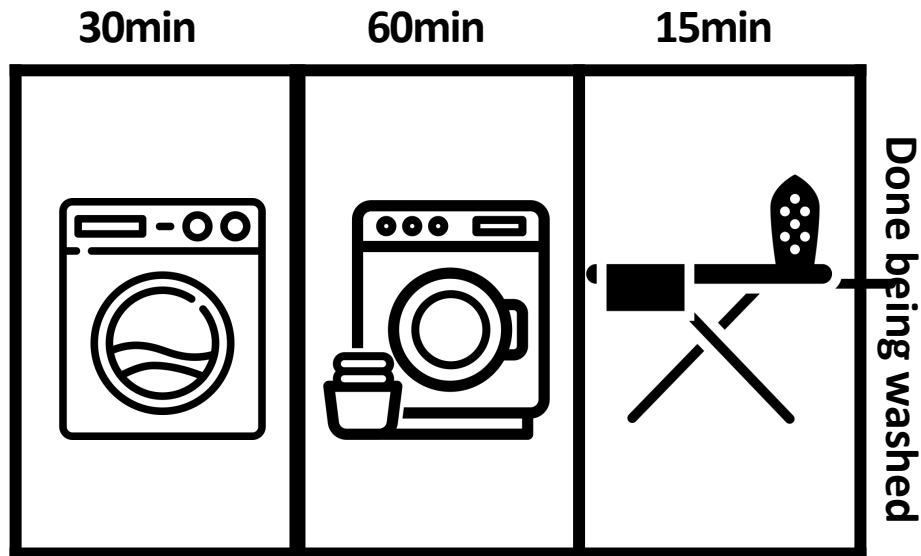
General observations about private laundry room model vs. shared laundry room model?

- Using machines *in parallel* in the shared laundry model
- At time step 3 (“steady state”) all machines are active
- Private model: always leaving 2/3 of laundry machines idle, despite laundry yet to wash!

Pipelining in the abstract: A laundry efficiency problem



To be washed

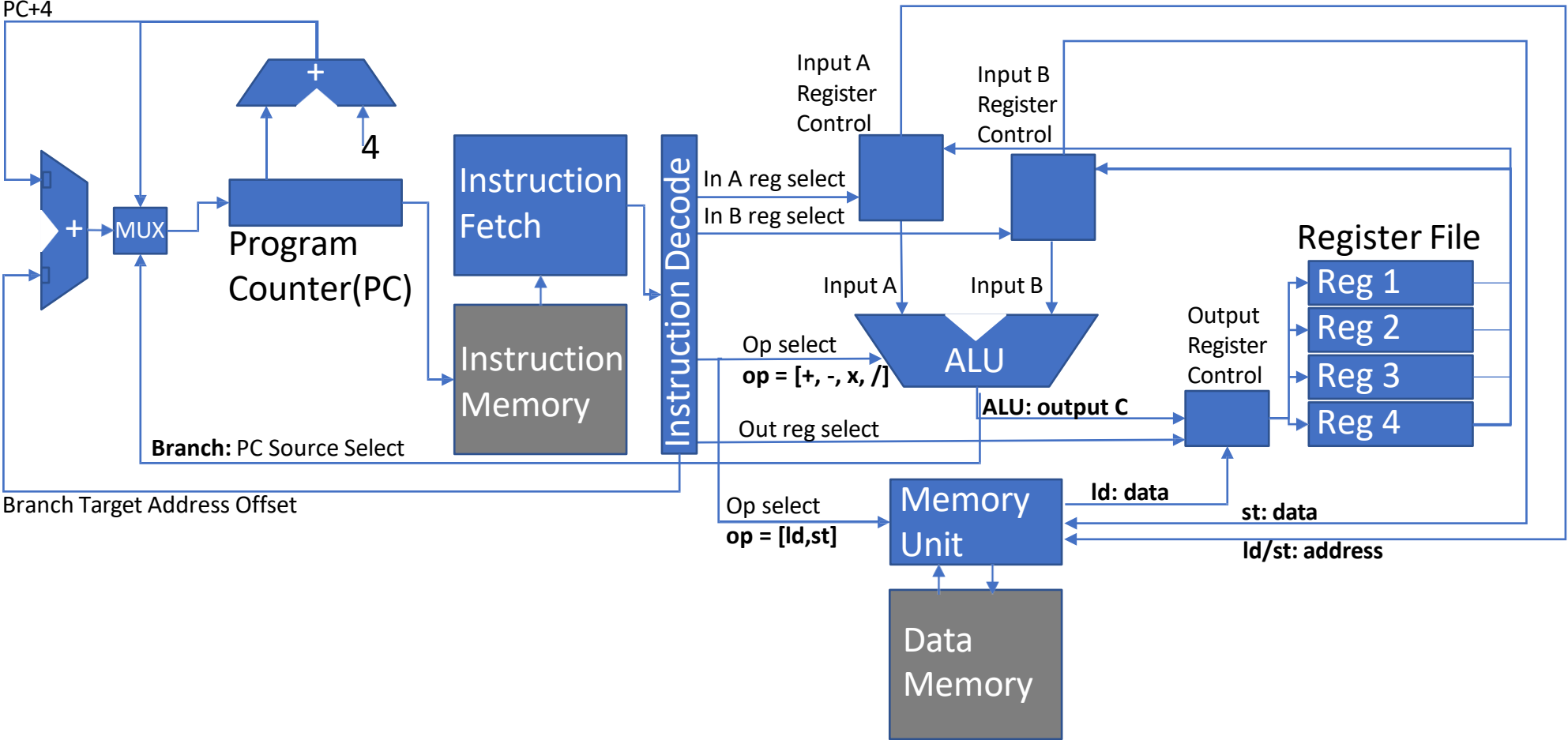


Shared Laundry Room: single laundry task uses single machine at a time, not entire room. Multiple roommates allowed in at once.

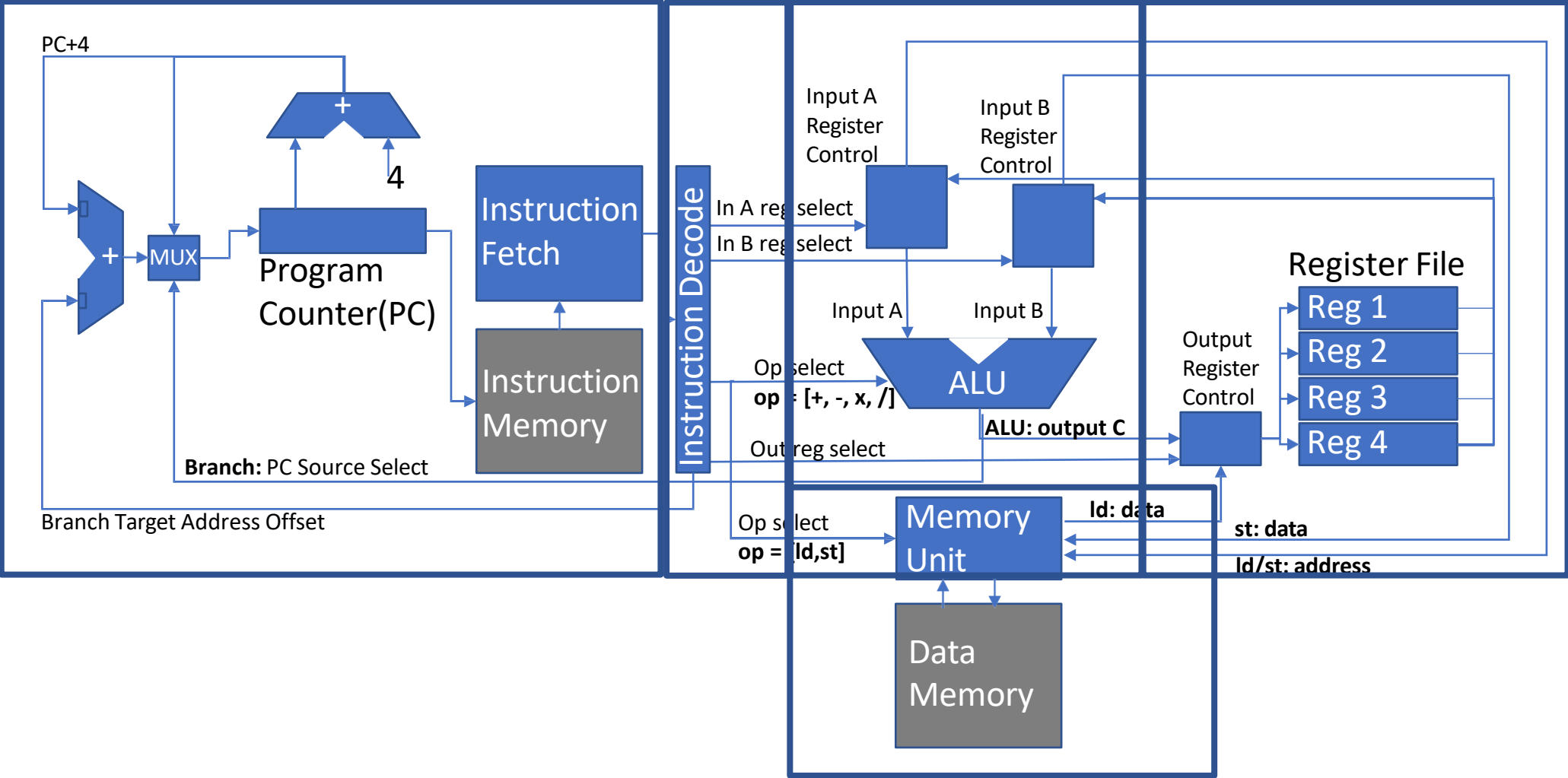
Improving Pipeline Performance

- If you could make washing take only 15 minutes what would be the impact upon throughput?
- What if you could make ironing take only 10 minutes?
- What if you could make drying take 45 minutes? Why is that different?
- *Hint: What (stage) limits the throughput? Why?*

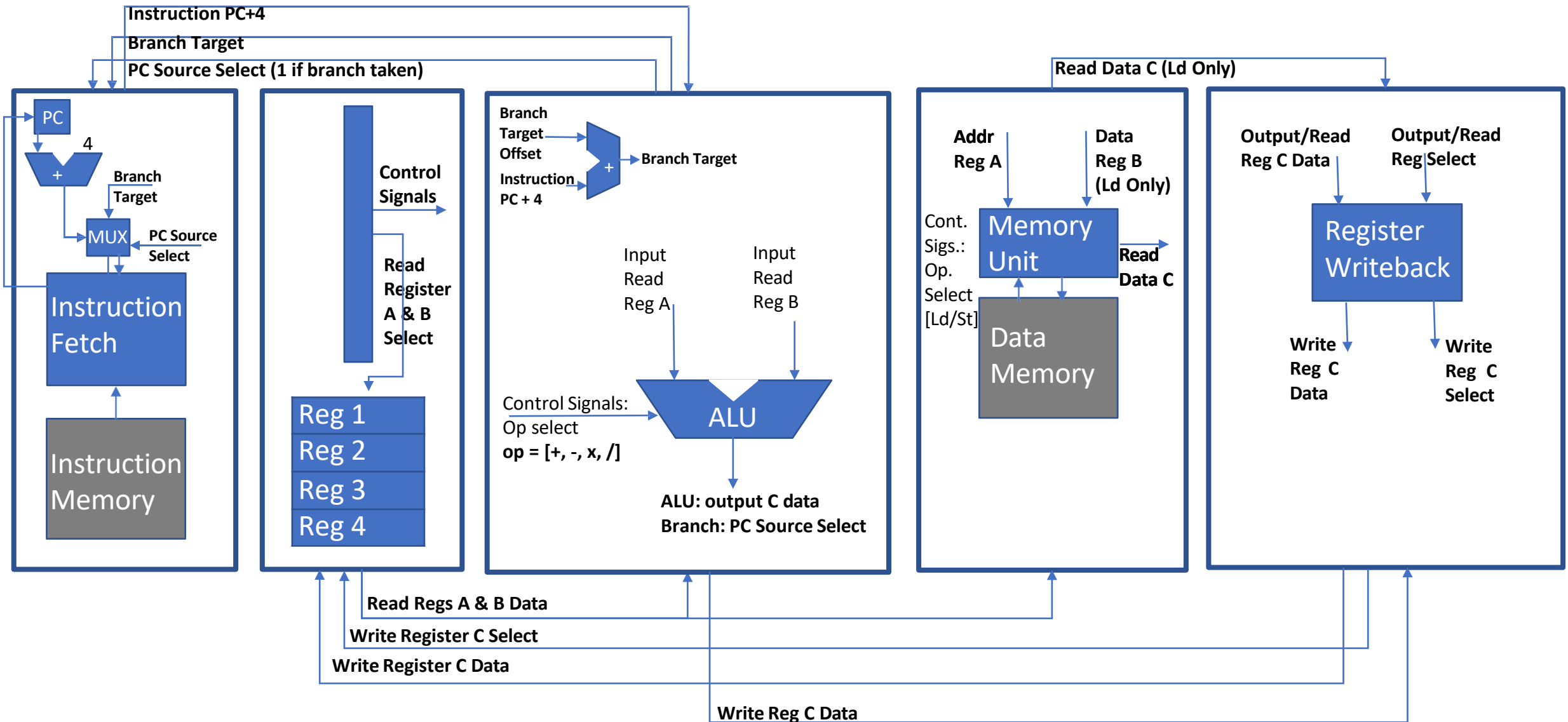
Let's do some grouping together of functionality



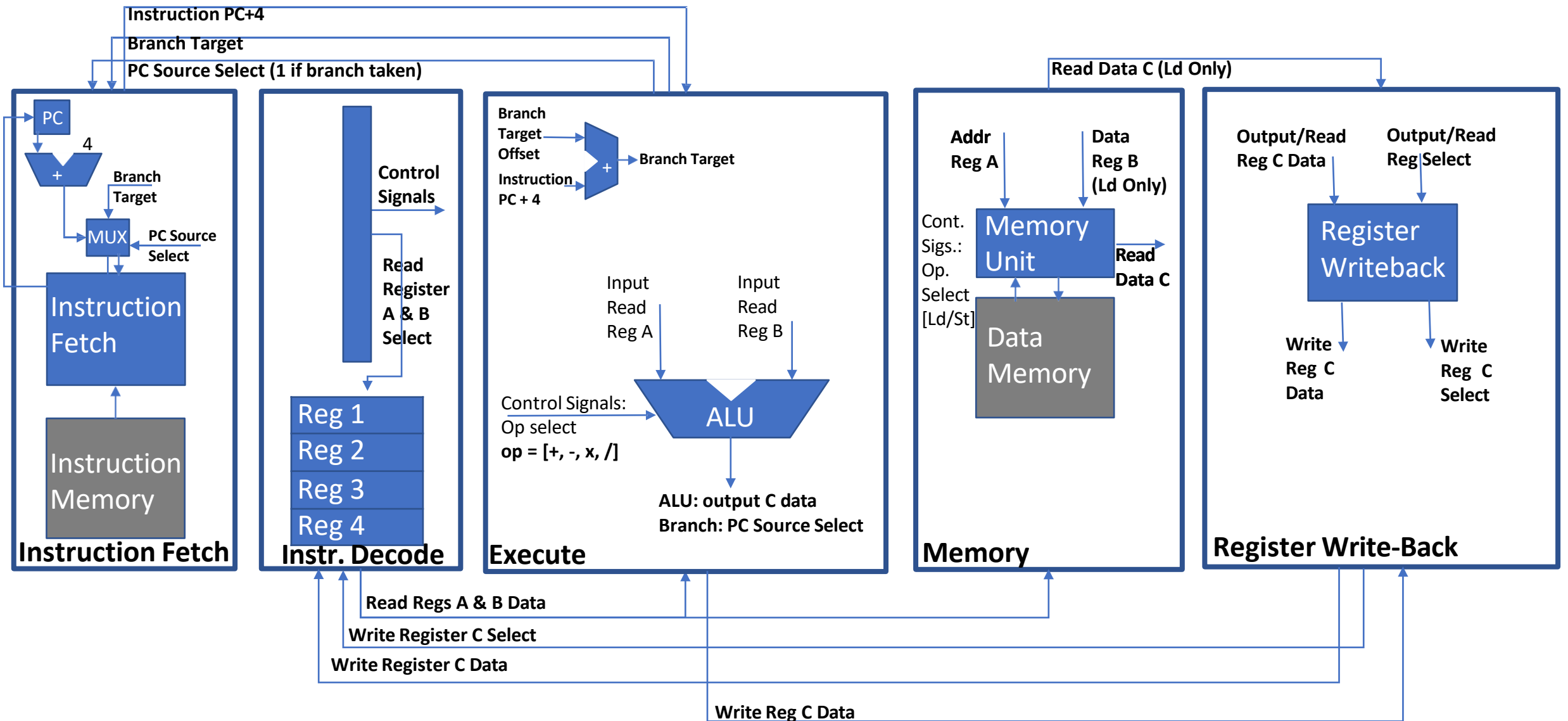
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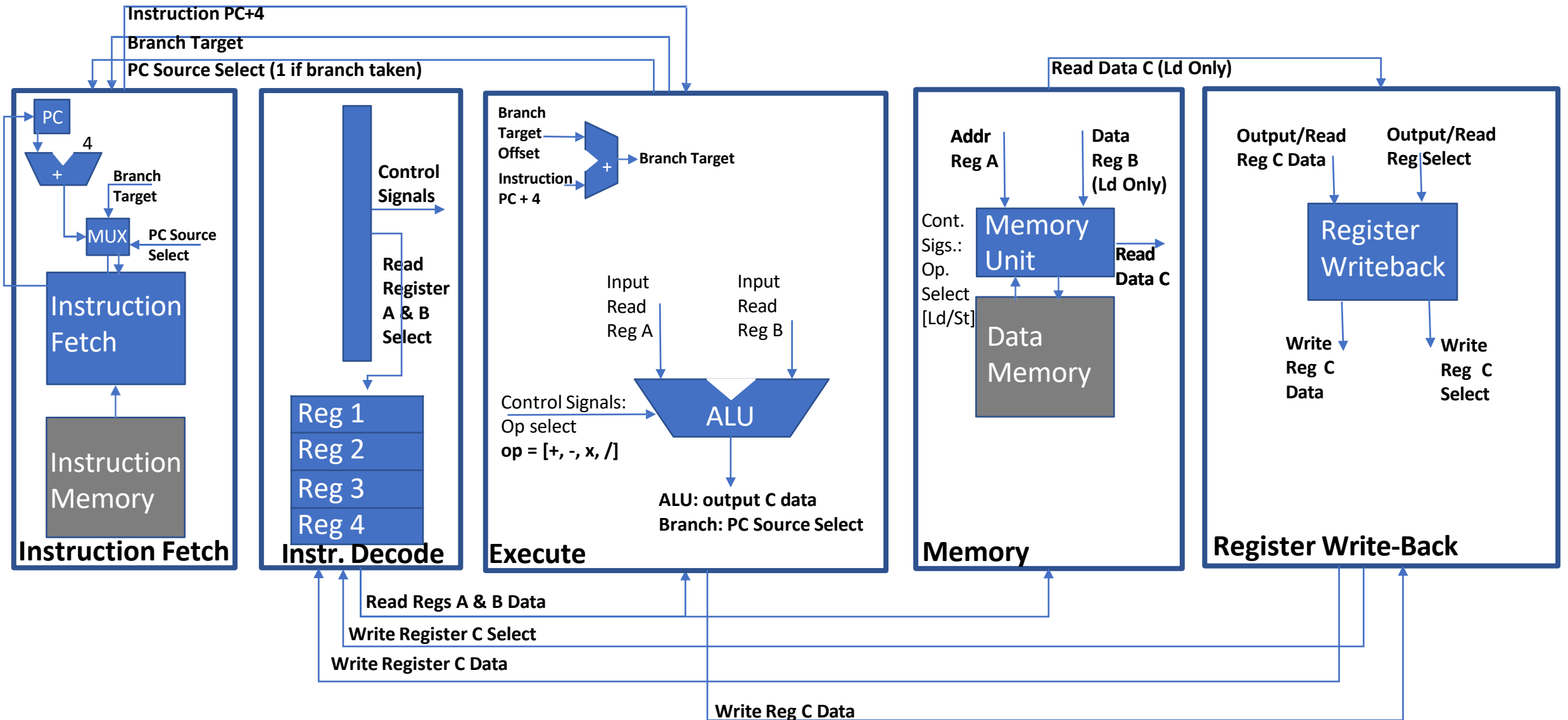
Let's do some grouping together of functionality



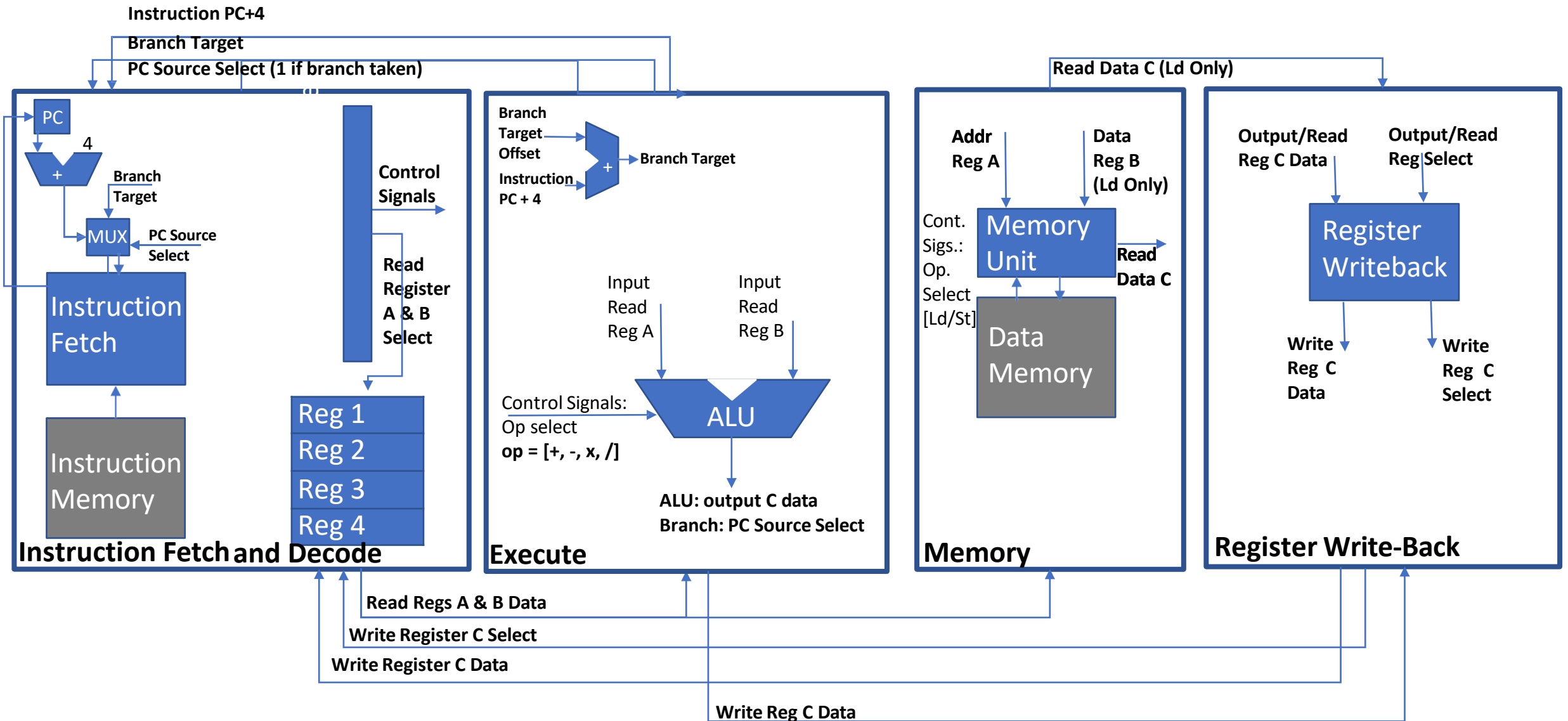
A Simple 5-Stage Pipelined Processor Datapath



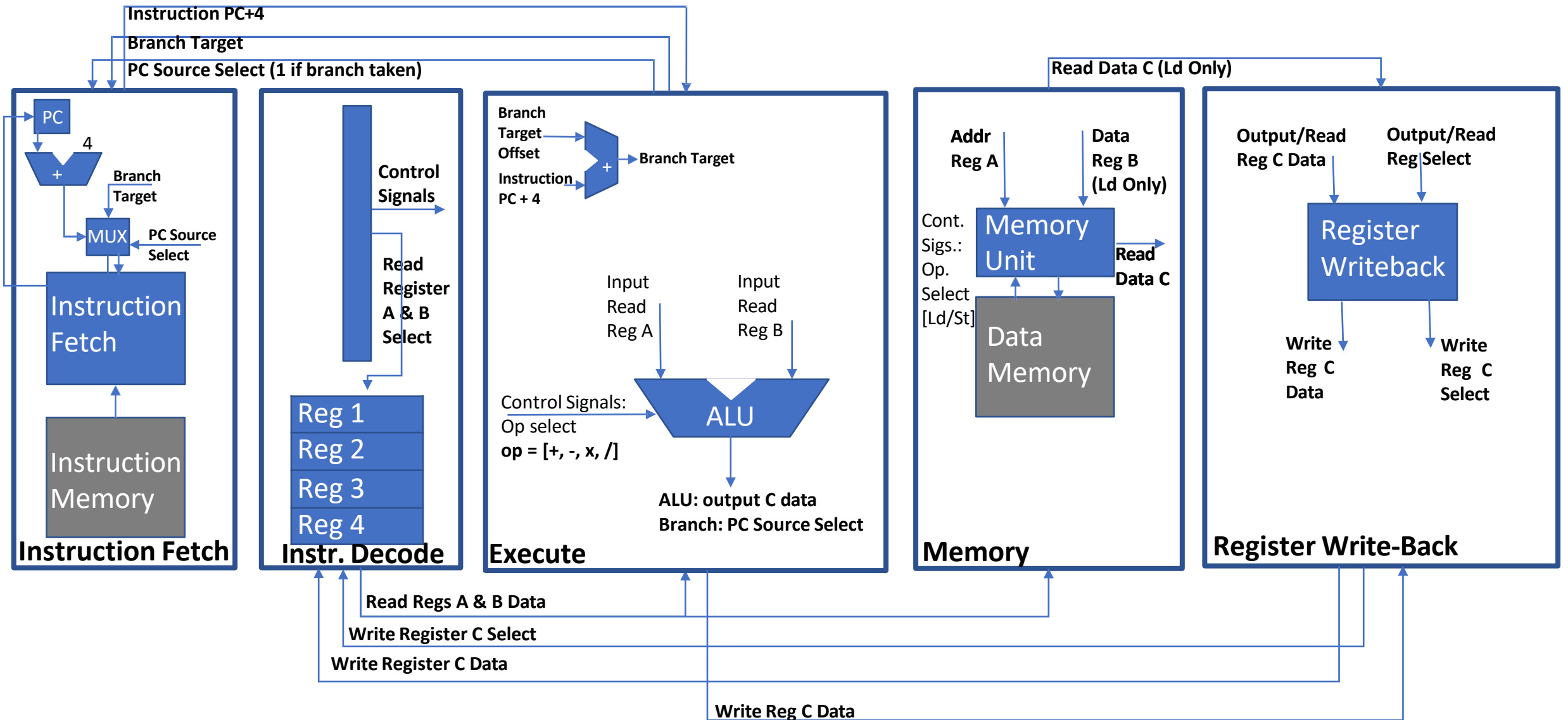
What about an alternative decomposition?



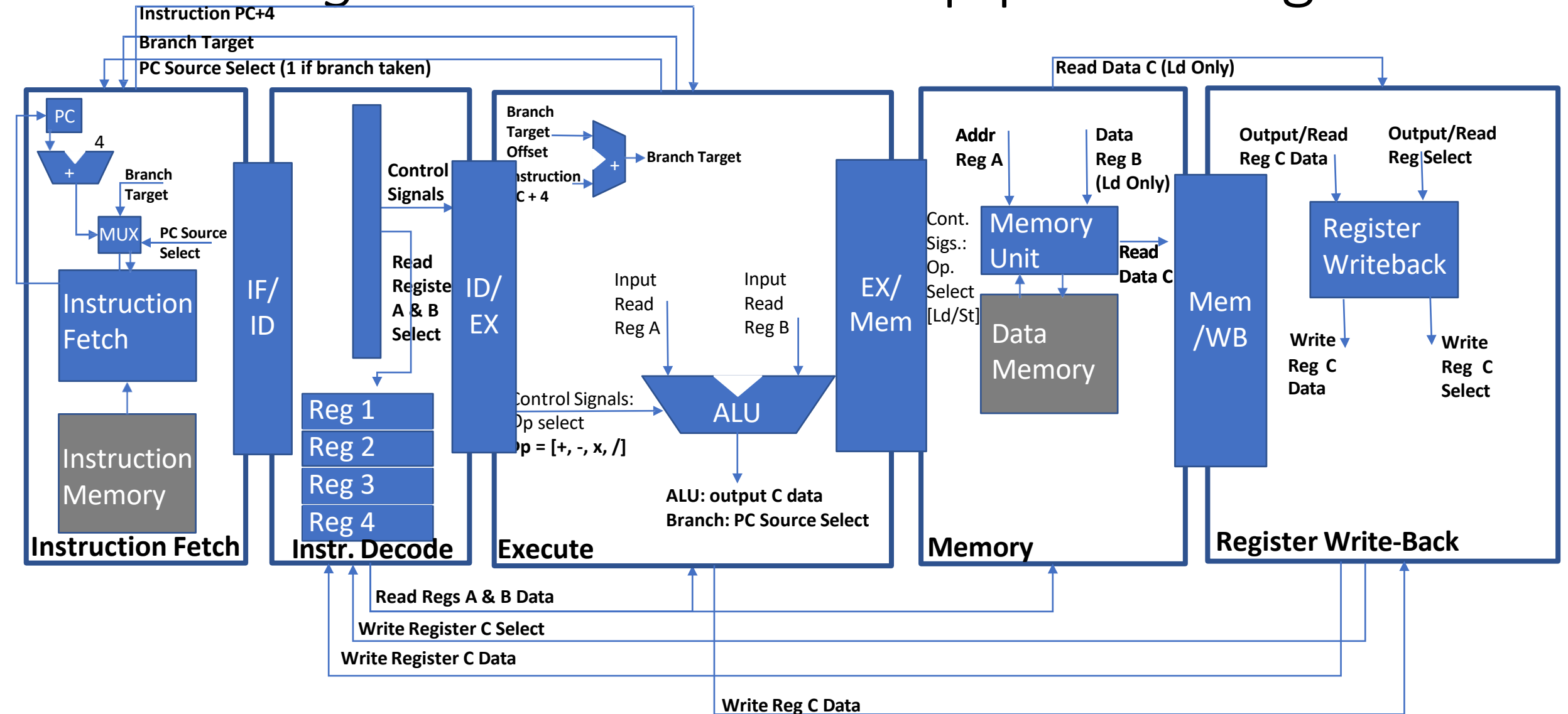
4-stage? Pro / con?



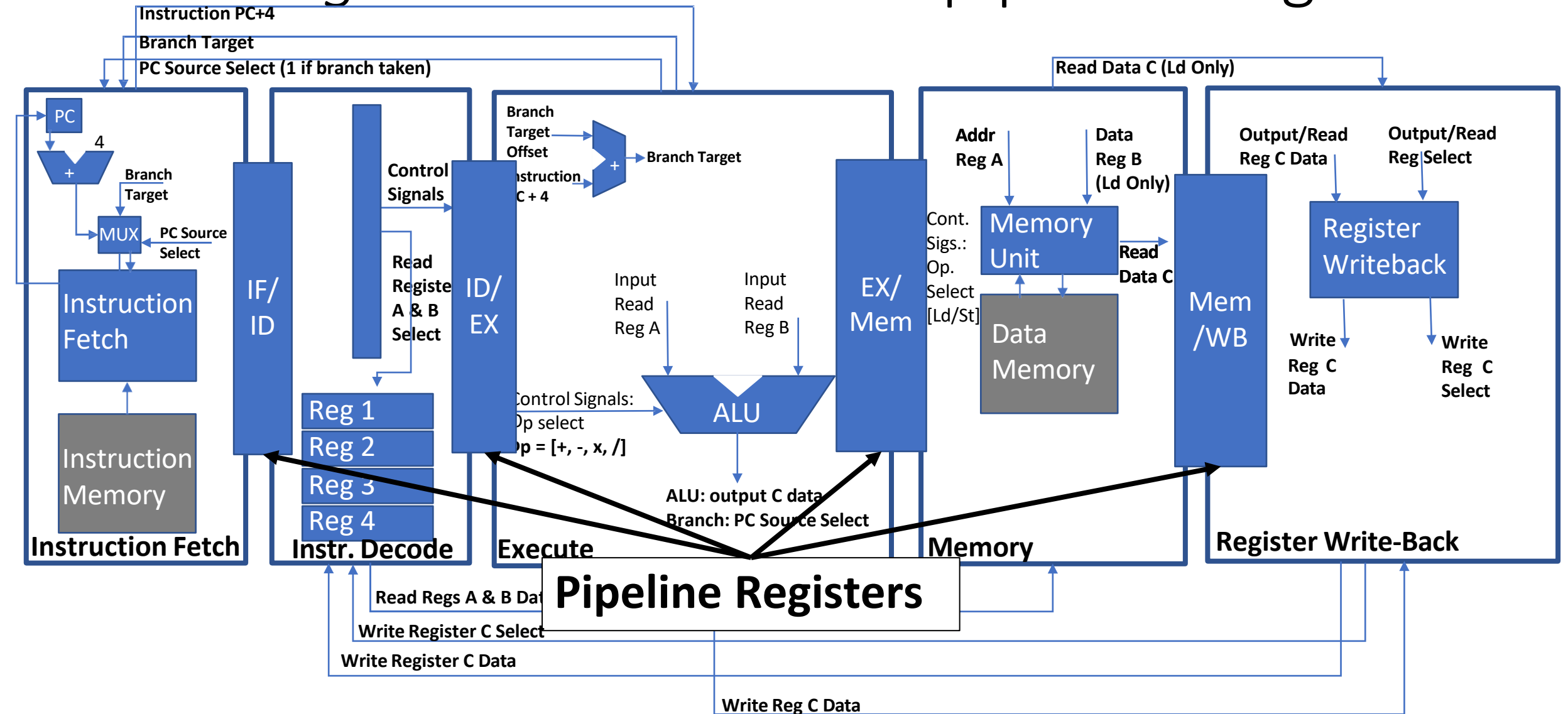
What does ALU op do in Mem? Memop in EX?



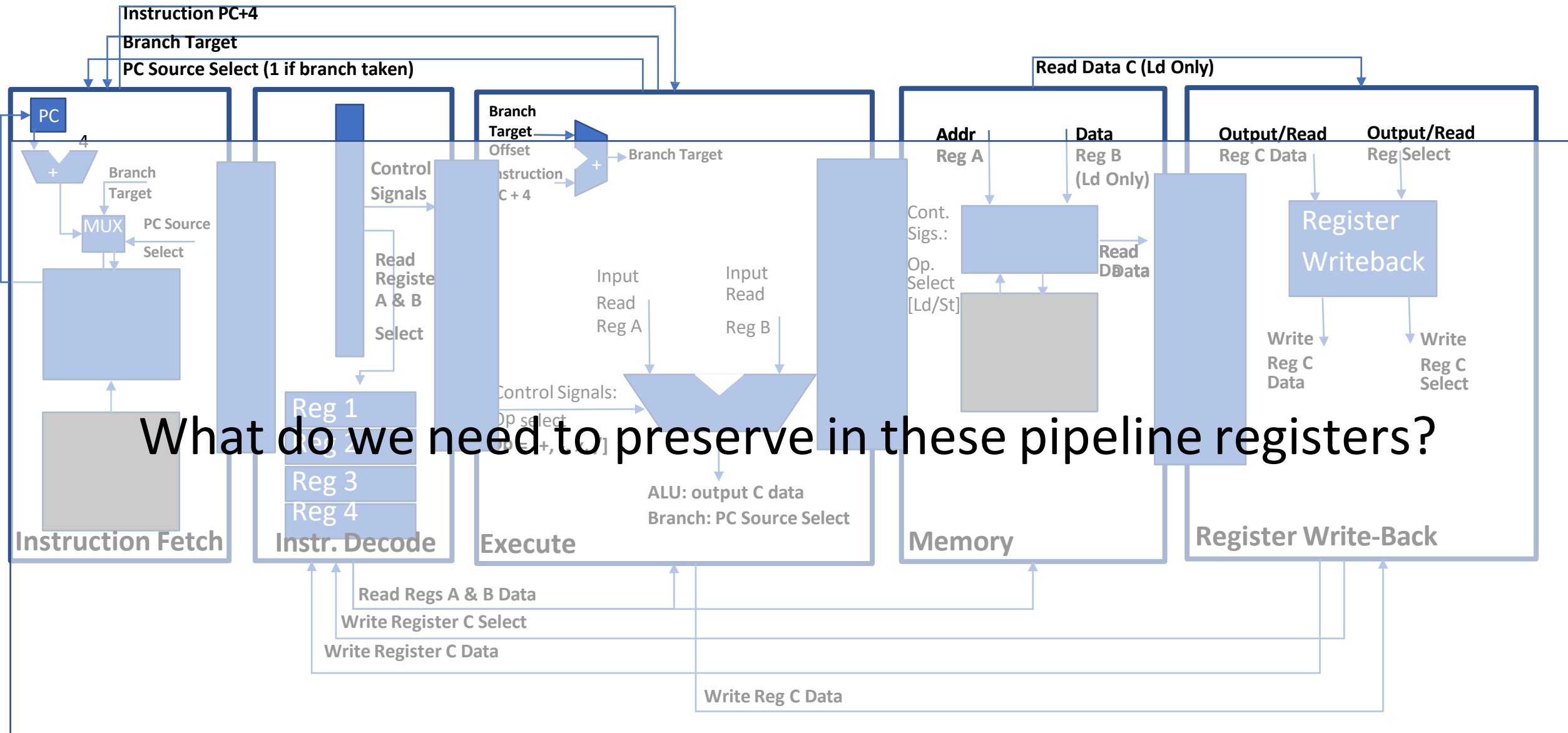
Cost of pipelining: Need to *register* state between pipeline stages



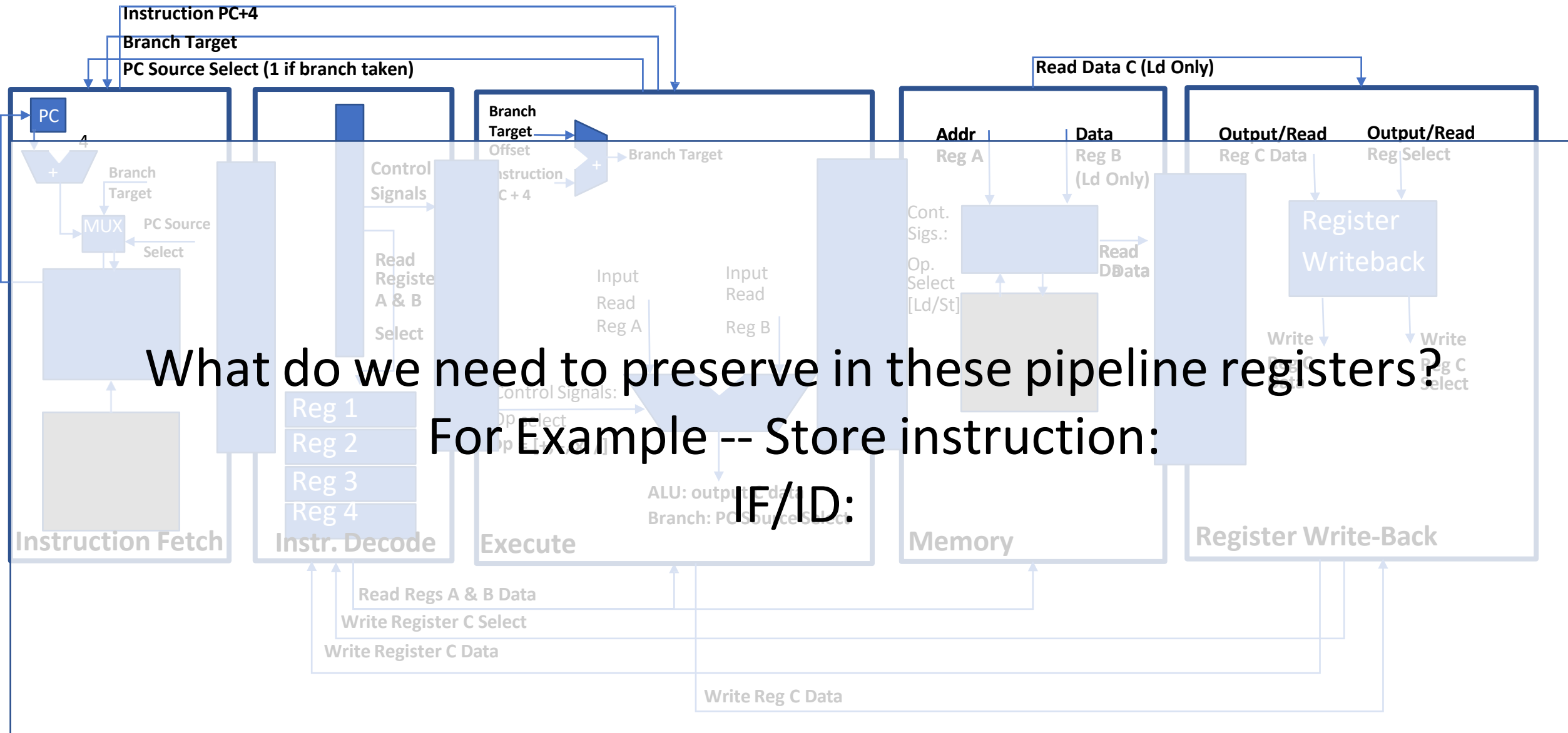
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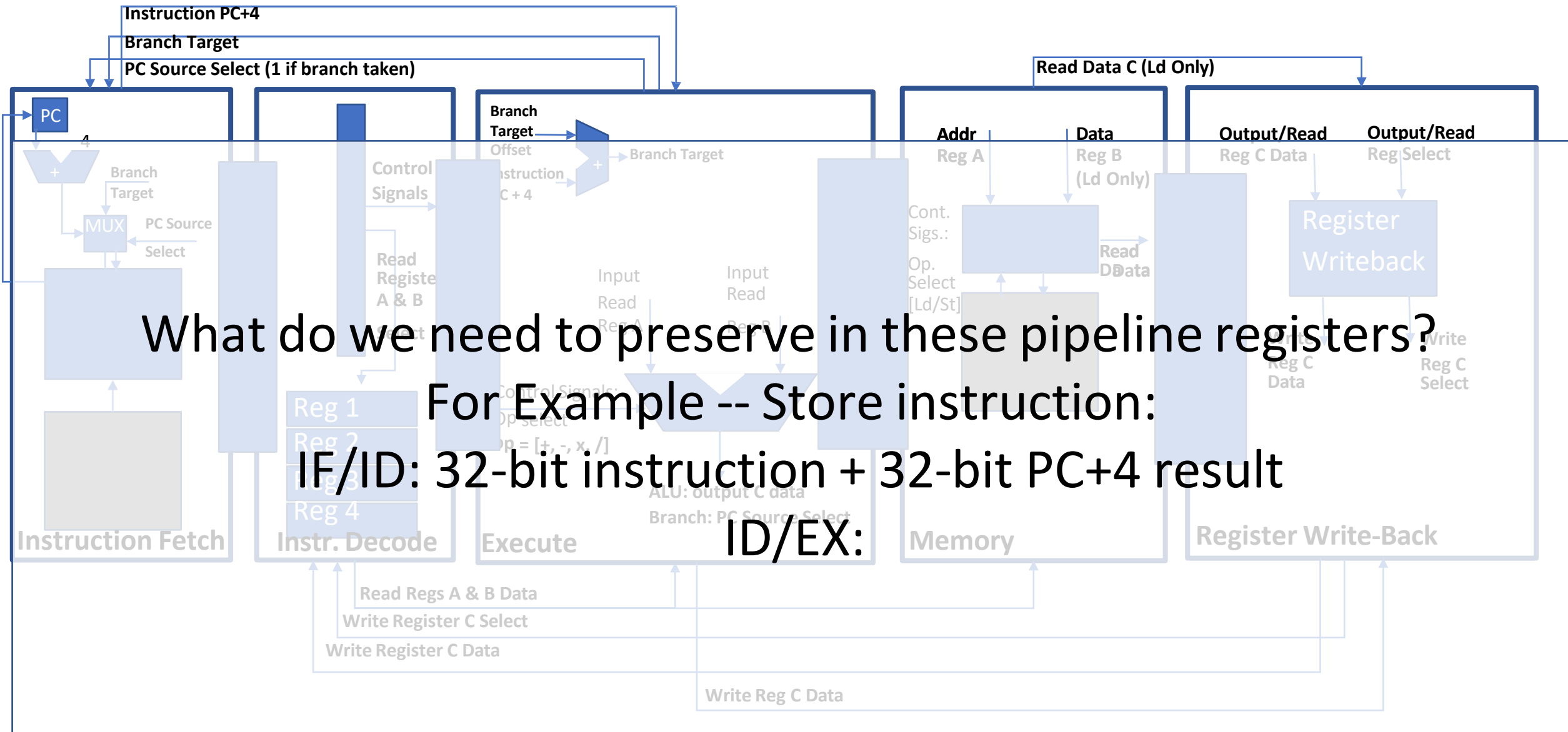
Need to store state between pipeline stages



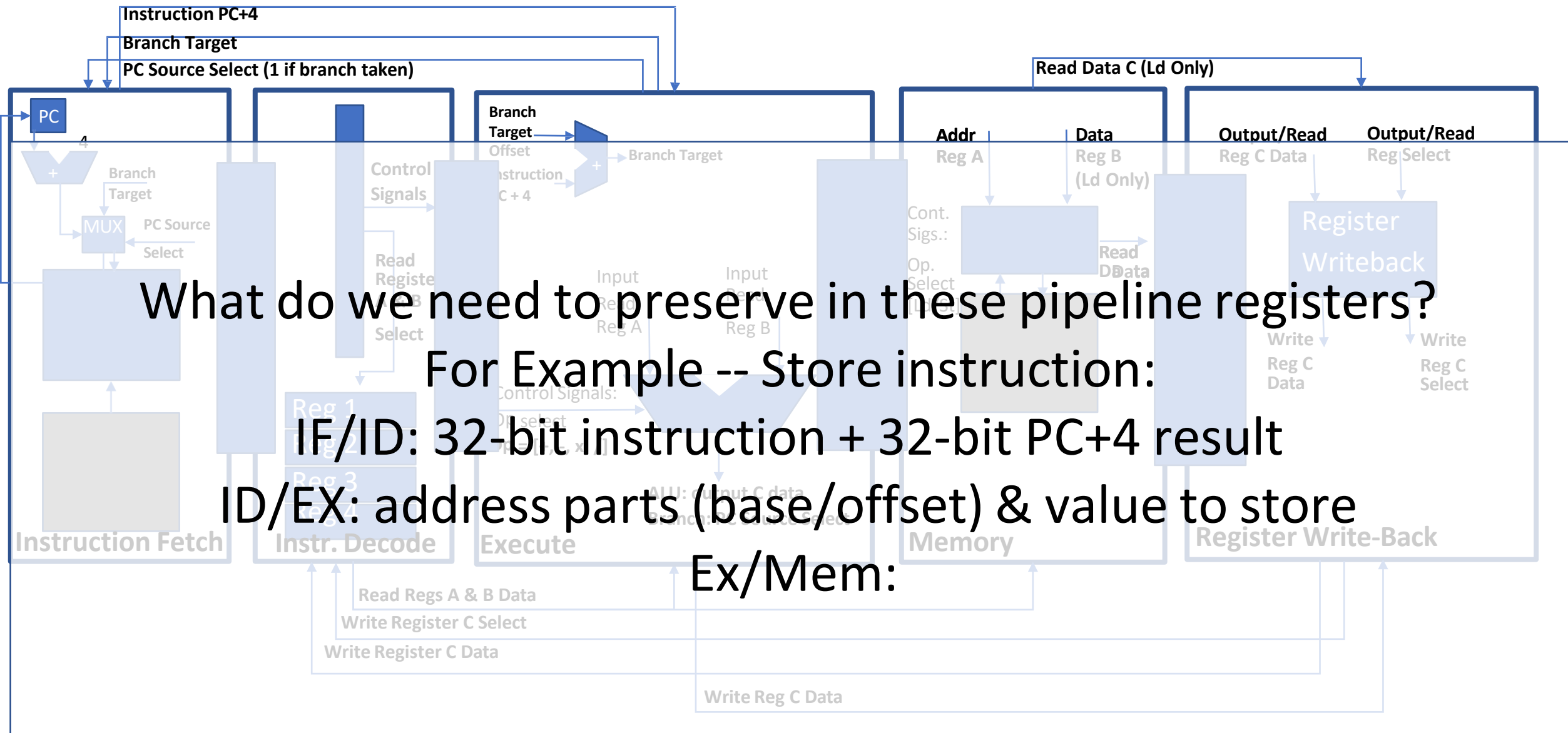
Need to store state between pipeline stages



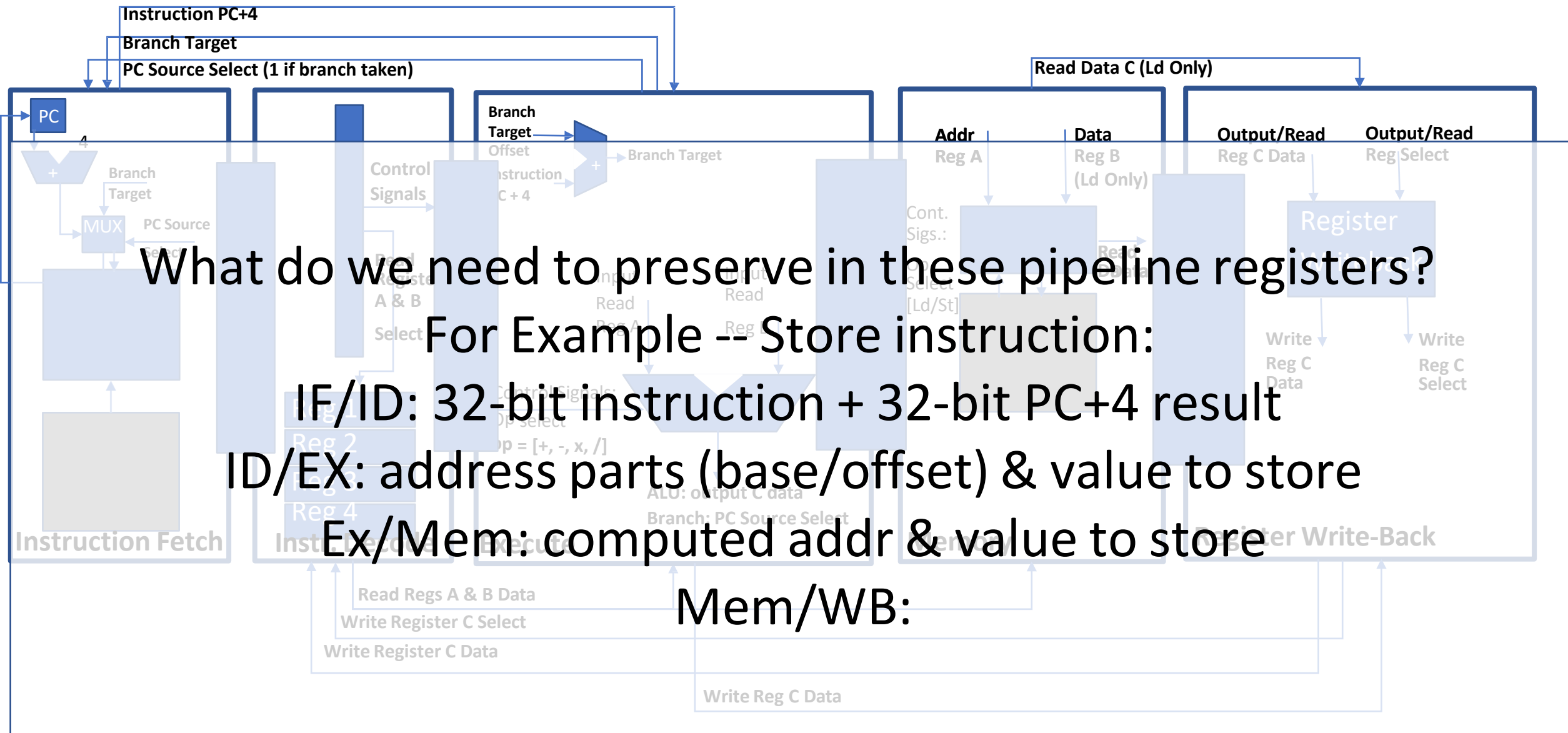
Need to store state between pipeline stages



Need to store state between pipeline stages



Need to store state between pipeline stages



What do we need to preserve in these pipeline registers?

For Example -- Store instruction:

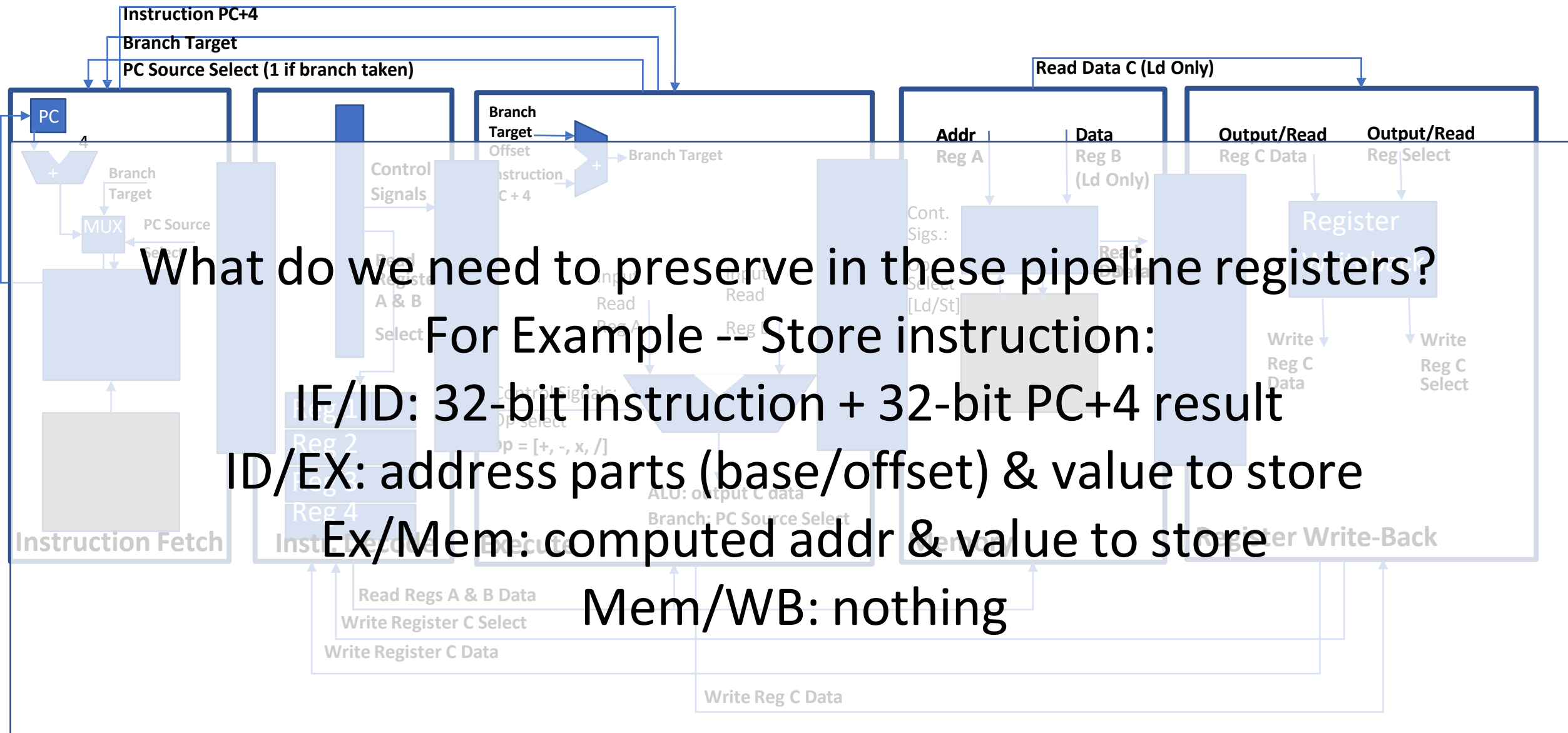
IF/ID: 32-bit instruction + 32-bit PC+4 result

ID/EX: address parts (base/offset) & value to store

Ex/Mem: computed addr & value to store

Mem/WB:

Need to store state between pipeline stages



Example: Pipelined Execution

`add x7 x8 x9`



Example: Pipelined Execution

lw x12 (x15) add x7 x8 x9



Example: Pipelined Execution

sw x0 (x13) lw x12 (x15) add x7 x8 x9

Fetch

Decode

Execute

Memory

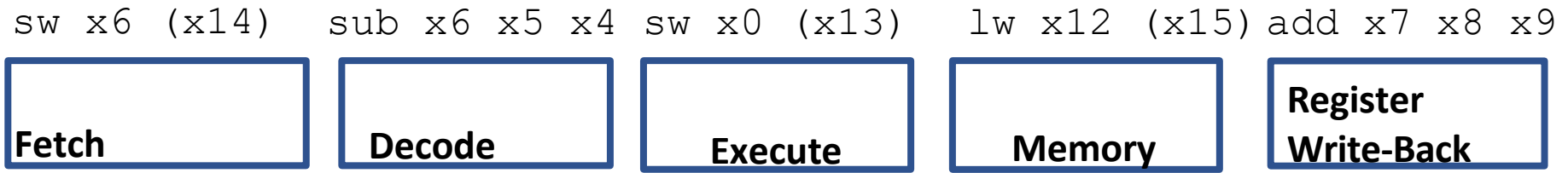
**Register
Write-Back**

Example: Pipelined Execution

sub x6 x5 x4 sw x0 (x13) lw x12 (x15) add x7 x8 x9



Example: Pipelined Execution



Example: Pipelined Execution



sw x6 (x14)

sub x6 x5 x4

sw x0 (x13)

lw x12 (x15)

add x7 x8 x9

Fetch

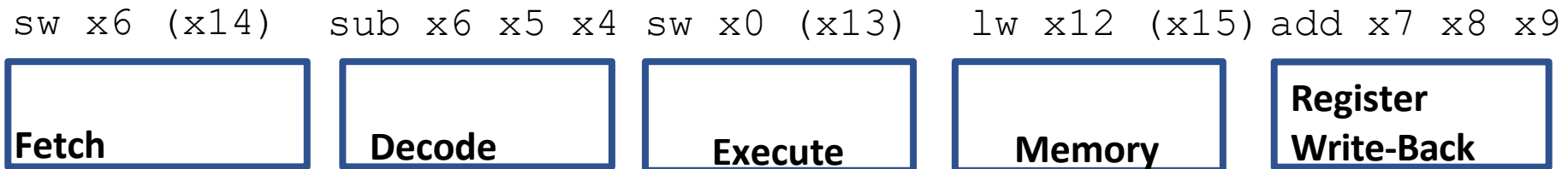
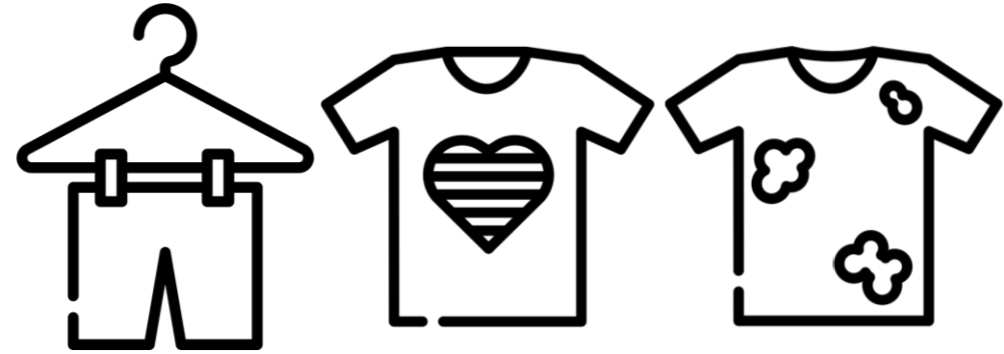
Decode

Execute

Memory

**Register
Write-Back**

Example: Pipelined Execution



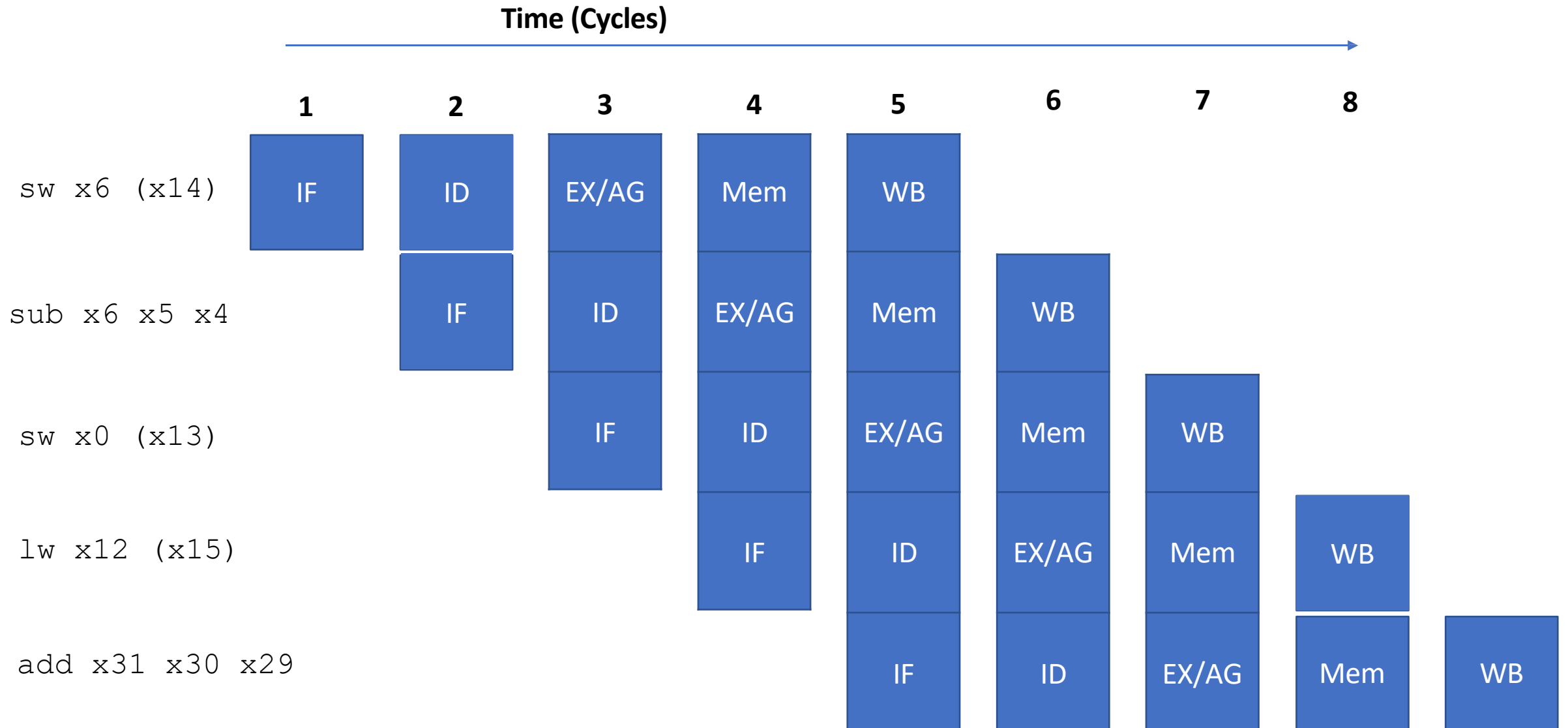
Key Idea: Pipelining unlocks

Instruction Level Parallelism (ILP)

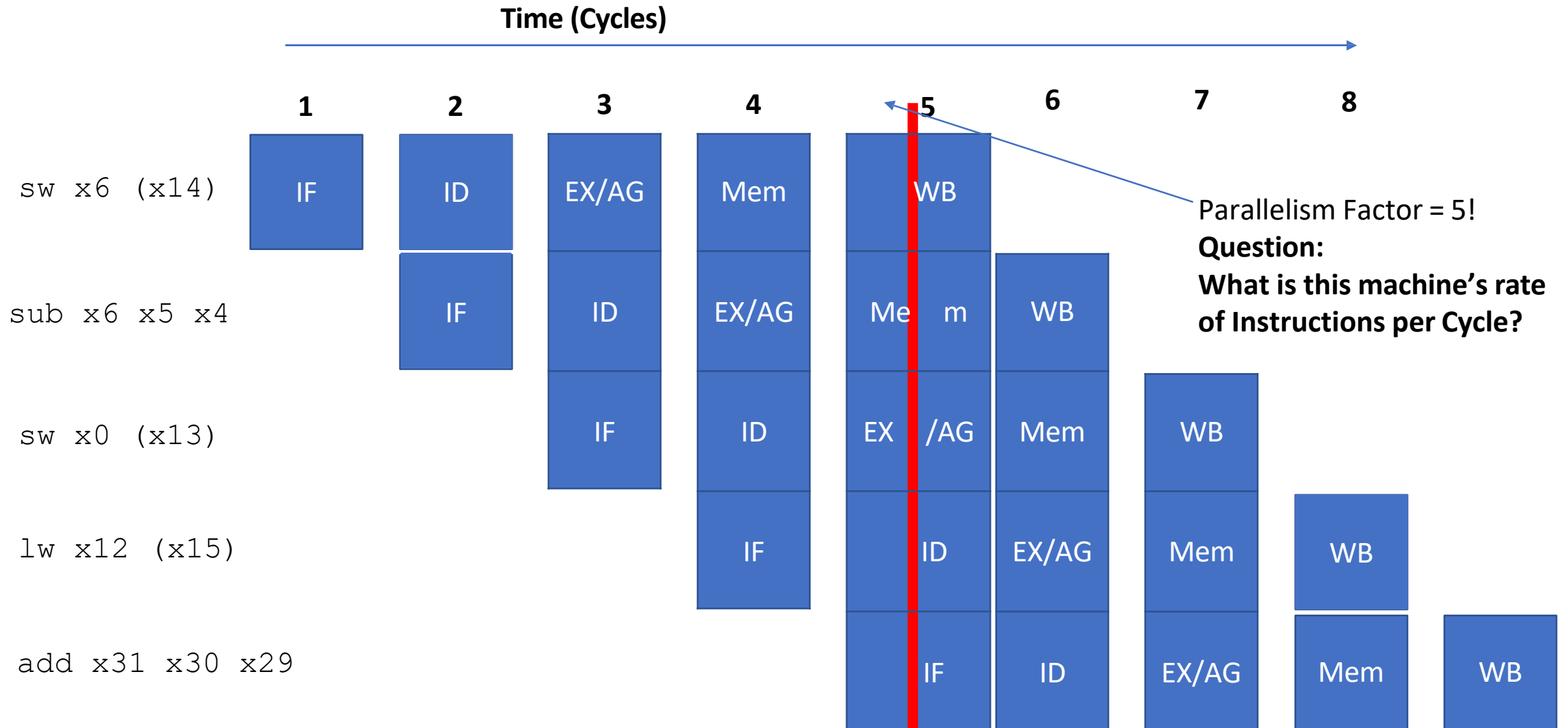
one of the great ideas in computer architecture

Practical Implications of adding ILP to the system?

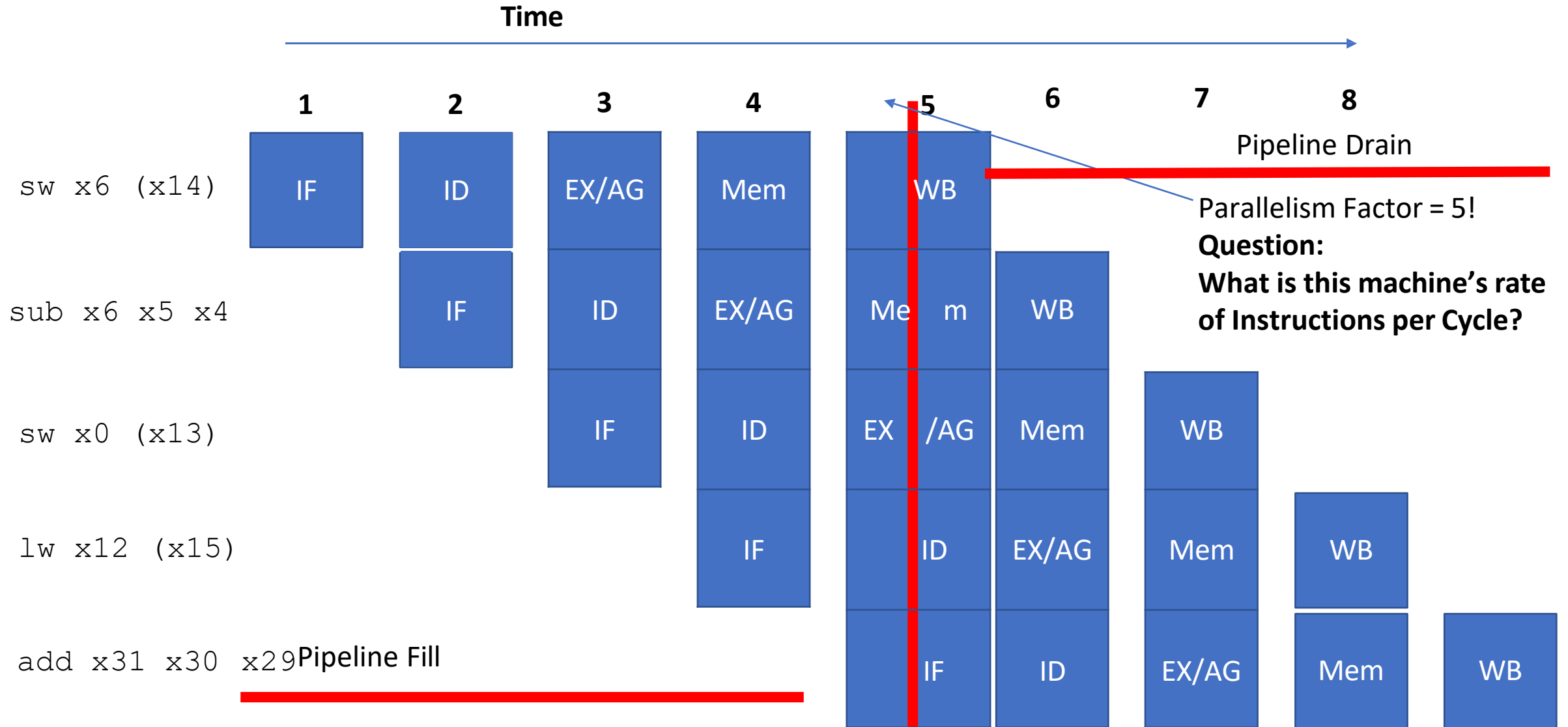
Pipeline Diagram Illustrates Parallelism



Pipeline Diagram Illustrates Parallelism



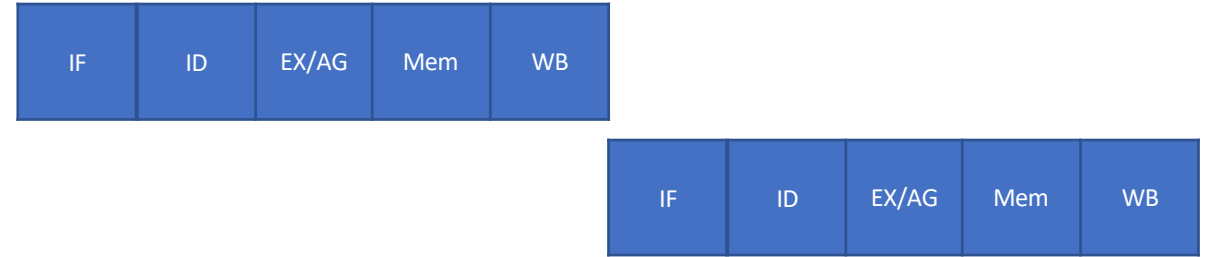
Pipeline Diagram Illustrates Parallelism



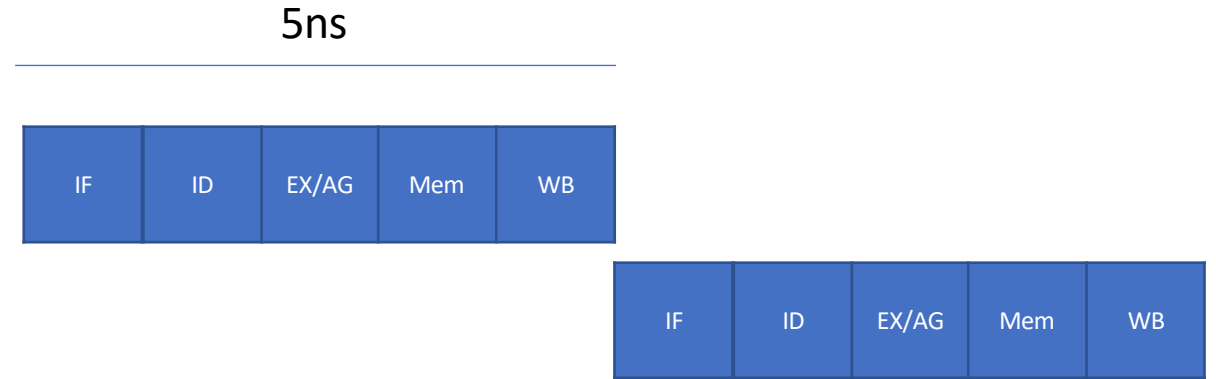
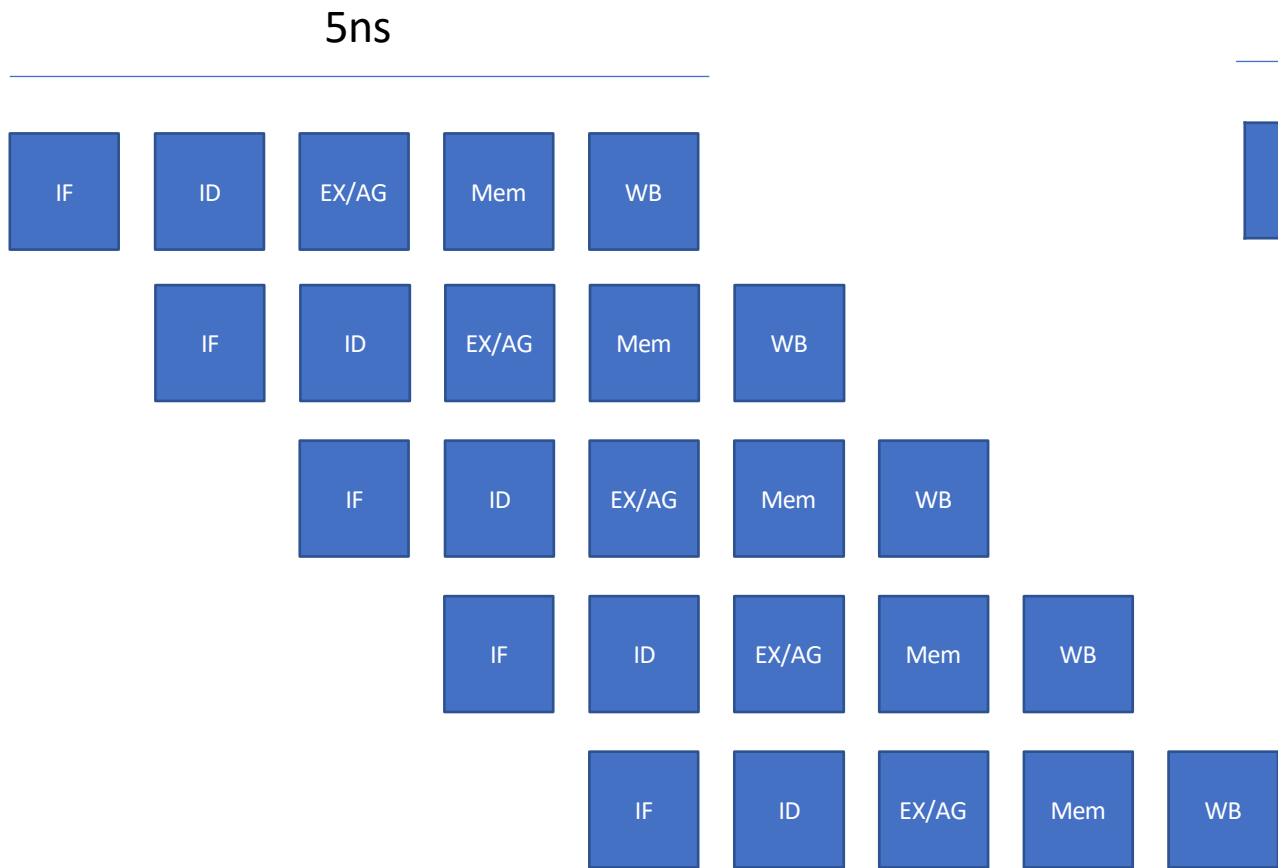
5ns



5ns



What gives? IPC is 1 for both and each instruction's *latency* is still 5ns.



**Instructions/Cycle
Vs
Instructions/Second**

**5ns = (up to) 25(*) stages of work
VS
5ns = 5 stages of work**

*** 15 here due to pipeline filling**

What gives? IPC is 1 in both cases!

Key Idea: Pipelined *Instruction Throughput* is higher.

**Shorter clock period + parallelism = 1 completed instruction per ns
even though *each* instruction takes 5ns to complete**

Iron Law of Computer Performance

instructions
/ program

X

cycles /
instruction

X

seconds
/ cycle

Iron Law of Computer Performance

instructions / program X cycles / instruction X seconds / cycle

Question: what term does pipelining optimize? how else might we approach optimization in light of this performance expression?

Pipelining Code Example

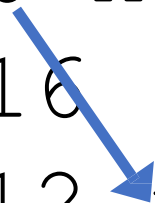
```
p = 0xabc;  
x = y - z  
m = *p;  
t = x + w;
```

```
sub x6 x5 x4  
lw x16 0xabc  
add x12 x6 x14
```

What is interesting about this short program?

Pipelining Code Example

```
sub  x6  x5  x4  
lw   x16 0xabc  
add  x12 x6  x14
```



What happens to `x6` as we execute this code?

Example: Pipelined Execution

sub x6 x5 x4

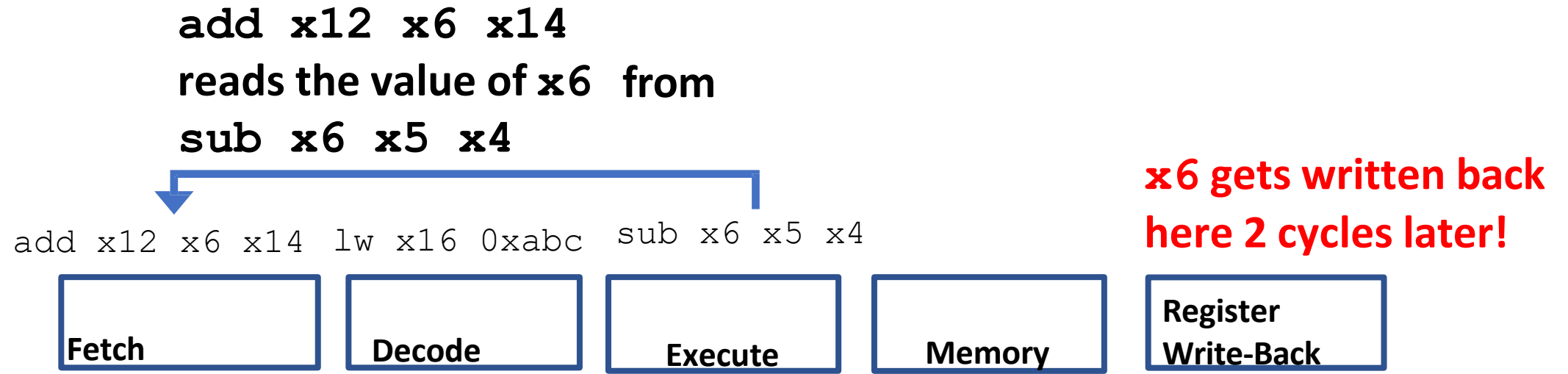


Example: Pipelined Execution

```
lw x16 0xabc sub x6 x5 x4
```



Example: Pipelined Execution



Read-After-Write (RAW) Hazard:

Input register does not contain updated data during register read cycle due to yet-to-be-completed register writeback from older instruction

Types of Data Hazards

```
sub  x6  x5  x4  
lw   x16 0xabc  
add  x12 x6  x14
```

Read-After-Write (RAW)

```
sub  x8  x16 x4  
add  x16 x6  x14  
lw   x16 0xabc
```

Write-After-Read (WAR)

```
lw   x6  0xabc  
sub  x6  x5  x4  
add  x12 x6  x14
```

Write-After-Write (WAW)

Only Read-After-Write (RAW) hazards are possible in our simple pipeline

Types of Data Hazards

```
lw   x6  0xabc  
sub  x6  x5  x4  
add  x12 x6  x14
```

Write-After-Write (WAW)

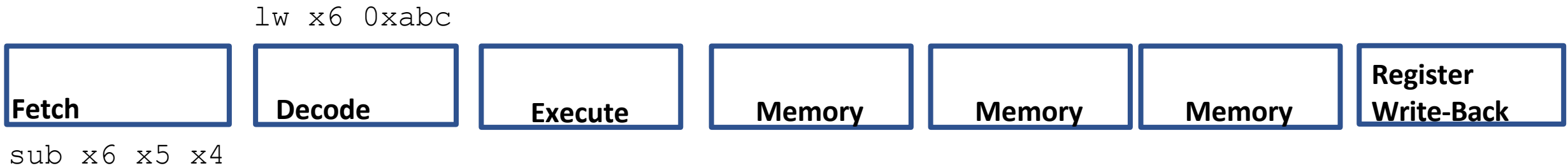
lw x6 0xabc



Types of Data Hazards

```
lw    x6  0xabc
sub   x6  x5  x4
add   x12 x6  x14
```

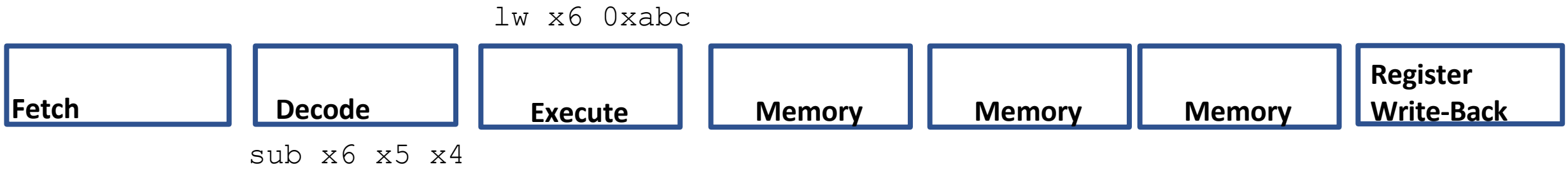
Write-After-Write (WAW)



Types of Data Hazards

```
lw    x6  0xabc
sub   x6  x5  x4
add   x12 x6  x14
```

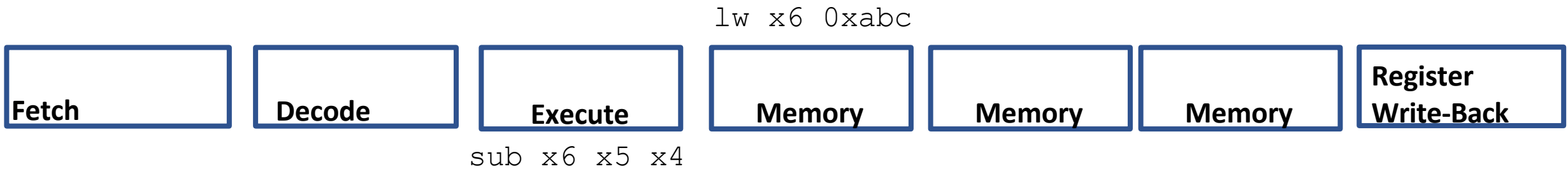
Write-After-Write (WAW)



Types of Data Hazards

```
lw   x6  0xabc
sub  x6  x5  x4
add  x12 x6  x14
```

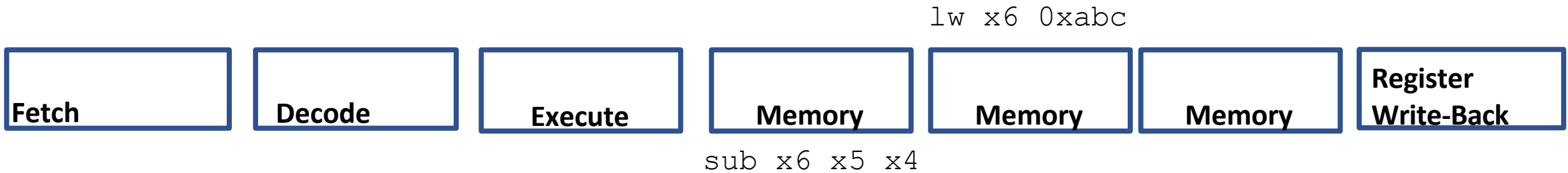
Write-After-Write (WAW)



Types of Data Hazards

```
lw   x6  0xabc
sub  x6  x5  x4
add  x12 x6  x14
```

Write-After-Write (WAW)



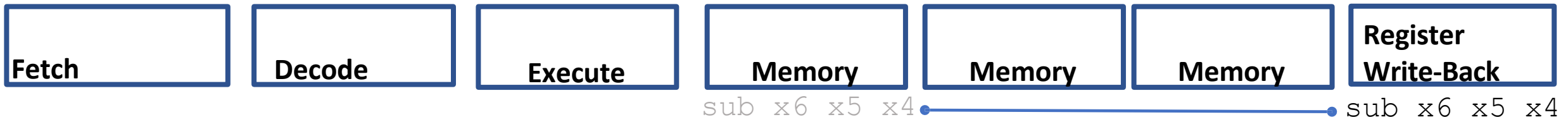
Types of Data Hazards

```
lw   x6  0xabc
sub  x6  x5  x4
add  x12 x6  x14
```

Write-After-Write (WAW)

Multi-cycle latency memory op

```
lw x6 0xabc lw x6 0xabc lw x6 0xabc
```




Non-mem-op, single memory cycle

Earlier `lw` instruction finishes after later `sub` instruction. Both write `x6`. Wrong final value in `x6`.
Explicitly handled with logic to maintain ordering in processors that allow this behavior (not our datapath)

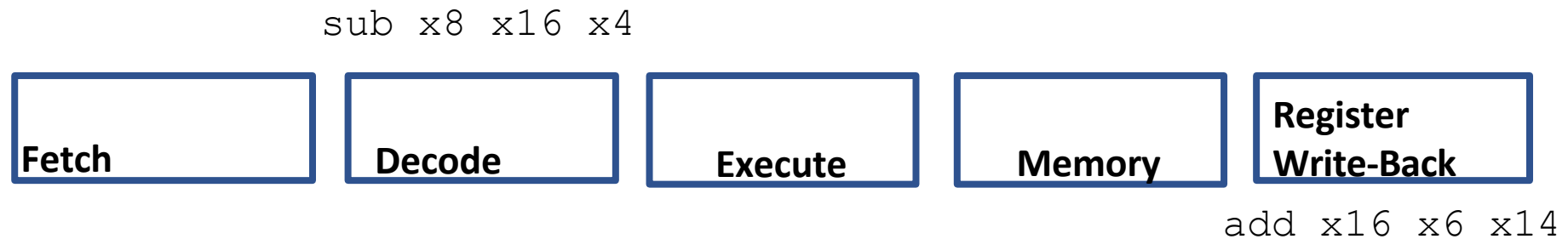
Types of Data Hazards

```
sub  x8  x16  x4
add  x16 x6  x14
lw   x11 0xabc
```



Write-After-Read (WAR)

Stalled at decode/reg. read
(why? wait a few lectures & more in 447)



Completes quickly and writes reg.

Later add instruction writes x16 before earlier sub instruction reads x16. sub sees wrong value!

What can we do about these data hazards?

```
sub  x6  x5  x4  
lw   x16 0xabc  
add  x12 x6  x14
```

Read-After-Write (RAW)

```
sub  x8  x16 x4  
add  x16 x6  x14  
lw   x16 0xabc
```

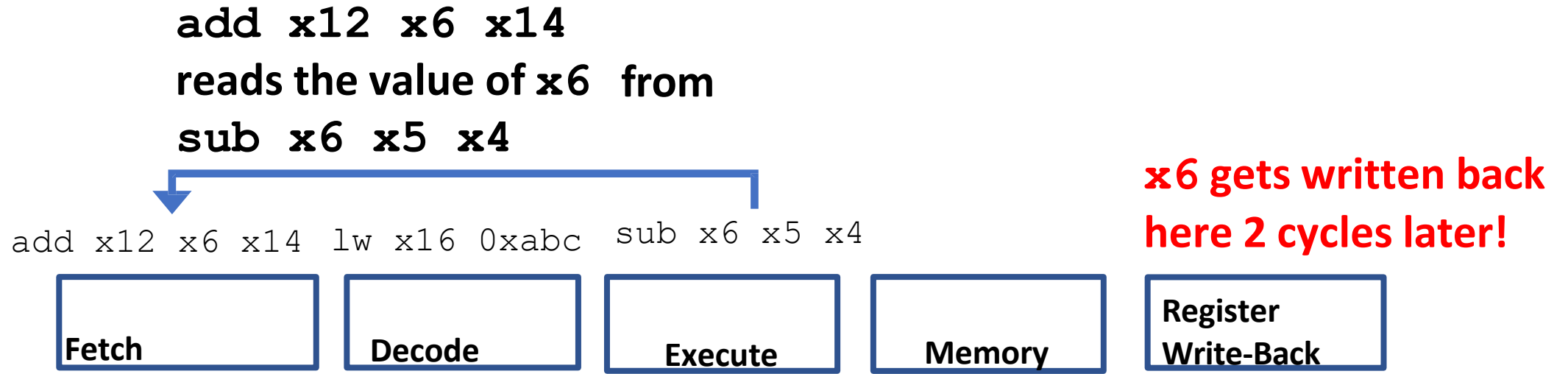
Write-After-Read (WAR)

```
lw   x6  0xabc  
sub  x6  x5  x4  
add  x12 x6  x14
```

Write-After-Write (WAW)

Only Read-After-Write (RAW) hazards are possible in our simple pipeline

Example: Pipelined Execution w/ RAW Hazard



Read-After-Write (RAW) Hazard:

Input register does not contain updated data during register read cycle due to yet-to-be-completed register writeback from older instruction

Example: Pipelined Execution w/ RAW Hazard

add x12 x6 x14 sub x6 x5 x4

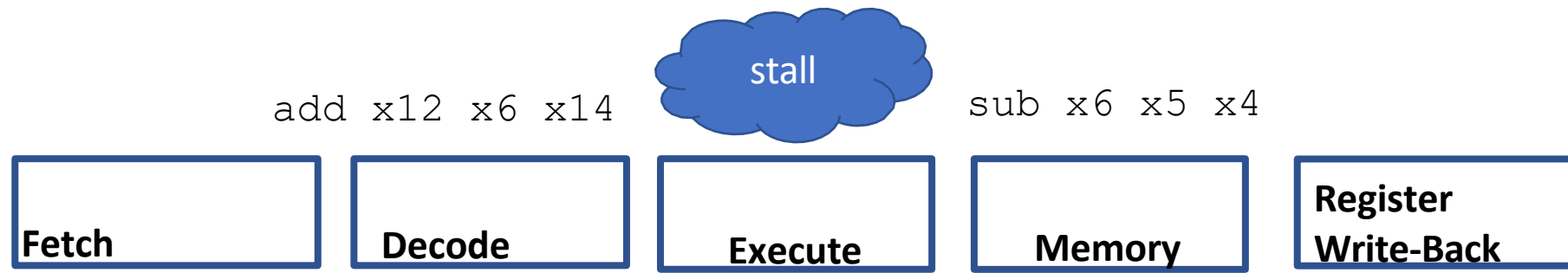


Example: Pipelined Execution w/ RAW Hazard

add x12 x6 x14 sub x6 x5 x4



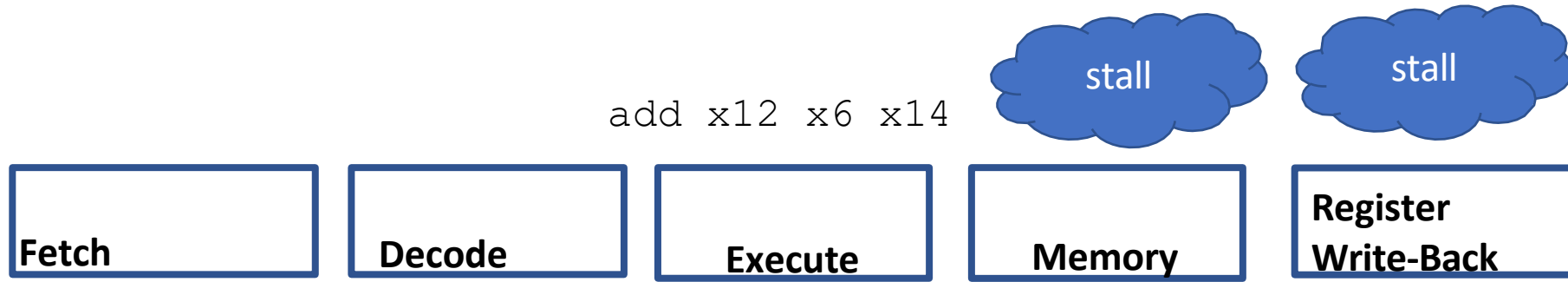
Example: Pipelined Execution w/ RAW Hazard



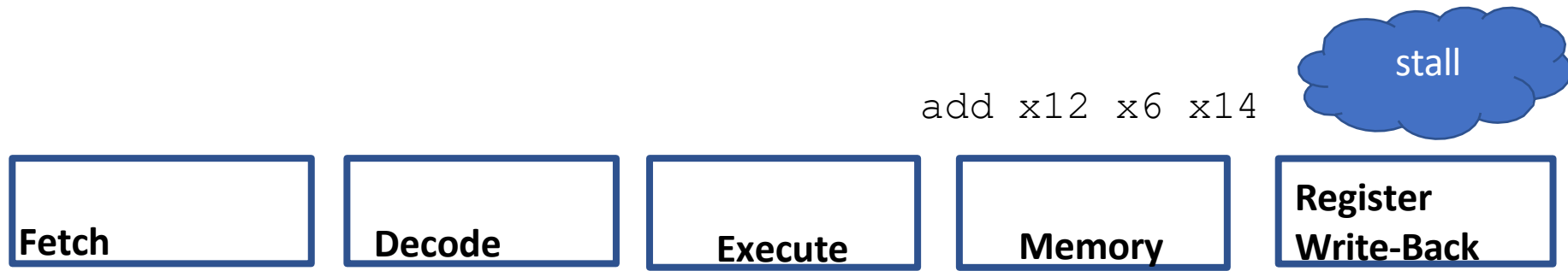
Example: Pipelined Execution w/ RAW Hazard



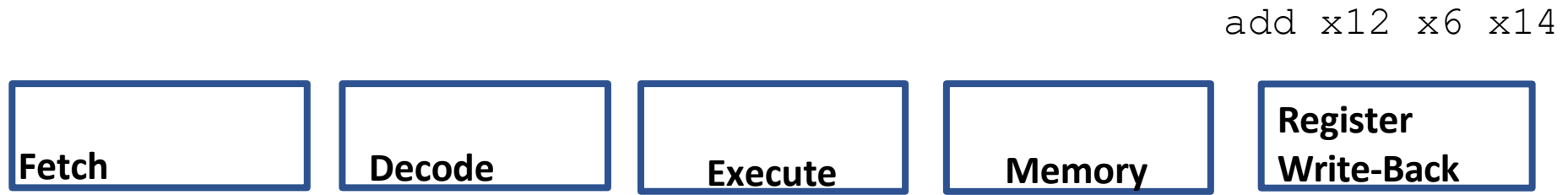
Example: Pipelined Execution w/ RAW Hazard



Example: Pipelined Execution w/ RAW Hazard



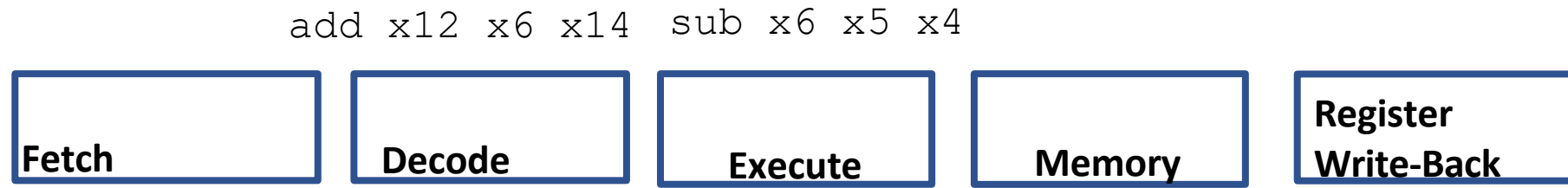
Example: Pipelined Execution w/ RAW Hazard



How do we avoid the stall cycles?



Example: Pipelined Execution w/ RAW Hazard

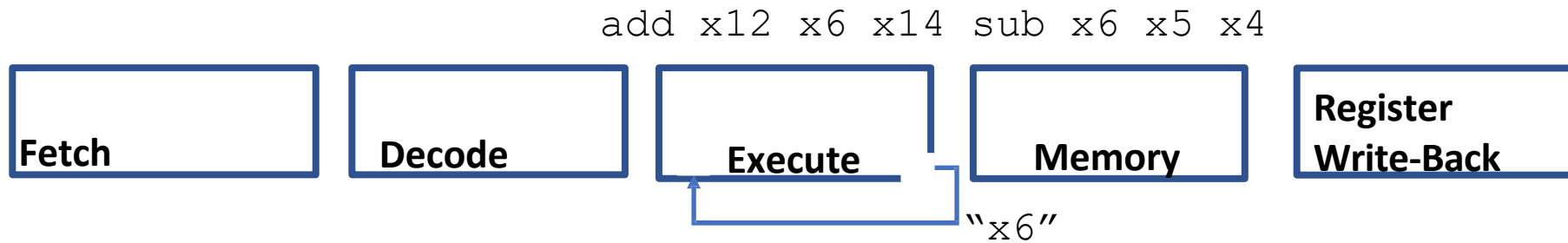


Value of x6 is available after `sub` Executes

We can forward the value to the add!

Forwarding to avoid a pipeline RAW Hazard

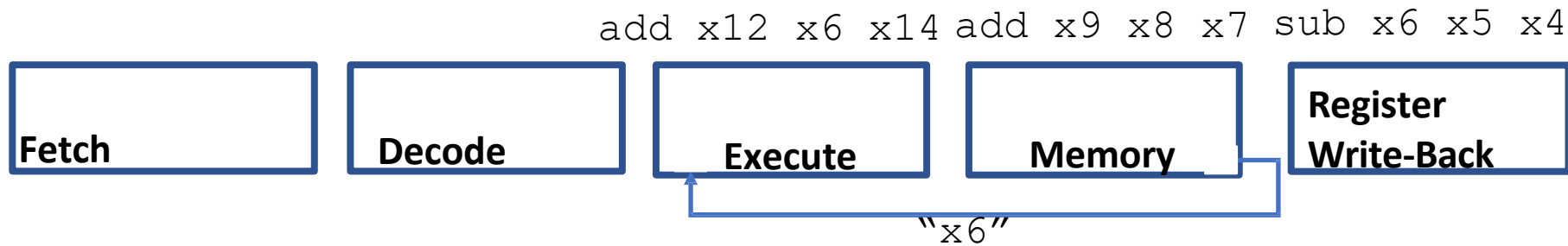
Value of x6 is available from Execute!



We can *forward* the value in the EX/MEM pipeline register from the sub back to Execute to act as the input operand for the add

Forwarding to avoid a pipeline RAW Hazard

Can also forward if there are intervening instructions

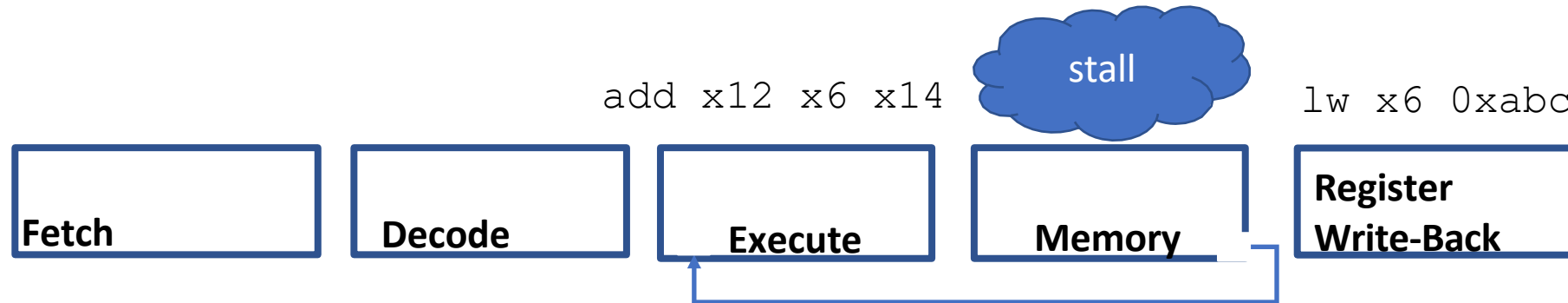


We can *forward* the value in the MEM/WB pipeline register from the sub back to Execute to act as the input operand for the add (going around the unrelated operation in the memory stage)

Pipeline Can Forward Between Different Stages

```
lw x6 0xabc  
add x12 x6 x14
```

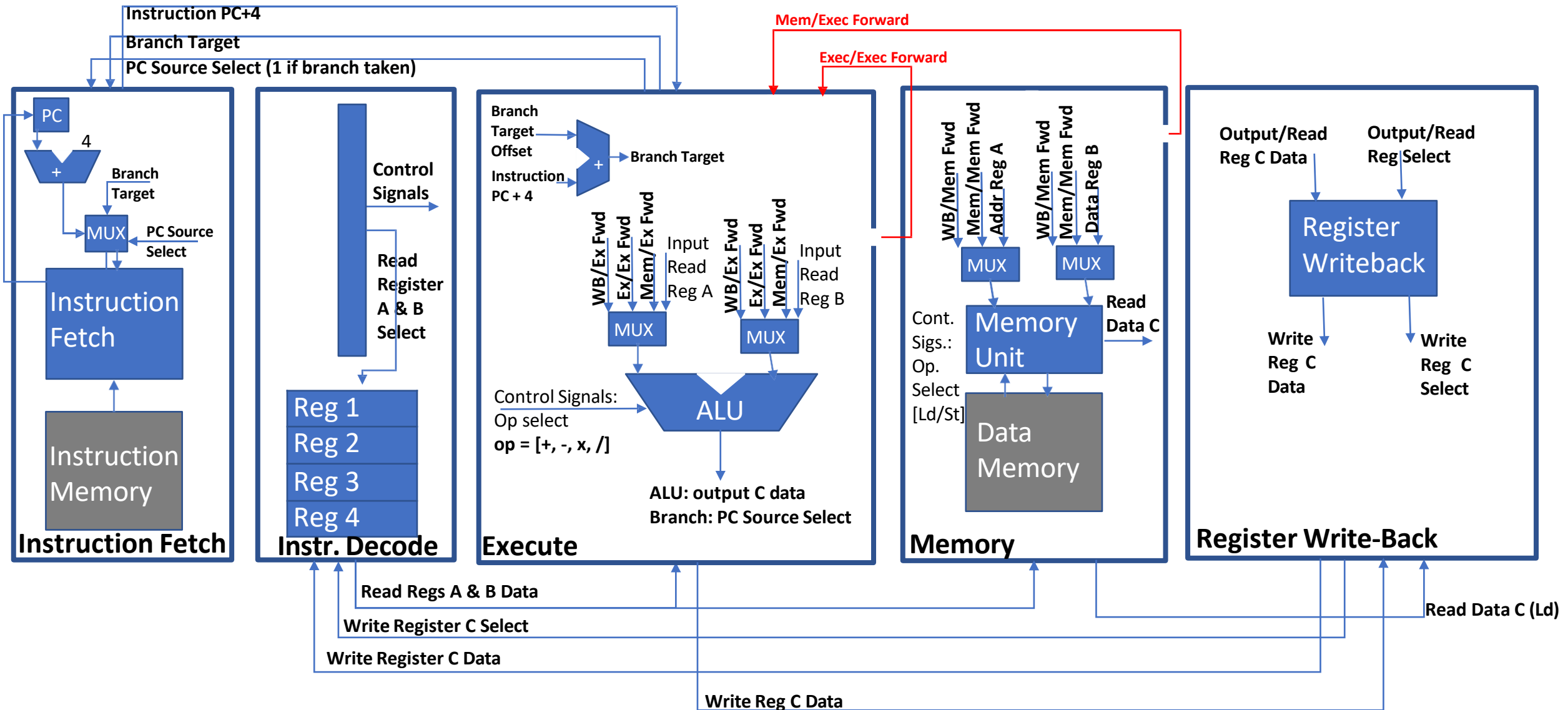
Value of x6 is available from Memory!



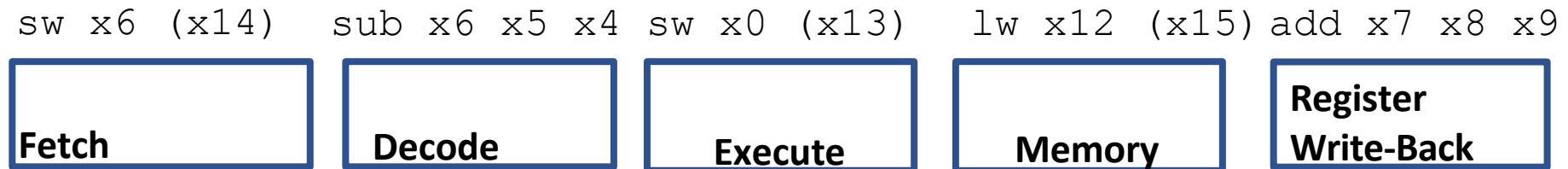
We can *forward* the value in Memory's pipeline register from the lw back to Execute's input for the add

(Still requires stalling...)

Adding Forwarding Support

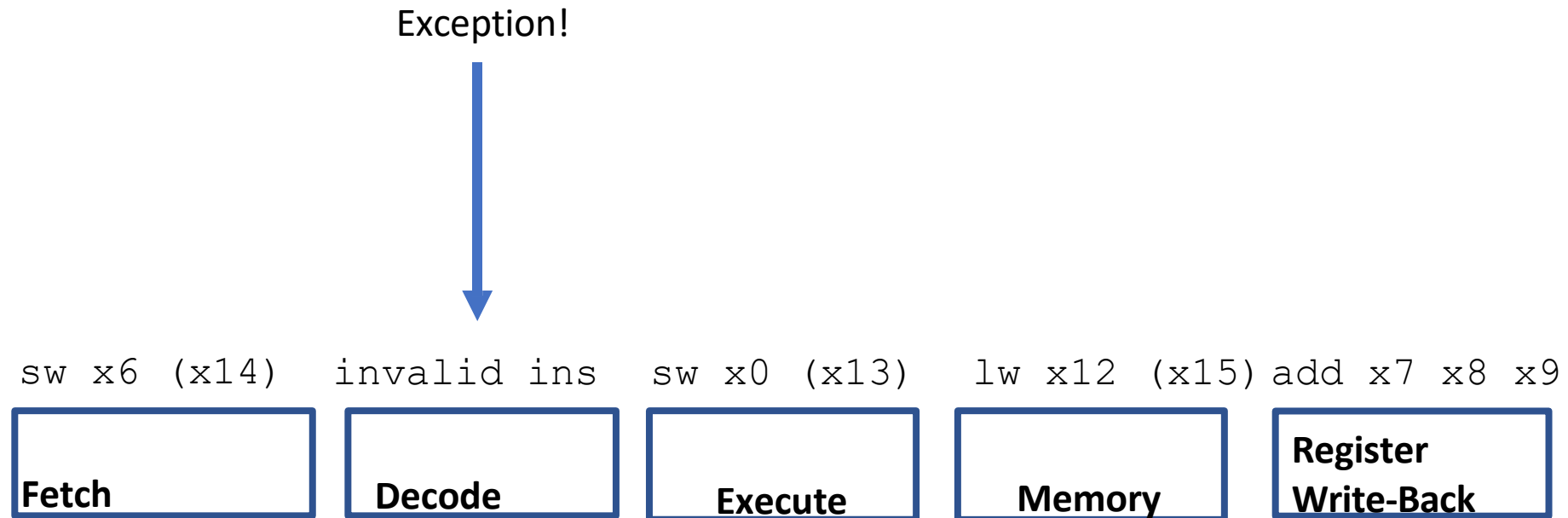


Question: What is time in a pipelined system?



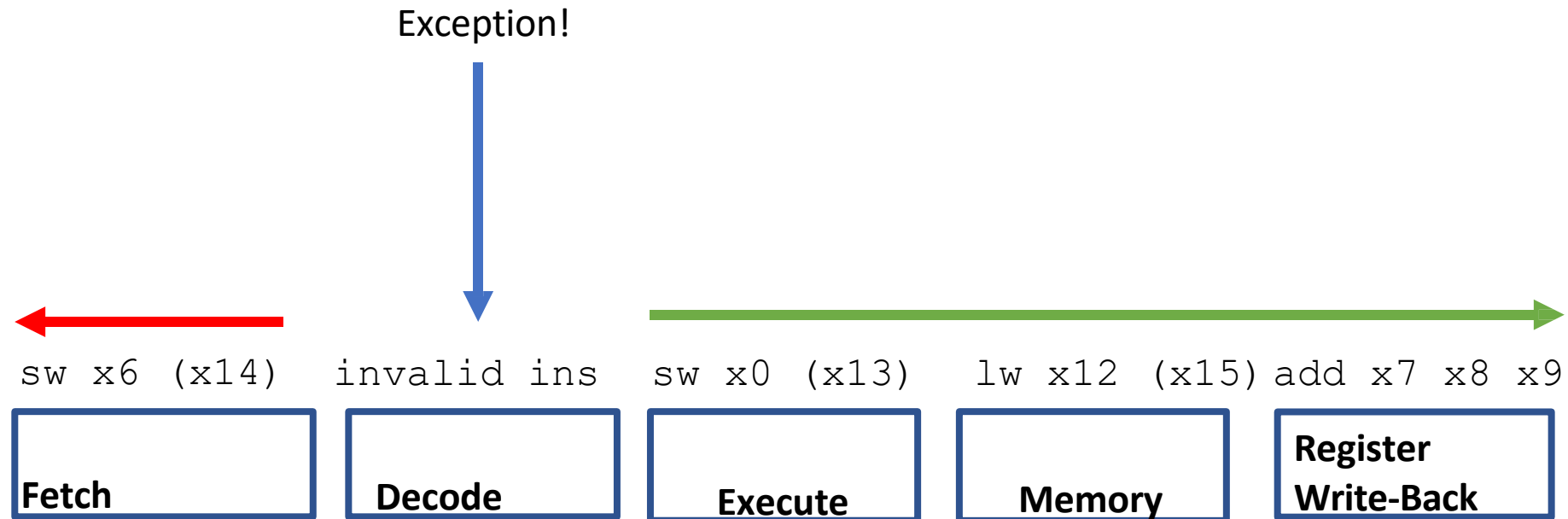
What if one of our instructions were to throw an exception (e.g., illegal instruction in decode or page fault on a memop)?

Exception Handling



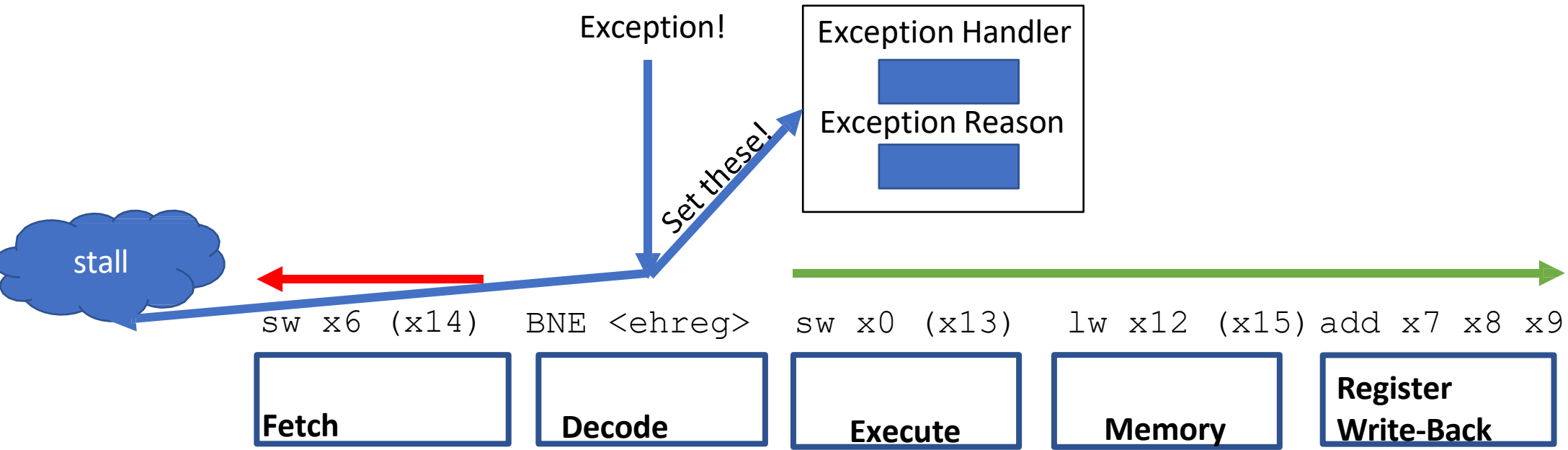
What if one of our instructions were to throw an exception (e.g., illegal instruction in decode or page fault on a memop)?

Exception Handling



Basic Exception Idea: Nuke everything that started after the current instruction, finish everything that started before the current instruction, jump to exception handler

Exception Handling



Basic Exception Idea: Nuke everything that started after the current instruction, finish everything that started before the current instruction, jump to exception handler, no new insns

What did we just learn?

- Basics of pipelining as a first technique for Instruction-level parallelism
- Datapath decomposition to support pipelined execution
- Hazards and their impediment to pipelined execution
- Forwarding in the pipeline to avoid stalling on data hazards

What to think about next?

- More microarchitectural concepts (next time)
 - Control hazards & branch prediction
- Caches as a microarchitectural optimization (next time)
 - Implementation of cache hierarchies
 - Cache design tradeoffs
- Performance Evaluation (next next time)
 - Design spaces, Pareto Frontiers, and design space exploration