



Collision resistance

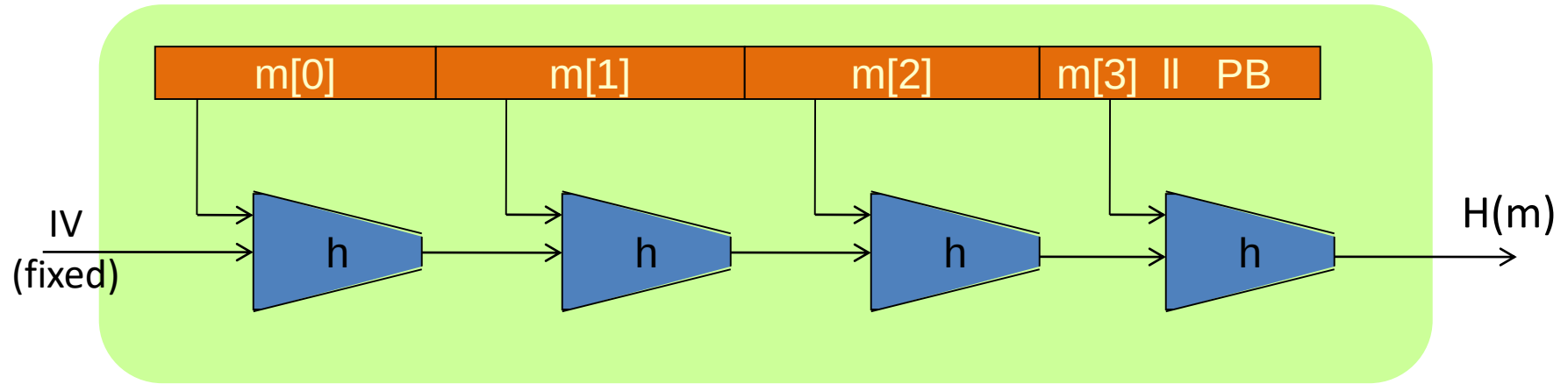


Collision resistance

HMAC:

a MAC from SHA-256

The Merkle-Damgård iterated construction



Thm: h collision resistant $\Rightarrow H$ collision resistant

Can we use $H(.)$ to directly build a MAC?

MAC from a Merkle-Damgard Hash Function

H: $X^{\leq L} \rightarrow T$ a C.R. Merkle-Damgard Hash Function

Attempt #1: $S(k, m) = H(k \parallel m)$

This MAC is insecure because:

- Given $H(k \parallel m)$ can compute $H(w \parallel k \parallel m \parallel \text{PB})$ for any w .
- Given $H(k \parallel m)$ can compute $H(k \parallel m \parallel w)$ for any w .
-  ○ Given $H(k \parallel m)$ can compute $H(k \parallel m \parallel \text{PB} \parallel w)$ for any w .
- Anyone can compute $H(k \parallel m)$ for any m .

Standardized method: HMAC (Hash-MAC)

Most widely used MAC on the Internet.

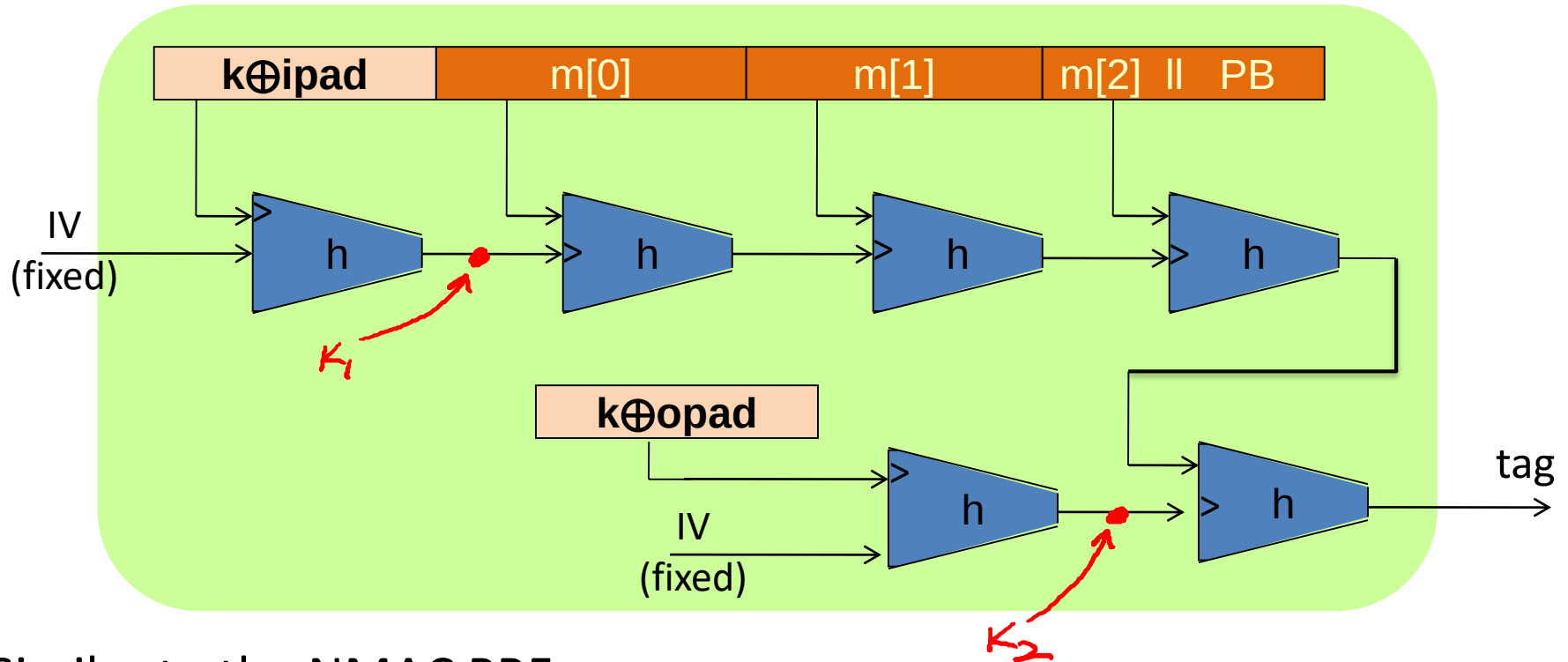
H: hash function.

example: SHA-256 ; output is 256 bits

Building a MAC out of a hash function:

$$\text{HMAC: } S(k, m) = H(k \oplus \text{opad} \parallel H(k \oplus \text{ipad} \parallel m))$$

HMAC in pictures



Similar to the NMAC PRF.

main difference: the two keys k_1 , k_2 are dependent

HMAC properties

Built from a black-box implementation of SHA-256.

HMAC is assumed to be a secure PRF

- Can be proven under certain PRF assumptions about $h(.,.)$
- Security bounds similar to NMAC
 - Need $q^2/|T|$ to be negligible ($q \ll |T|^{1/2}$)

In TLS: must support HMAC-SHA1-96

End of Segment



Collision resistance

Timing attacks on MAC
verification

Warning: verification timing attacks [L'09]

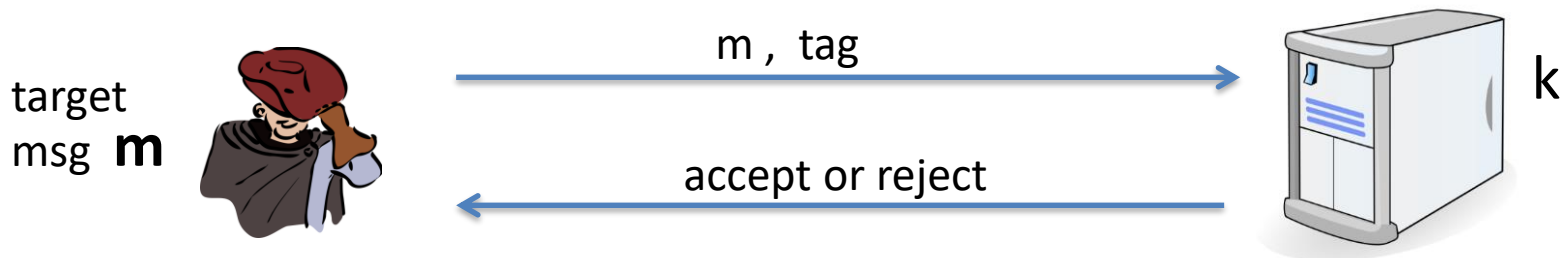
Example: Keyczar crypto library (Python) [simplified]

```
def Verify(key, msg, sig_bytes):  
    return HMAC(key, msg) == sig_bytes
```

The problem: '==' implemented as a byte-by-byte comparison

- Comparator returns false when first inequality found

Warning: verification timing attacks [L'09]



Timing attack: to compute tag for target message m do:

Step 1: Query server with random tag

Step 2: Loop over all possible first bytes and query server.

stop when verification takes a little longer than in step 1

Step 3: repeat for all tag bytes until valid tag found



Defense #1

Make string comparator always take same time (Python) :

```
return false if sig_bytes has wrong length  
result = 0  
for x, y in zip( HMAC(key,msg) , sig_bytes):  
    result = result | (ord(x) ^ ord(y))  
return result == 0
```

Can be difficult to ensure due to optimizing compiler.

Defense #2

Make string comparator always take same time (Python) :

```
def Verify(key, msg, sig_bytes):  
    mac = HMAC(key, msg)  
    return HMAC(key, mac) == HMAC(key, sig_bytes)
```

Attacker doesn't know values being compared

End of Segment